



INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

<p>(51) International Patent Classification ⁶ : H04L 29/06, 29/12</p>	A1	<p>(11) International Publication Number: WO 99/38303</p> <p>(43) International Publication Date: 29 July 1999 (29.07.99)</p>		
<table style="width: 100%; border: none;"> <tr> <td style="width: 50%; vertical-align: top; border: none; padding: 5px;"> <p>(21) International Application Number: PCT/US99/01195</p> <p>(22) International Filing Date: 21 January 1999 (21.01.99)</p> <p>(30) Priority Data: 60/072,175 22 January 1998 (22.01.98) US 09/057,578 9 April 1998 (09.04.98) US</p> <p>(71) Applicant (for all designated States except US): NORTHERN TELECOM LIMITED [CA/CA]; Station A, P.O. Box 6123, Montreal, Quebec H3C 3J5 (CA).</p> <p>(72) Inventors; and</p> <p>(75) Inventors/Applicants (for US only): BROTHERS, John, David, West [US/US]; 1620 Grand Junction, Alpharetta, GA 30004 (US). SMITH, Jeffrey, G. [US/US]; 12250 Stevens Creek Drive, Alpharetta, GA 30005 (US).</p> <p>(74) Agents: GARRETT, Arthur, S. et al.; Finnegan, Henderson, Farabow, Garrett & Dunner, L.L.P., 1300 I Street, Washington, DC 20005-3315 (US).</p> </td> <td style="width: 50%; vertical-align: top; border: none; padding: 5px;"> <p>(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, US, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).</p> <p>Published <i>With international search report. Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</i></p> </td> </tr> </table>			<p>(21) International Application Number: PCT/US99/01195</p> <p>(22) International Filing Date: 21 January 1999 (21.01.99)</p> <p>(30) Priority Data: 60/072,175 22 January 1998 (22.01.98) US 09/057,578 9 April 1998 (09.04.98) US</p> <p>(71) Applicant (for all designated States except US): NORTHERN TELECOM LIMITED [CA/CA]; Station A, P.O. Box 6123, Montreal, Quebec H3C 3J5 (CA).</p> <p>(72) Inventors; and</p> <p>(75) Inventors/Applicants (for US only): BROTHERS, John, David, West [US/US]; 1620 Grand Junction, Alpharetta, GA 30004 (US). SMITH, Jeffrey, G. [US/US]; 12250 Stevens Creek Drive, Alpharetta, GA 30005 (US).</p> <p>(74) Agents: GARRETT, Arthur, S. et al.; Finnegan, Henderson, Farabow, Garrett & Dunner, L.L.P., 1300 I Street, Washington, DC 20005-3315 (US).</p>	<p>(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, US, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).</p> <p>Published <i>With international search report. Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</i></p>
<p>(21) International Application Number: PCT/US99/01195</p> <p>(22) International Filing Date: 21 January 1999 (21.01.99)</p> <p>(30) Priority Data: 60/072,175 22 January 1998 (22.01.98) US 09/057,578 9 April 1998 (09.04.98) US</p> <p>(71) Applicant (for all designated States except US): NORTHERN TELECOM LIMITED [CA/CA]; Station A, P.O. Box 6123, Montreal, Quebec H3C 3J5 (CA).</p> <p>(72) Inventors; and</p> <p>(75) Inventors/Applicants (for US only): BROTHERS, John, David, West [US/US]; 1620 Grand Junction, Alpharetta, GA 30004 (US). SMITH, Jeffrey, G. [US/US]; 12250 Stevens Creek Drive, Alpharetta, GA 30005 (US).</p> <p>(74) Agents: GARRETT, Arthur, S. et al.; Finnegan, Henderson, Farabow, Garrett & Dunner, L.L.P., 1300 I Street, Washington, DC 20005-3315 (US).</p>	<p>(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, US, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).</p> <p>Published <i>With international search report. Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</i></p>			
<p>(54) Title: PROXY SERVER FOR TCP/IP NETWORK ADDRESS PORTABILITY</p> <div style="text-align: center; margin: 20px 0;"> <pre> graph LR 12[192.5.144.8] --- 18[Hub] 14[10.3.28.5] --- 18 16[47.31.128.195] --- 18 18 --- 20[Proxy Server] 20 --- 22((Router)) 10(()) --- 18 </pre> </div>				
<p>(57) Abstract</p> <p>System and method for fully transparent IP mobility services for clients in a dynamic LAN Ethernet environment. The functionality within a proxy server, a combination of network address translation, proxy address resolution protocol, and proxy domain name service, allow the proxy server to support and provide full IP client functionality to any IP-enabled network device in any proxy server enabled LAN. The proxy server may be added to an existing Ethernet (or Ethernet-emulated) LAN. Once configured with the necessary subnet range, DNS, and IP pools, the proxy server provides support for any mobile device that enters the LAN.</p>				

FOR THE PURPOSES OF INFORMATION ONLY

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece			TR	Turkey
BG	Bulgaria	HU	Hungary	ML	Mali	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MN	Mongolia	UA	Ukraine
BR	Brazil	IL	Israel	MR	Mauritania	UG	Uganda
BY	Belarus	IS	Iceland	MW	Malawi	US	United States of America
CA	Canada	IT	Italy	MX	Mexico	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NE	Niger	VN	Viet Nam
CG	Congo	KE	Kenya	NL	Netherlands	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NO	Norway	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	NZ	New Zealand		
CM	Cameroon	KR	Republic of Korea	PL	Poland		
CN	China	KZ	Kazakstan	PT	Portugal		
CU	Cuba	LC	Saint Lucia	RO	Romania		
CZ	Czech Republic	LI	Liechtenstein	RU	Russian Federation		
DE	Germany	LK	Sri Lanka	SD	Sudan		
DK	Denmark	LR	Liberia	SE	Sweden		
EE	Estonia			SG	Singapore		

PROXY SERVER FOR TCP/IP NETWORK ADDRESS PORTABILITY

Background of the Invention

The present invention relates generally to address portability and, more particularly,
5 to a method and apparatus for address portability to provide fully transparent internet
protocol (IP) mobility services to IP-enabled network devices in any Ethernet local area
network (LAN).

The transmission control protocol/internet protocol (TCP/IP protocol), a suite of
communications protocols used by host computers to exchange information between
10 application processes over LANs or wide area networks (WANs), was designed when
laptops and other mobile IP devices were essentially nonexistent. As a result, there was
no issue with mobility, since each IP network device was typically a workstation,
minicomputer, or the like. The movement of devices from place to place in such a static
environment was expected to be a very rare occurrence, and one that could be adequately
15 handled by manual intervention. This assumption, in conjunction with various resource
constraints, influenced the development of the IP protocol such that each LAN only
operated with a limited range (a subnet) of IP addresses. Any device with an IP address
outside of that range was simply ignored by the LAN's router, rendering it unable to
communicate with any device within that network.

20 Over the last several years, the IP protocol has become the primary data
communications protocol on virtually every computer in the world. This includes a
substantial number of laptops and other portable computer devices. As the prevalence of
laptops increases, IP mobility issues have substantially increased. For example, it is now
common for customers, vendors, and even business associates that have laptops or other
25 mobile IP devices to attempt to hook into a "foreign" LAN and attempt to use its facilities.
Typically, this results in significant frustration since the amount and complexity of
reconfiguration to permit the connection is not insubstantial.

One attempted solution to this problem, dynamic host configuration protocol
(DHCP), evolved over the last couple of years. Under DHCP, a computer configured to
30 use that protocol may retrieve local IP configuration data automatically when the mobile
IP device is connected to the network. While this is a reasonable solution to mobility
problems, its scope is somewhat limited. For example, the mobile network device must

- 2 -

be configured to use DHCP, and the LAN must have a DHCP server enabled. Moreover, the duration of DHCP "timeouts" within the mobile network device must be short enough to allow the device to request a new address at the new location. As a result of at least these limitations, DHCP has not sufficiently solved the problem. In some cases, DHCP
5 has proven unacceptable to the network clients who may not have DHCP pre-configured or to network administrators who wish to have more knowledge of and control over the mobile IP devices that enter and leave the network.

It is, therefore, desirable to provide fully transparent IP mobility services for clients in a dynamic network environment.

10 Summary of the Invention

Systems and methods consistent with the present invention satisfy this and other needs by supporting and providing full IP client functionality to any IP-enabled network device in any mobility-enabled LAN. The present invention provides full functionality regardless of both the IP address of the mobile device and subnet restrictions of the LAN.

15 A method for use with a proxy server consistent with the present invention establishes communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network. The method includes the step of generating an address resolution protocol packet to identify the arbitrary address for the destination device. The proxy server receives the address
20 resolution protocol packet and generates an address resolution protocol response packet including the arbitrary address of the destination device. The method also includes the step of transmitting the address resolution protocol response packet from the proxy server to the device in the first network.

Another method consistent with the present invention is for use with a proxy server
25 and establishes communications between a random device and a destination device having an arbitrary address on a network. The method includes the steps of generating an address resolution protocol packet to identify the arbitrary address for the destination device and receiving, by the proxy server, the address resolution protocol packet. The method also includes the steps of generating an address resolution protocol response packet including
30 the arbitrary address of the destination device and transmitting the address resolution protocol response packet from the proxy server to the random device.

- 3 -

Yet another method consistent with the present invention is for use with a proxy server which is in communication with a mobile device and remote name server. The method permits obtaining an internet protocol address from the remote name server for the mobile device and includes the steps of generating a query packet including a request for
5 an address associated with a domain name and receiving the query packet from the mobile device in the proxy server. The method also includes the steps of forwarding the query packet to the remote name server and generating a response packet including the requested address. The method also includes transmitting the response packet to the proxy server and transmitting the response packet to the mobile device.

10 Another method consistent with the present invention is for use with a proxy server and provides for communications between a random device and a destination device in a network. The method includes the steps of performing a proxy address resolution protocol to initiate communications between the random device and the destination device, performing a proxy domain name service to identify a destination address in the second
15 network associated with a domain name, and performing a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the destination device in the network. Use of this combination allows a system to support and provide full client functionality to mobile network devices.

20 Systems are also provided for carrying out these and other methods consistent with the present invention.

Several advantages accrue to methods and systems consistent with the present invention. For example, these systems and methods provide a secure and complete mobility solution, including the various cases where prior art solutions were inadequate.
25 Such systems and methods are completely transparent to the end-user, who may or many not use DHCP, but will still be able to communicate with a LAN or even with a WAN. They are also more "administrator-friendly", especially when the acceptance protocol involves e-mail notification to the network administrator that a new device has joined the network. Security is enhanced by reducing the network's exposure to foreign snooping.

30 The above and additional features and advantages of the present invention will be readily appreciated by one of ordinary skill in the art from the following detailed description.

- 4 -

Brief Description of the Drawings

Figure 1 is a diagram of a random Ethernet LAN environment consistent with the present invention;

Figure 2 is a high level system diagram of a proxy server consistent with the present invention;

Figures 3 and 4 illustrate network address translation associated with the routing of traffic from a random LAN to a legal LAN consistent with the present invention;

Figures 5 and 6 illustrate network address translation associated with the routing of traffic from a legal LAN to a random LAN consistent with the present invention;

Figures 7 and 8 illustrate generation of a proxy address resolution protocol (ARP) packet and generation of a proxy ARP response packet consistent with the present invention;

Figures 9-12 are flowcharts depicting steps for proxy ARP consistent with the present invention;

Figures 13 and 14 illustrate generation of a proxy domain name service (DNS) query packet and generation of a proxy DNS response packet consistent with the present invention;

Figures 15 and 16 are flowcharts depicting steps for proxy DNS consistent with the present invention;

Figure 17 illustrates an alternative proxy server implementation consistent with the present invention;

Figure 18 illustrates normal traffic flow in the alternative proxy server implementation of Figure 16;

Figures 19-22 illustrate network address translation for use with the alternative proxy server implementation consistent with the present invention; and

Figures 23-24 illustrate proxy ARP for use with the alternative proxy server implementation consistent with the present invention.

Detailed Description of the Preferred Embodiments

Reference will now be made in detail to embodiments consistent with this invention that are illustrated in the accompanying drawings. The same reference numbers in different drawings generally refer to the same or like parts.

Two-Armed Proxy Server

- 5 -

Figure 1 shows a "random" LAN 10. For purposes of this discussion, "random" simply means there are a random number of mobile IP network devices, and those devices each use a random IP address. One example of a "random" LAN would be an IP Ethernet network in a hotel. LAN 10 includes a plurality of interconnected mobile IP network devices, including laptops 12, 14, and 16, having IP addresses scattered across the range of known addresses. As shown, the plurality of network devices communicate through a hub 18 and a proxy server 20 with network router 22. The links used to interconnect the various network elements shown may, for example, be Ethernet links. In the LAN 10, proxy server 20 may be referred to as a "two-armed" (TA) proxy server since it possesses two network interfaces, *e.g.*, two Ethernet links.

Normally, this type of network would be extremely difficult to manage, since a standard router expects all IP addresses that it serves to fall within a limited range. Consistent with the present invention, however, traffic from each of these devices may be modified so that the information presented to the network router is acceptable. The modification may be accomplished by proxy server 20 using a combination of network address translation (NAT), a proxy address resolution protocol (ARP) service, and a proxy domain name system (DNS) service, as discussed below. NAT is a well-known process by which traffic received by and transmitted from a particular device with an arbitrary IP address is modified to present the correct IP address to a network router. NAT service may be specifically configured to translate particular IP addresses. A proxy mobility server consistent with the present invention may translate random IP addresses dynamically.

Figure 2 illustrates a high level diagram of proxy server 20. As shown, proxy server 20 includes a processor in communication with a hard drive, a system memory, and a user memory. The system and user memories may include read-only and/or random access types of memories. These memories are useful for storing packet contents, which may include addresses and the like, as well as data content and packet length, to name a few. Proxy server 20 also includes interfaces, which may take the form of cards, through which proxy server communicates with networks. In the example shown, proxy server 20 is interfaced with a legal/public network and a random network through Ethernet interface 0 and Ethernet interface 1, respectively.

Figures 3 and 4 depict the routing of information, such as an IP packet 26, from a device within a random LAN 28 to a device within a "legal" LAN 30. For purposes of this

- 6 -

discussion, a "legal" LAN is simply a public LAN, *i.e.*, one with a legal set of IP addresses. In Figure 3, packet 26 from the device, which has an IP address of 47.31.128.195, is routed from random LAN 28 to proxy server 20. As shown, proxy server 20 is a server for a subnet (here denoted 137.118.199.X) which, as is known, is a set of machines that are
5 physically connected together in an Ethernet LAN, *e.g.*, the legal LAN 30 in Figure 3. Proxy server 20 performs a network address translation and, as shown in Figure 4, the translated packet 32 is transmitted to the legal LAN. The translation may be performed using known techniques, such as those specified in Internet Engineering Task Force Request for Comments (IETF RFC) 1631. A packet following the opposite path, *i.e.*,
10 packet 34 routed from a device within legal LAN 30 to a device within random LAN 28, would also undergo network address translation (to translated packet 36) in proxy server 20 (see Figures 5 and 6).

Proxy ARP

In addition to NAT, proxy server 20 may employ a proxy address resolution
15 protocol (ARP) service to provide mobile functionality consistent with the present invention. ARP is a known protocol which may be used by a network device to discover what other devices are connected to the local network. Proxy ARP service allows TA proxy server 20 to "spoof" mobile network IP devices having random IP addresses into thinking that server 20 is the device with which those mobile IP devices wish to
20 communicate. This is necessary when the mobile IP device first boots and attempts to determine its gateway. As is known, existing proxy ARP implementations are limited in their use since only traffic from certain select specific addresses can be handled. Proxy ARP consistent with the present invention is not so limited and may be used to identify any arbitrary address.

25 Figure 7 illustrates an ARP packet 38 being transmitted by a mobile device in random LAN 28 to proxy server 20. As shown, ARP packet 38 includes the address of the sending device, as well as a query from the device regarding the whereabouts of its gateway, which has the address of the gateway sought. The gateway, which may be a network router, connects the smaller LAN (*e.g.*, a random LAN) to a larger WAN (*e.g.*,
30 a public LAN, such as the "legal" LAN 30 of Figure 7) and passes traffic from the LAN to the WAN. When a mobile device in the random LAN wishes to send traffic to a device having an arbitrary address in a second network (*e.g.*, the WAN) outside of the LAN, the

- 7 -

sending device needs to know to which gateway device the traffic should be sent. While normally the gateway is on the same network as the mobile device that is searching for it (using the ARP), this is not possible in a random LAN. Accordingly, the proxy server pretends that it is the gateway, and the mobile device will use it to reach the WAN. In
5 response to receiving ARP packet 38, proxy server 20 generates an ARP response packet 40 destined for the sending device in random LAN 28. As shown in Figure 8, response packet 40 includes the sending device's IP source address as the destination address and informs the sending device that proxy server 20 is the device's gateway.

Figure 9 depicts steps for proxy ARP consistent with the present invention. During
10 initialization (step 80), IP addresses and network masks are determined, as shown in greater detail in Figure 10. First, a raw socket is created to examine ARP packets (step 100). This socket is a communications programming interface created between proxy server 20 and the random and public LANs. In the Redhat Linux operating system, only one such socket is needed, whereas other operating systems, such as Sun Microsystem's
15 Solaris operating system, must create one socket per Ethernet interface. Next, IP network masks for both the public LAN interface and the random LAN interface are identified (steps 102, 104). As is known, these interfaces may be cards in proxy server 20. This information, used by the proxy server in combination with other information to determine what set of devices are part of the networks and, therefore, whether it can send packets
20 directly to any mobile device or whether the packets must be sent through a router, is typically maintained in a stable storage device, such as hard disk, flash memory, and the like, in proxy server 20. Similarly, IP addresses of the public and random LAN interfaces are identified (steps 106, 108), as is the medium access control (MAC) address of the random interface (step 110). Like the IP network masks, these addresses are typically
25 stored or built into proxy server 20.

With continuing reference to Figure 9, when a new ARP packet from a random IP device, such as ARP packet 38 of Figure 7, arrives at the random network interface from a mobile IP device in the random LAN (step 82), proxy server 20 retrieves the packet contents from the operating system (step 84). The packet, as is known, has a header and
30 data, which includes *inter alia* the packet source address (*i.e.*, the address of the mobile device that sent the ARP packet) and the packet destination address (*i.e.*, the address of that device's gateway). The server then applies the ARP data format to the IP packet (step 86).

The ARP data format is defined in IETF Standard 37, and the application of the format to the data may be done using the standard method of casting.

Next, the proxy server performs a proxy ARP network determination to determine the network to which the IP packet is destined for (step 88). Figure 11 shows a flowchart
5 detailing steps for this determination consistent with the present invention. A public network ID (PubNetID) is determined first (step 120). In one embodiment, the public network ID is derived from the public IP interface address and the IP network mask of the public LAN, *e.g.*, logically “anding” the address with the mask. The proxy server uses the PubNetID to discover what network it is part of, *i.e.*, what IP devices are local and which
10 are not local. Devices that not local are reached through a router. Next, the proxy server determines if the IP destination address of the incoming packet is for a “local” network, *i.e.*, the public LAN. In one embodiment, a network ID associated with the incoming packet (NewNetID) is determined (step 122) based on the destination IP address sub-component of the ARP data structure associated with the ARP packet and the IP network
15 mask of the public LAN interface, *e.g.*, logically “anding” the destination address with the public IP network mask. If PubNetID is equal to NewNetID (step 124), the incoming packet is destined for a device that has an IP address on the public LAN but is physically part of or on the random LAN. The proxy server discards this incoming packet (step 126) because another device in the random LAN will receive the packet by Ethernet. The proxy
20 server simply ignores the packet because it does not need to create a response packet; the intended device physically on the random LAN should respond.

If the incoming packet is not destined for the public LAN, a random network ID (RandNetID) may be determined based on the random IP address and the IP network mask of the random LAN, *e.g.*, logically “anding” the address with the mask (step 128). The
25 proxy server uses RandNetID to discover what devices are local to the random LAN. Devices that not local are reached through a router. A network ID associated with the packet (NewNetID2) is also determined (step 130). This may be determined based on the destination IP address sub-component of the ARP data structure associated with the ARP packet and the IP network mask of the random LAN interface, *e.g.*, logically “anding” the
30 destination address with the random IP network mask. The proxy server uses this information to determine if the IP destination address of the ARP packet is for a “local” network, *i.e.*, the random LAN. RandNetID may then be compared to NewNetID2 (step

132). If RandNetId is equal to NewNetID2, the incoming packet is destined for a device that has an IP address on the random LAN and physically part of or on the random LAN. Again, the proxy server discards the incoming packet (step 134), since another device in the random LAN should respond and will receive the packet by Ethernet.

5 If the incoming packet has not been discarded by the proxy server based on these comparisons, the packet is destined for a device outside of a local network, *i.e.*, from a device in a random LAN, such as a hotel, to a device outside of that LAN and physically part of, for example, a public LAN. Consistent with the present invention, the proxy server prepares an ARP response packet (step 136) to the incoming packet to convince the
10 sending device that the proxy server is the device's gateway. The response ARP packet format is defined in IETF Standard 37.

Referring once again to Figure 9, once an ARP response packet has been created, a proxy ARP address exchange is performed (step 90), as shown in Figure 12. Consistent with the present invention, the IP source address component from the incoming packet is
15 copied into the destination address component of the response packet (step 140). This directs the response packet to the appropriate mobile IP device. Similarly, the source MAC address component of the incoming packet is copied to the destination MAC address of the response packet (step 142). The destination address component of the incoming packet is copied into the source address component of the response packet (step 144). Address
20 exchange consistent with the present invention also contemplates filling in the source MAC address component of the response packet with the MAC address of the random interface (step 146). By inserting the operation component of the response packet with the appropriate value in network-byte order (*e.g.*, the value "2" for RedHat Linux 5.0), the packet is considered a response packet for purposes of the ARP protocol.

25 Once the address exchange is completed, the response packet may be written to the random Ethernet interface using a standard system call (Figure 9, step 92). When the mobile IP device receives the response packet, it will believe proxy server 20 is its gateway from the random LAN to outside public LANs. Subsequent traffic destined for public LANs will be routed there by the proxy server.

30 Proxy DNS

Normally, a mobile network device communicates with a nearby network element commonly referred to as a DNS server. The DNS server functions to translate an IP name

- 10 -

input by a user, such as "undefined.etherloop.com," into a corresponding IP address, such as 137.118.199.33. Thus, when a mobile IP device user inputs an IP name as an intended destination, the device communicates with the DNS server, which then performs a translation, *e.g.*, a name lookup. However, with a random LAN, this is not possible.

5 Figure 13 shows the transmission of a DNS packet 42 from a mobile IP device within random LAN 28 to proxy server 20. Consistent with the present invention, the proxy server pretends that it is the correct DNS server and handles the DNS translation activities. In general, DNS packet 42 includes the IP address of the sending mobile IP device, a destination address, *i.e.*, the address of the DNS server, which is typically
10 preconfigured in the device, and the IP name the device is seeking an IP address for. Proxy server 20 redirects the request to a local process, *i.e.*, a process within the proxy server which, in turn, performs the required translation. After the translation is performed, proxy server 20 generates a DNS response packet 44 (see Figure 14) which includes the address of the sending network device as a destination address and the IP address corresponding
15 to the IP name.

Figure 15 depicts steps for proxy DNS consistent with the present invention. First, the proxy DNS server is initialized (step 150). As shown in Figure 16, during initialization, proxy server 20 creates an unreliable datagram protocol (UDP) datagram socket (step 170) using known methods. This socket is a communications interface
20 between the local process (proxy DNS) and the operating system. Proxy server 20 also establishes a firewall rule such that any packet headed for any destination in any LAN or WAN (*i.e.*, any IP address) from any source in any LAN or WAN with a destination port of "53" is delivered to the proxy server 20 at the same port, *i.e.*, port "53" of the server (step 172). Port "53" is the port for the DNS server of the host machine (a machine with
25 an IP address). The socket can then be bound to port "53" using standard methods (step 174), such that any DNS query with a destination of port "53" is routed to the local process (proxy DNS). The identity of a remote name server physically located outside the random LAN on the Internet is also identified by the proxy server using, for example, a configuration file or some other known method (step 176). After creating a UDP socket
30 association between the proxy DNS server and the identified remote name server (step 178), proxy server 20 enters a loop waiting for new connections, such as new DNS queries

from the random LAN or DNS responses from the remote name server, using standard methods (step 180).

After proxy DNS initialization, a "random" client (*i.e.*, a device in the random LAN having an IP address, such as 1.1.1.1) makes a DNS query to the identified remote name
5 server (*e.g.*, having an IP address, such as 2.2.2.2) at port "53" (step 152) using standard DNS protocol, *e.g.*, IETF standard 13. Proxy server 20 receives the packet and, based on the port address used as the destination port (*i.e.*, port "53"), redirects the packet to the proxy DNS server (the local process within the proxy server) (step 154) using, for example, firewall redirection code built into Linux Redhat 5.0. The proxy DNS server in turn
10 receives the packet and determines the original destination address (*i.e.*, 2.2.2.2) and port (*i.e.*, port "53") of the intended destination, storing them in appropriate variables. In Redhat Linux 5.0, this may be accomplished using an appropriate system call to collect the packet data from the kernel.

The proxy DNS server may then send the DNS query packet to the remote name
15 server identified during initialization (step 158). The proxy DNS server creates a new UDP socket, a communications interface between it and the remote name server. Typically, the proxy DNS server uses a separate socket and port (such as port "2001," in this example) which may be arbitrarily assigned) for each mobile device IP address so as to be able to identify to which device in the random LAN the response should be sent. The remote
20 name server then processes and responds to the DNS query from the proxy DNS server, as defined in IETF Standard 13 (step 160). A DNS response packet, which includes the requested address, is generated by the remote name server and sent to the proxy server using the port defined in step 158 (*i.e.*, port "2001" in this example). Proxy server (such as proxy server 20) receives the DNS response from the remote name server on the
25 specified port (*i.e.*, port "2001") using, for example, standard Unix system calls and determines the IP address of the client (*i.e.*, 1.1.1.1 in this example) using that port (step 164). The DNS response is then sent to the client by the proxy server (step 166; see also Figure 14), which performs a source address and port modification. In one embodiment, the proxy server modifies the source address (the address of the proxy server, *e.g.*, 3.3.3.3.
30 in this example) and source port of the response packet (the port of the link back to the client, which may be arbitrarily assigned) to the original destination address (*i.e.*, 2.2.2.2.

- 12 -

in this example) and the original destination port (*i.e.*, port "53" in this example) of the DNS query packet, respectively.

One-Armed Proxy Server

The TA proxy server described above is extremely well-suited for operation in any
5 environment where a fully random assortment of users may attempt to connect to the LAN. As previously noted, one such environment is a hotel environment. Today, hotel guests frequently have mobile network devices, such as laptops, and wish to connect via Ethernet, EtherLoop, or the like, into the hotel's network and from there to the Internet to retrieve electronic mail and conduct other business.

10 A TA proxy server, however, may not fit well into a LAN networking community since most end-user network devices will have a stable IP address, correctly configured and assigned for the LAN, unlike the hotel environment. Moreover, a TA proxy server increases latency (delay) on a network. For at least these reasons, there is a need for proxy routers that do not interfere with the normal operation of a LAN. These services will exist
15 on a normally configured LAN and only begin operation only when a random interloper, *i.e.*, an individual device, connects to the network. The device that supports this service will be lower in cost and will not create any performance problems for the standard network traffic.

Figure 17 illustrates an alternative proxy server implementation consistent with the
20 present invention that satisfies this need. In this implementation, LAN 200 includes a plurality of mobile IP network devices 202, 204, and 206 which communicate with a router 210 through hub 212. The links used to interconnect the various network elements shown may, for example, be Ethernet links. Proxy server 214 communicates with the various network elements through a single link, *e.g.*, a single Ethernet interface. As such, proxy
25 server 214 may be referred to as one-armed proxy server. Generally, proxy server 214 includes the same hardware as proxy server 20 (Figure 2). The same Ethernet interface receives data from LAN 200 and delivers translated data back to the LAN. In Figure 17, the majority of the devices on the network are on the same network as the router. These devices operate normally without any interference from proxy server 214, which is always
30 listening to the traffic on this network. Device 208, an interloper with an address of 47.31.128.195 will not be able to directly communicate with the router, which ignores all traffic from an address that is not in its network domain (*i.e.*, 137.118.199.X). One of the

- 13 -

benefits of this implementation is that the standard network traffic is not interrupted by the proxy server, as shown in Figure 18. As with the TA proxy server of Figure 1, one-armed proxy server 214 performs network address translation, proxy ARP, and proxy DNS services to provide similar fully transparent client functionality to any IP-enabled network device.

Network Address Translation

Figure 19 depicts a network environment useful for discussing NAT performed by one-armed proxy server 64. Network 220, an Ethernet type of network, includes a plurality of interconnected network elements including router 222, workstation 224, hub 226, and proxy server 214. Interloper 228 is also connected to LAN 220, although it is a foreign IP-enabled network device relative to the network, *i.e.*, it is not part of LAN 220.

As shown in Figure 19, interloper 228 generates a packet 230 for a LAN/WAN other than the LAN 220. Packet 230 includes the IP address of the transmitting device, interloper 228, as well as an IP address of the destination device, which is not particularly shown. Due to the nature of Ethernet networks, every element on the LAN receives packet 230 from hub 226. The IP protocols within workstation 224 and similar network elements recognize that the destination is not a local network device and consequently ignore the packet. Router 222 may or may not accept packet 230 depending on the source IP address. Even if the router accepts the packet and passes it on to the LAN/WAN, the packet will not return to the router (address 137.118.199.1) since the source address is 47.31.128.195.

Proxy server 214, however, recognizes that it is capable of properly translating packet 230 into translated packet 232 having an acceptable format utilizing known address translation methods (see Figure 20). Since translated packet 232 has an associated IP address recognizable by router 222, the router believes packet 232 originated from within LAN 220 and not by a foreign mobile device, *i.e.*, interloper 228. As shown, packet 232 includes the same destination address as packet 230. Once router 222 receives this packet, it can send the packet on to the Internet as normal.

After the remote device receives the translated packet, it may generate a response packet. If so, the response packet 234 must be translated by proxy server 214 so that it can be delivered to interloper 228 (see Figure 21). Router 222 knows that the .55 IP address is associated with a device on its network. In this case, the device happens to be proxy server 214, but the router is unaware of the presence of the server nor does it matter to the

- 14 -

router that the destination is the proxy server and not a "normal" network device such as a workstation. Instead, router 222 simply forwards response packet 234 to proxy server 214, just like it would forward any other packet to other network devices. At about the same time the proxy server receives response packet 234, interloper 228 also receives the response packet from hub 226. However, since the interloper 228 knows that its IP address is 47.31.128.195, it throws the response packet away. Consistent with the present invention, proxy server 214 receives the response packet, performs a reverse network address translation, and sends a translated packet 236 back out on the LAN through hub 226 (see Figure 22). Packet 236 is broadcast across the entire LAN, but since only one device on the network has IP address 47.31.128.195 (interloper 228), only that device will not discard the packet. Proxy server is able to specifically target the interloper by using the interloper's medium access control (MAC) address as the destination.

Proxy ARP

Figures 23 and 24 show a network environment useful for discussing the proxy ARP capabilities of proxy server 214. Interloper 228 may generate an ARP packet 238 in order to discover the MAC Address of its gateway device, *i.e.*, 47.31.128.1. Normally, since no device in LAN 220 has the appropriate MAC address, packet 238 would be ignored and interloper 228 would be unable to function. Consistent with the present invention, proxy server 214 will however recognize that packet 238 does not belong on the 137.118.199.X network and will automatically generate a response (see Figure 24). Response packet 240 includes, as a destination address, the IP address of interloper 228, as well as a reply to the gateway query. Once interloper 228 receives packet 240, it considers the proxy server 214 to be its gateway device and will use the server for all further communications outside of the local LAN. The steps for proxy ARP performed by TA proxy server 20 discussed above are equally applicable to proxy server 214.

Proxy DNS

Proxy server 214 is also capable of performing proxy DNS. If, from the point of view of the interloper, the DNS server is usually on the same LAN as the interloper, the interloper will generate an ARP request for the DNS server. Consistent with the present invention, the proxy server 214 will respond to the ARP request with its own address. Future DNS queries will be delivered directly to proxy server 214, which can then answer them. Similarly, if, from the point of view of the interloper, the DNS server is outside of

- 15 -

the local LAN, *i.e.*, on a WAN, the proxy server will automatically receive the DNS query, since it is the interloper's gateway. In this case, proxy serve 214 will see the DNS query packets arrive and will be able to response to them locally. The steps for proxy DNS performed by TA proxy server 20 discussed above are equally applicable to the proxy DNS
5 service performed proxy server 214.

In addition to the features described above, proxy servers consistent with the present invention support certain security functions to improve network administration. For example, each time an interloper connects to an proxy server-enabled network, the proxy server will be able to provide connectivity for that user. To improve the security of
10 the network, the proxy server will deliver a message to a specified network administrator e-mail account to the alert the administrator to the presence of this new user. While this is not ironclad security, it is a reasonable first step in network security.

Properly configured, proxy servers consistent with the present invention can provide interlopers with a secondary gateway. This has at least two benefits, including
15 reduced congestion on the standard router and improved control over the interloper's internet access. Reduced congestion is a relatively straightforward concept, *i.e.*, by using a different router than the standard network traffic, it reduces the possibility of excessive demand on router resources, that in turn might affect the performance of the standard network users. Further, by specifying a secondary gateway, the network administrator can
20 funnel interlopers into a less open corporate environment, preventing those users from reaching sensitive material within the standard corporate network. This is one of the major benefits of the present invention over the straight DHCP model, for two reasons. First, if the network is served by a proxy server, a stand-alone DHCP server is not needed. The standard network users will not need to use DHCP on a day-to-day basis. Given the first
25 constraint, everyone who uses DHCP is, by implication, an interloper and can be treated with additional security restrictions.

As for other protocols, one of ordinary skill will appreciate that proxy servers consistent with the present invention are only capable of supporting IP-based translation services. This is primarily because IP is both a LAN and WAN protocol. Other common
30 LAN protocols, such as Apple Talk, or IPX are significantly limited in scope, and are not capable of the "long range" communications that make the proxy server translation services possible. However, it may be possible for "bridges" to be built to allow IPX-based

- 16 -

computers to communicate with their "home" networks. However, this must currently be resolved on a case-by-case basis.

It will be appreciated by those skilled in this art that various modifications and variations can be made to the IP mobility service strategy consistent with the present invention described herein without departing from the spirit and scope of the invention. Other embodiments of the invention will be apparent to those skilled in this art from consideration of the specification and practice of the invention disclosed herein. It is intended that the specification and examples be considered exemplary only, with a true scope and spirit of the invention being indicated by the following claims.

- 17 -

We claim:

1. A method, for use with a proxy server, for establishing communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network, the method comprising the steps of:
 - 5 generating an address resolution protocol packet to identify the arbitrary address for the destination device;
 - receiving, by the proxy server, the address resolution protocol packet;
 - generating an address resolution protocol response packet including the arbitrary address of the destination device; and
 - 10 transmitting the address resolution protocol response packet from the proxy server to the device in the first network.
2. The method of claim 1 further comprising the step of:
 - 15 determining if the proxy server should respond to the address resolution protocol packet.
3. The method of claim 2 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:
 - 20 determining if the address resolution protocol packet is destined for a device that has an address outside of the first network but that is part of the first network; and
 - discarding the address resolution protocol packet if it is destined for a device that has an address outside of the first network but that is part of the first network.
4. The method of claim 3 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:
 - 25 determining if the address resolution protocol packet is destined for a device that has an address on the first network and that is part of the first network; and
 - discarding the address resolution protocol packet if it is destined for a device that has an address on the first network and that is part of the first network.
5. The method of claim 4 wherein the step of generating an address resolution protocol response packet includes the substep of:

- 18 -

generating a response packet if the address resolution protocol packet is destined for a device that is outside of the first network.

6. The method of claim 5 further comprising the step of:
5 performing an address exchange to direct the response packet to the appropriate device.

7. The method of claim 6 wherein the step of performing an address exchange includes the substeps of:
10 writing an internet protocol source address component of the address resolution protocol packet into a destination address component of the response packet;
writing a source MAC address component of the address resolution protocol packet into a destination MAC address component of the response packet; and
writing a MAC address associated with the first network into a source MAC
15 address component of the response packet.

8. A system for establishing communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network, the system comprising:
20 a proxy server in communication with the first network and the second network, the proxy server including:
memory for receiving an address resolution protocol packet generated by the device in the first network; and
a processor, in communication with the memory, for determining if the
25 proxy server should respond to the address resolution protocol packet with an address resolution protocol response packet, the processor also for generating the response packet and transmitting the response packet to the device in the first network.

9. A method, for use with a proxy server, for establishing communications
30 between a random device and a destination device having an arbitrary address on a network, the method comprising the steps of:

- 19 -

generating an address resolution protocol packet to identify the arbitrary address for the destination device;

receiving, by the proxy server, the address resolution protocol packet;

generating an address resolution protocol response packet including the arbitrary
5 address of the destination device; and

transmitting the address resolution protocol response packet from the proxy server to the random device.

10. The method of claim 9 further comprising the step of:

10 determining if the proxy server should respond to the address resolution protocol packet.

11. The method of claim 10 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:

15 determining if the address resolution protocol packet is destined for a device that has an address associated with the network but that is physically not part of the network; and

discarding the address resolution protocol packet if it is destined for a device that has an address associated with the network but that is physically not part of the network.

20

12. The method of claim 11 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:

determining if the address resolution protocol packet is destined for a random device that is not part of the network; and

25 discarding the address resolution protocol packet if it is destined for a random device that is not part of the network.

13. The method of claim 12 wherein the step of generating an address resolution protocol response packet includes the substep of:

30 generating a response packet if the address resolution protocol packet is destined for a device that is part of the network.

- 20 -

14. The method of claim 13 further comprising the step of:
performing an address exchange to direct the response packet to the appropriate device.

5 15. The method of claim 14 wherein the step of performing an address exchange includes the substeps of:

writing an internet protocol source address component of the address resolution protocol packet into a destination address component of the response packet;

writing a source MAC address component of the address resolution protocol packet
10 into a destination MAC address component of the response packet; and

writing a MAC address associated with the random device into a source MAC address component of the response packet.

16. A system for establishing communications between a random device and
15 a destination device having an arbitrary address on a network, the system comprising:

a proxy server in communication with the random device and the destination device, the proxy server including:

memory for receiving an address resolution protocol packet generated by the random device; and

20 a processor, in communication with the memory, for determining if the proxy server should respond to the address resolution protocol packet with an address resolution protocol response packet, the processor also for generating the response packet, and transmitting the response packet to the random device.

25 17. A method, for use with a proxy server, for communicating with a mobile device and remote name server and for obtaining an internet protocol address from the remote name server for the mobile device, the method comprising the steps of:

generating a query packet including a request for an address associated with a domain name;

30 receiving the query packet from the mobile device in the proxy server;

forwarding the query packet to the remote name server;

generating a response packet including the requested address;

- 21 -

transmitting the response packet to the proxy server; and
transmitting the response packet to the mobile device.

18. The method of claim 17 wherein the query packet includes a destination
5 port and wherein the step of forwarding the query packet to the remote name server
includes the substeps of:

directing the query packet to a proxy DNS server based on the destination port;
determining, by the proxy DNS server, the address of the remote name server and
the destination port; and
10 identifying a port associated with the mobile device.

19. The method of claim 18 wherein the step of transmitting the response
packet to the proxy server includes the substeps of:

receiving the response packet on the port associated with the mobile device; and
15 determining, by the proxy server, the address of the mobile device based on the port
associated with the mobile device.

20. The method of claim 19 wherein the response packet includes a source
address and a source port, and wherein the step of transmitting the response packet to the
20 mobile device includes the substeps of:

modifying the source address of the response packet to the address of the remote
name server; and

modifying the source port of the response packet to the destination port.

25 21. A system for obtaining an internet protocol address from a remote name
server for a mobile device, the system comprising:

a proxy server in communication with the mobile device and the remote name
server, the proxy server including:

memory for receiving an address query packet generated by the mobile
30 device; and

- 22 -

a processor, in communication with the memory, for forwarding the query packet to the remote name server and for transmitting a response packet, generated by the remote name server and including the requested address, to the mobile device.

5 22. A method, for use with a proxy server, for providing communications between a random device and a destination device in a network, the method comprising the steps of:

performing a proxy address resolution protocol to initiate communications between the random device and the destination device;

10 performing a proxy domain name service to identify a destination address in the second network associated with a domain name; and

performing a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the destination device in the network.

15

23. A system for providing communications between a random device and a device in a network, the system comprising:

a proxy router including

memory; and

20

a processor programmed to

(a) perform a proxy address resolution protocol to initiate communications between the random device and the device in the network;

(b) perform a proxy domain name service to identify a destination address in the network based on a destination name; and

25

(c) perform a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the device in the network.

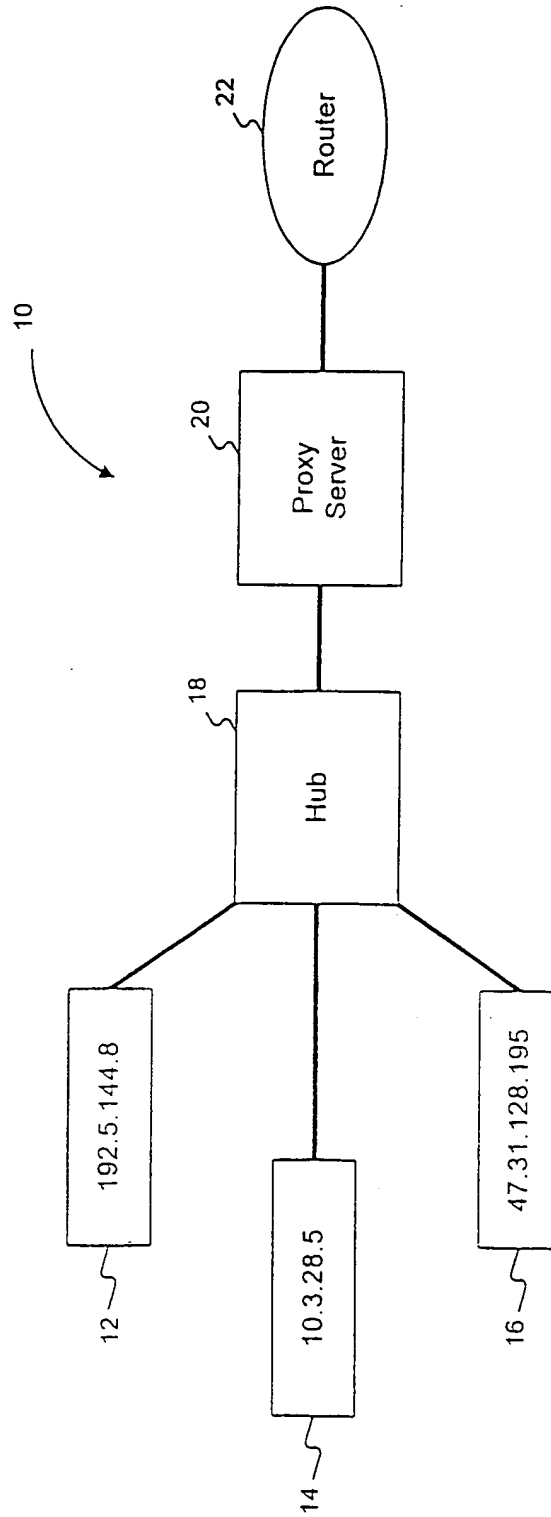


Figure 1

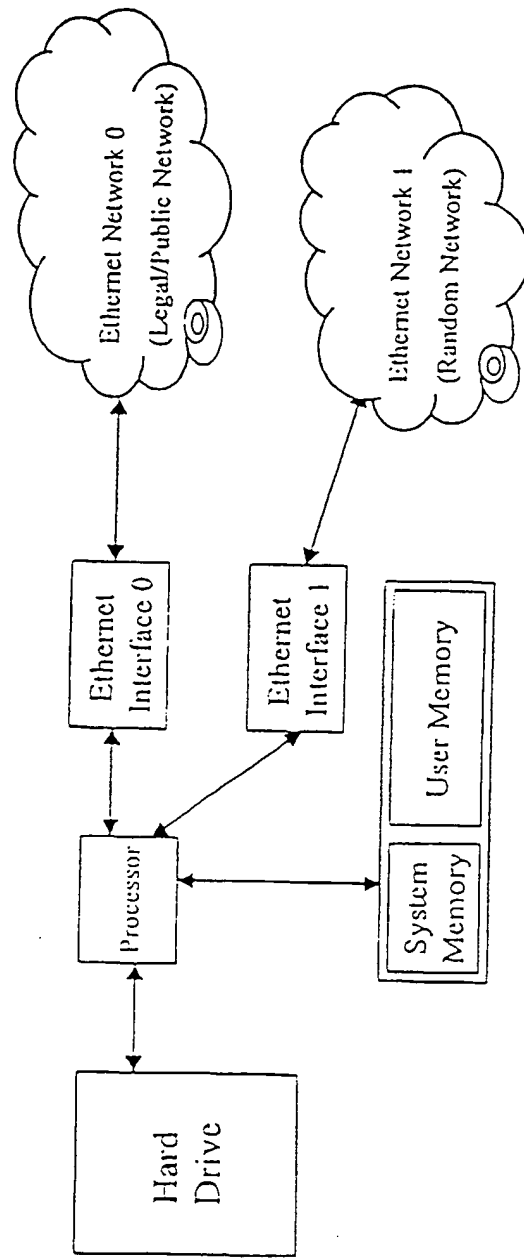


Figure 2

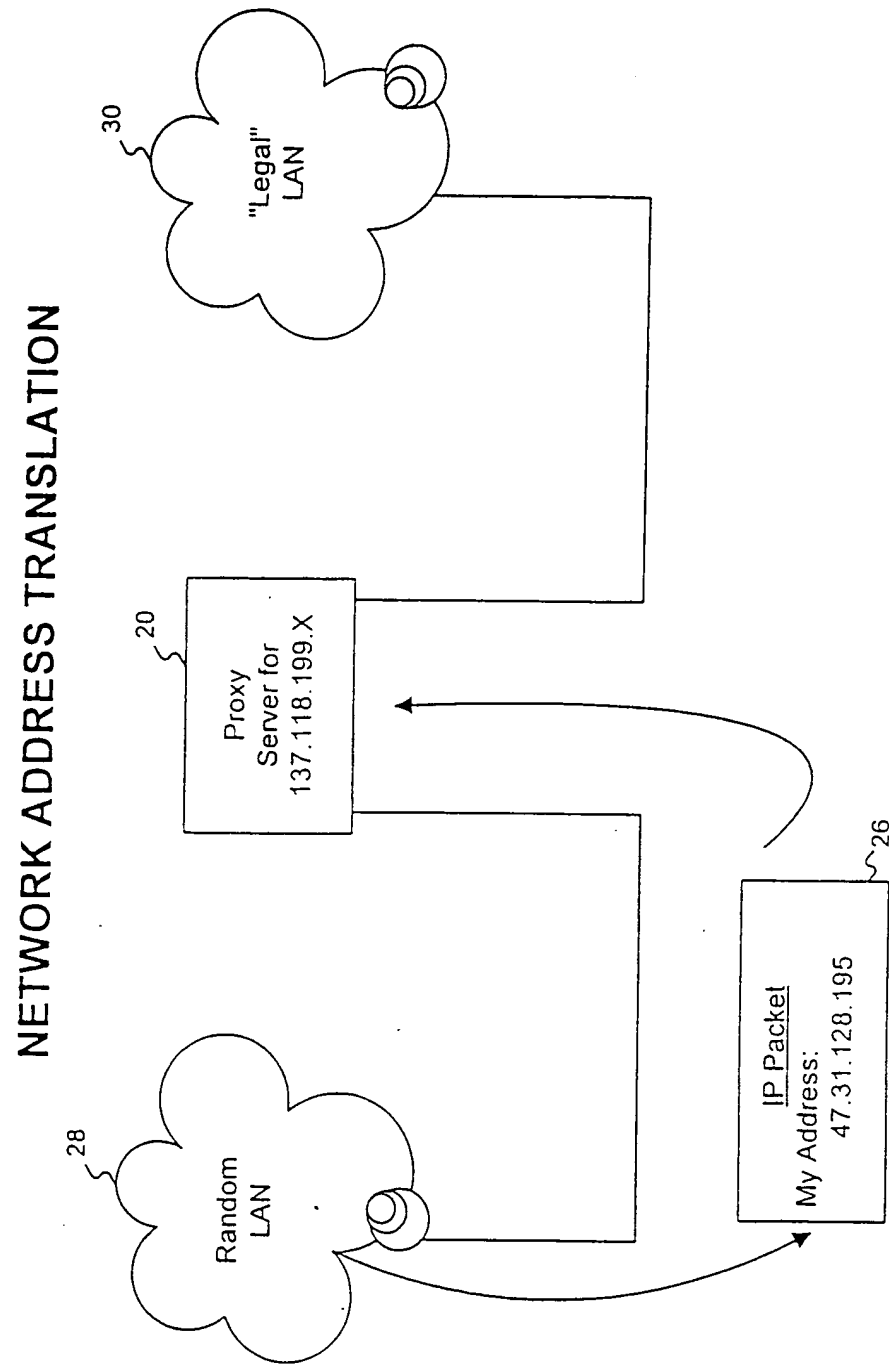


Figure 3

NETWORK ADDRESS TRANSLATION

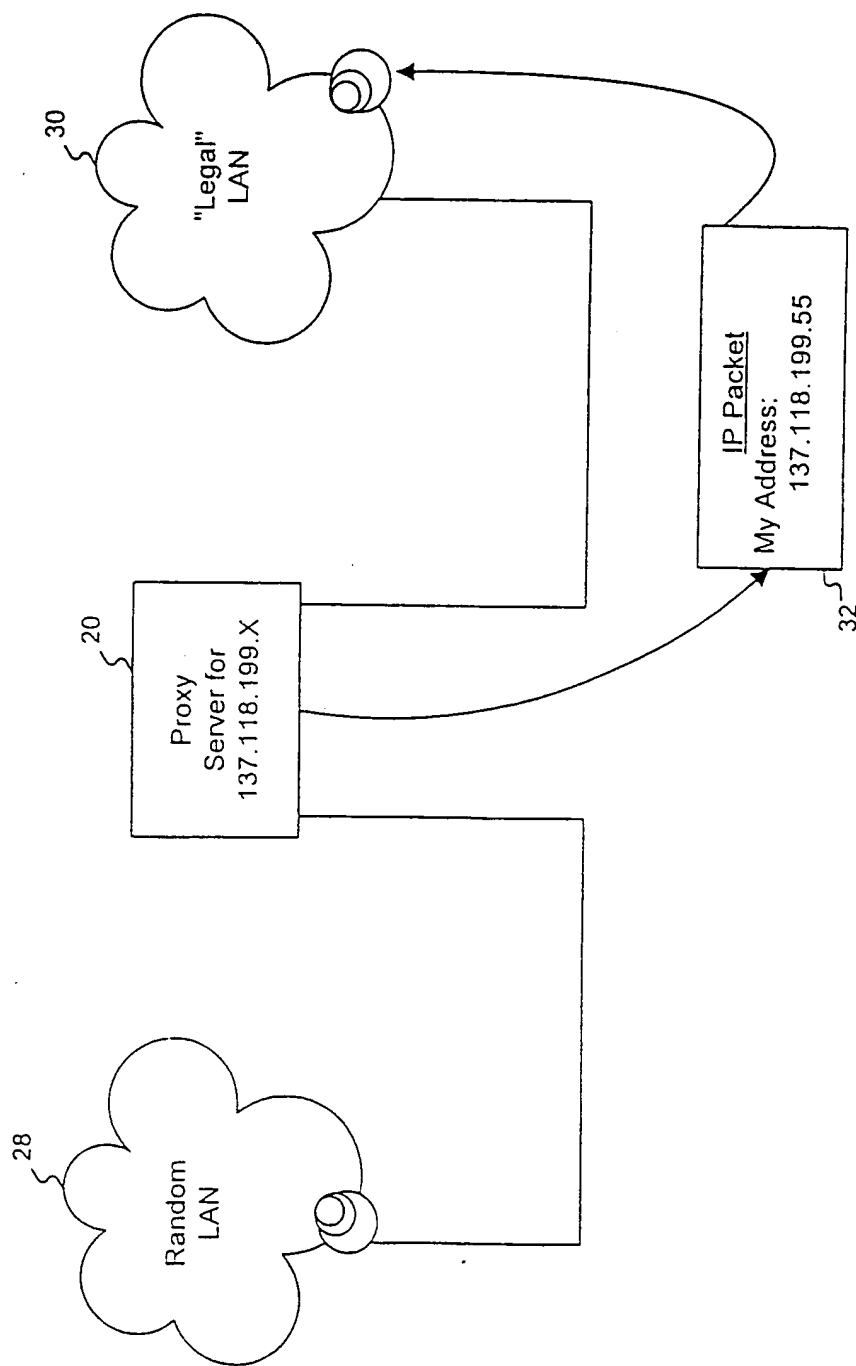


Figure 4

NETWORK ADDRESS TRANSLATION

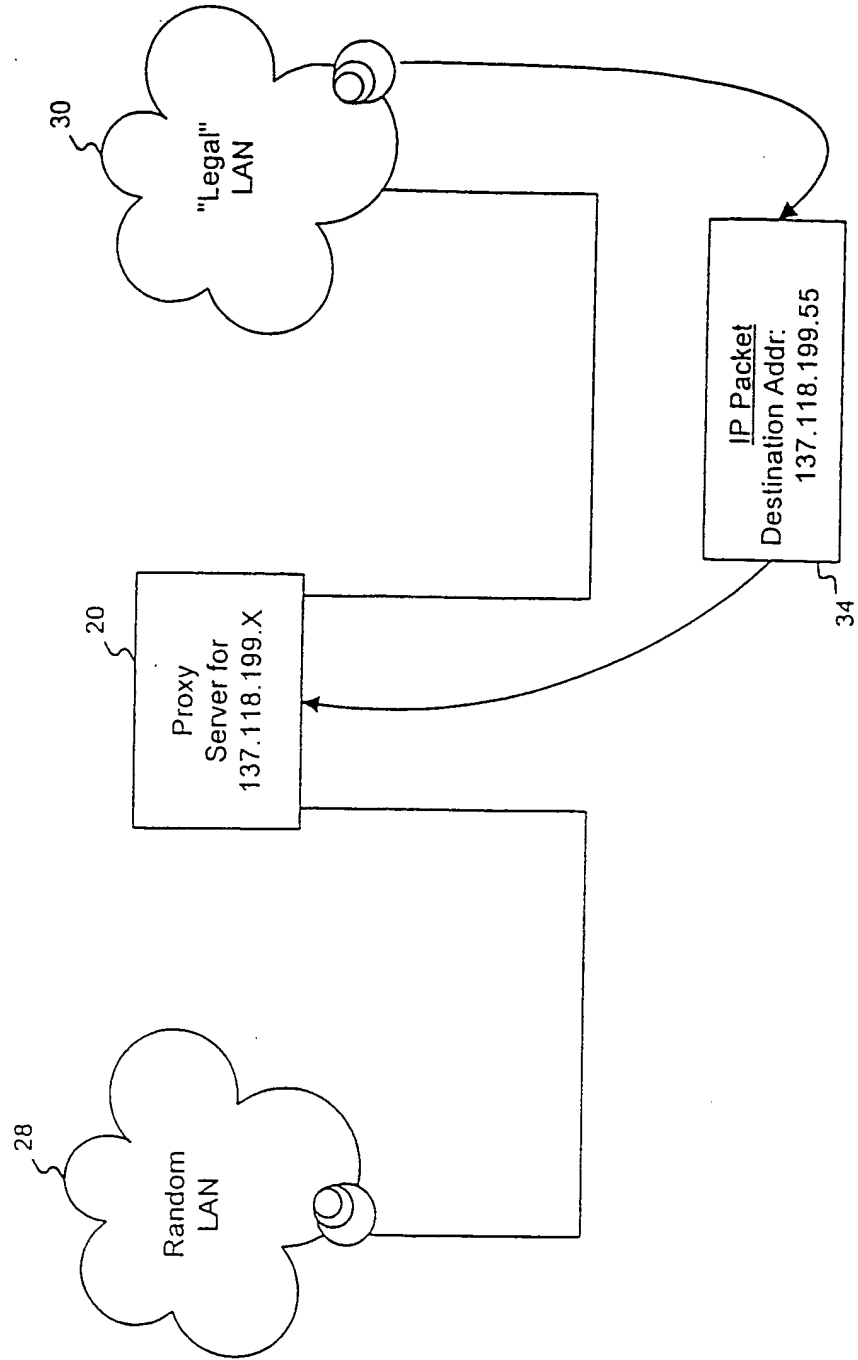


Figure 5

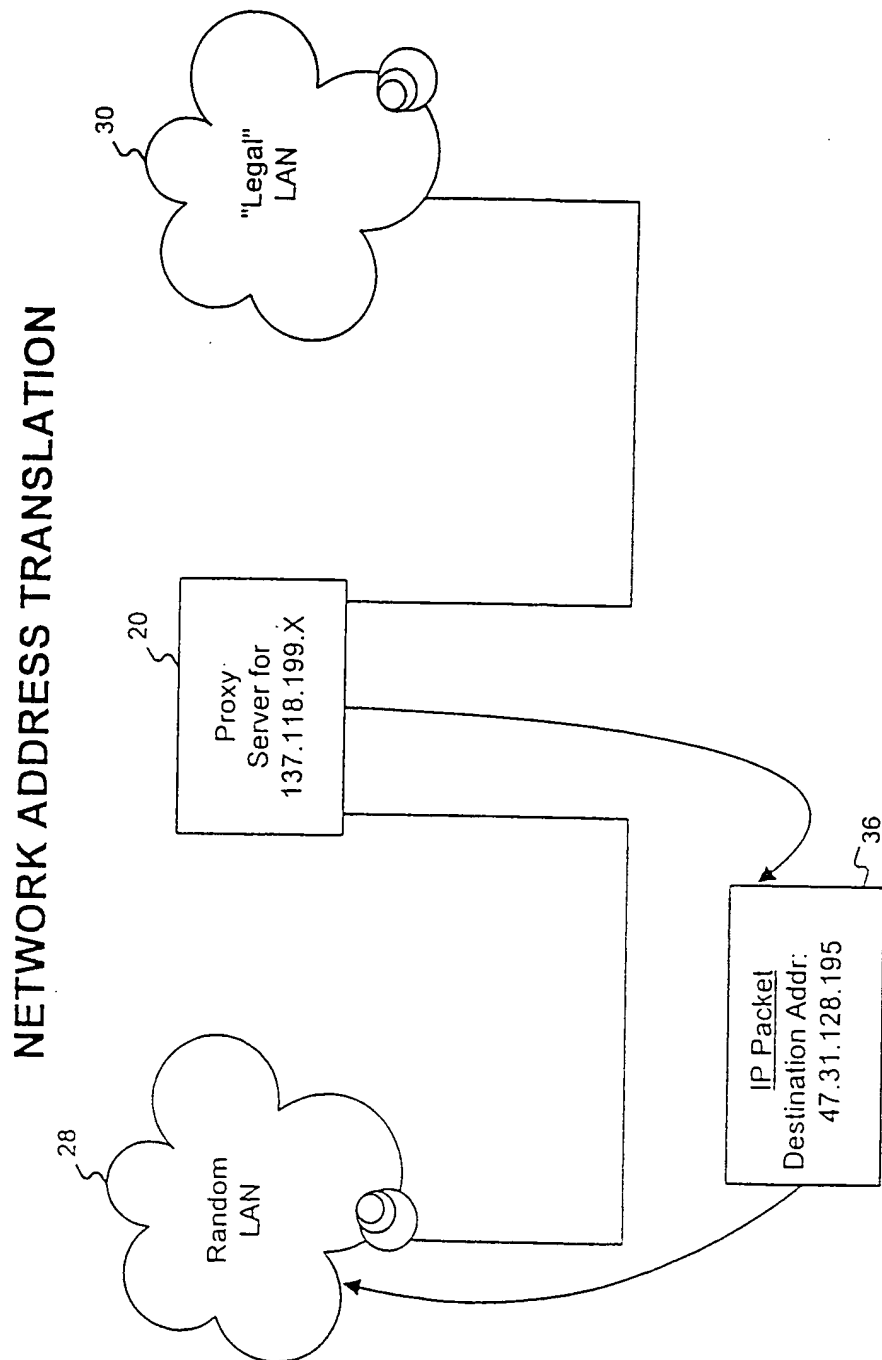


Figure 6

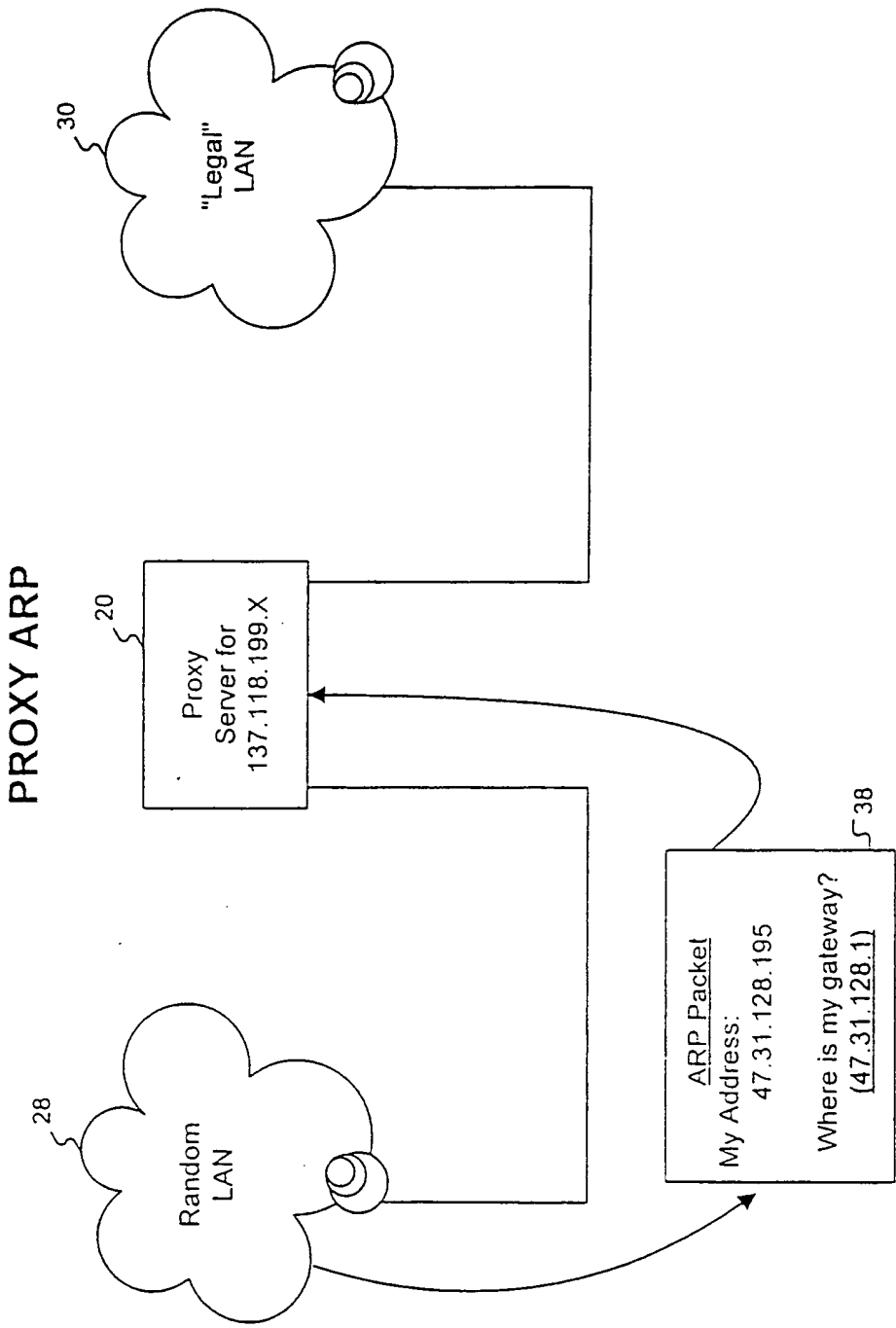


Figure 7

PROXY ARP

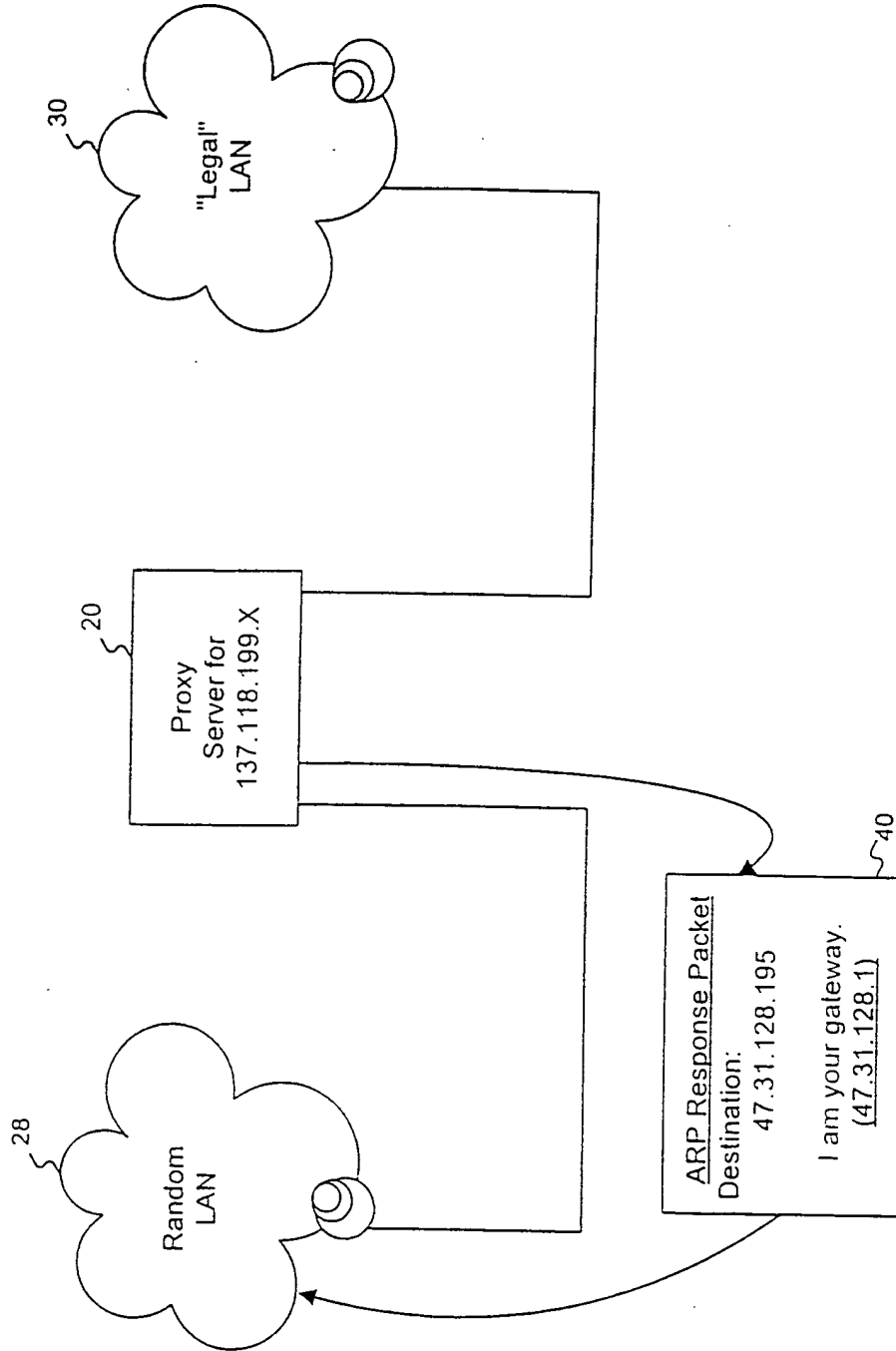


Figure 8

PROXY ARP

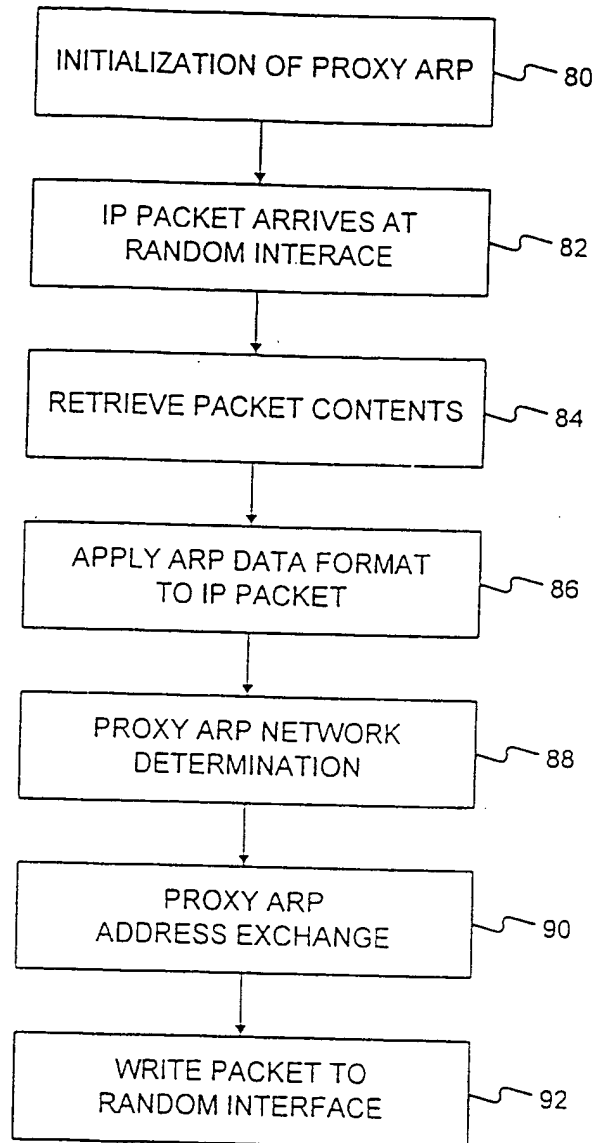


Figure 9

PROXY ARP INITIALIZATION

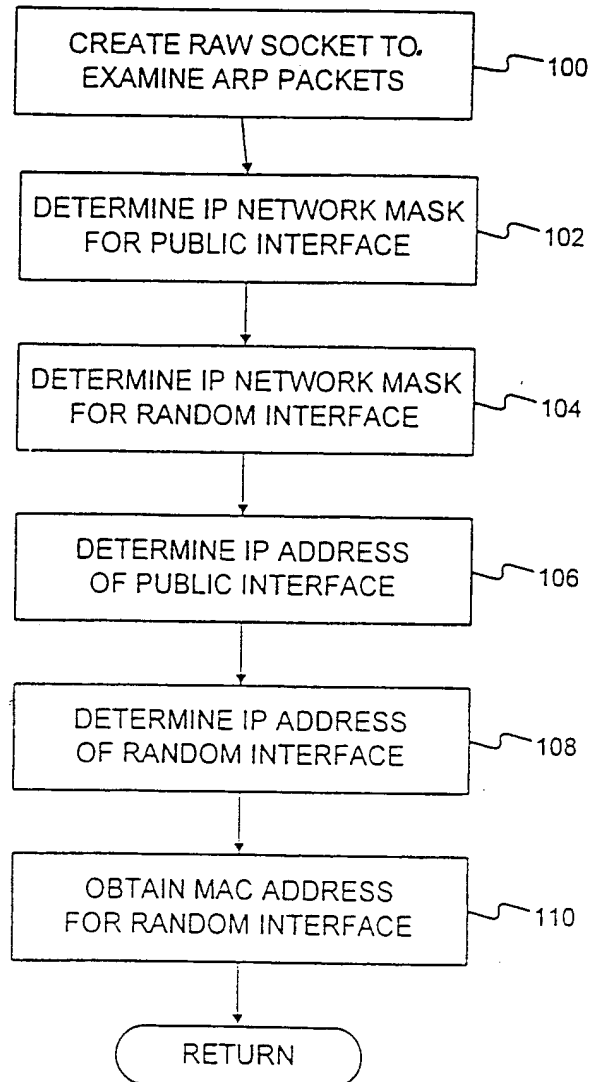


Figure 10

NETWORK DETERMINATION

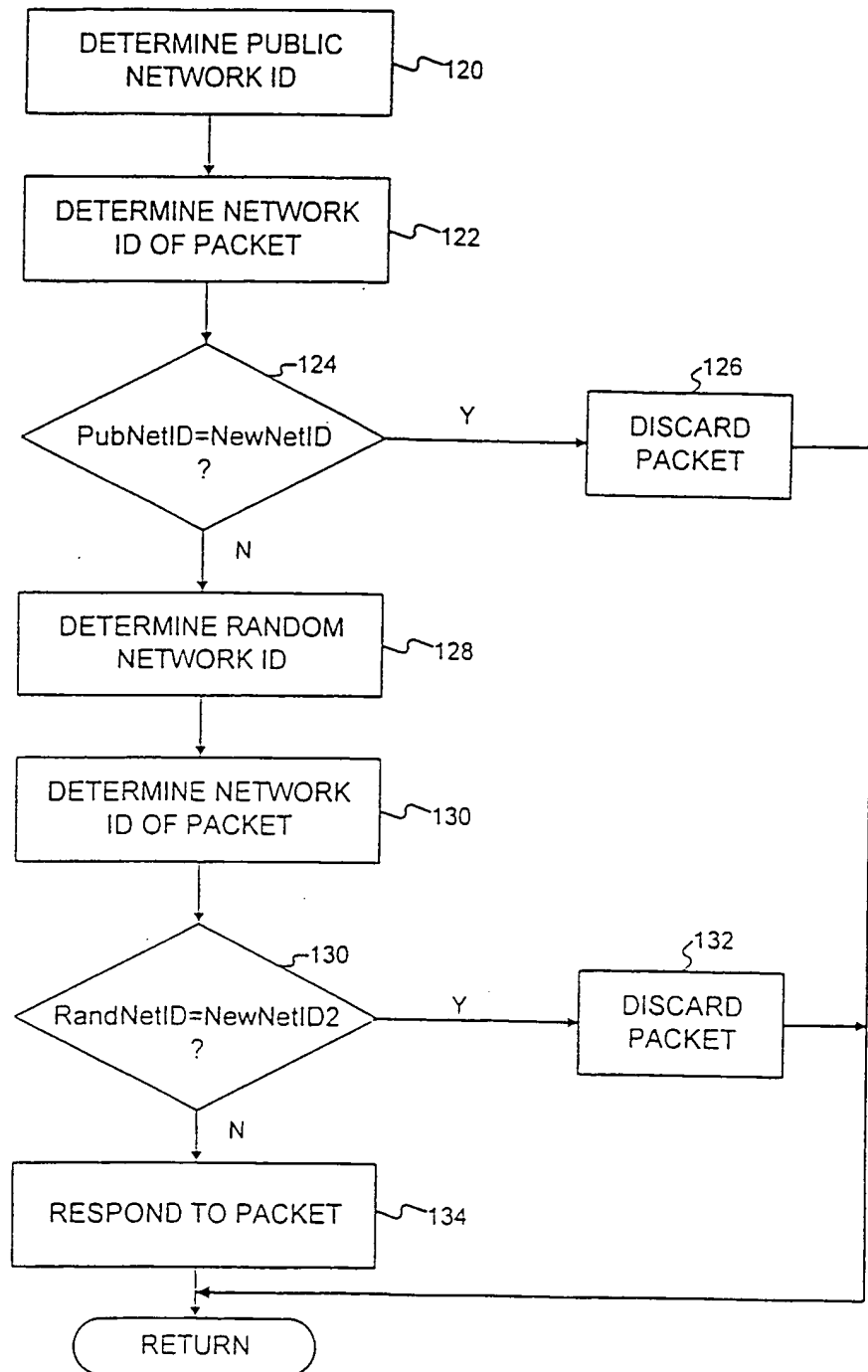


Figure 11

ADDRESS EXCHANGE

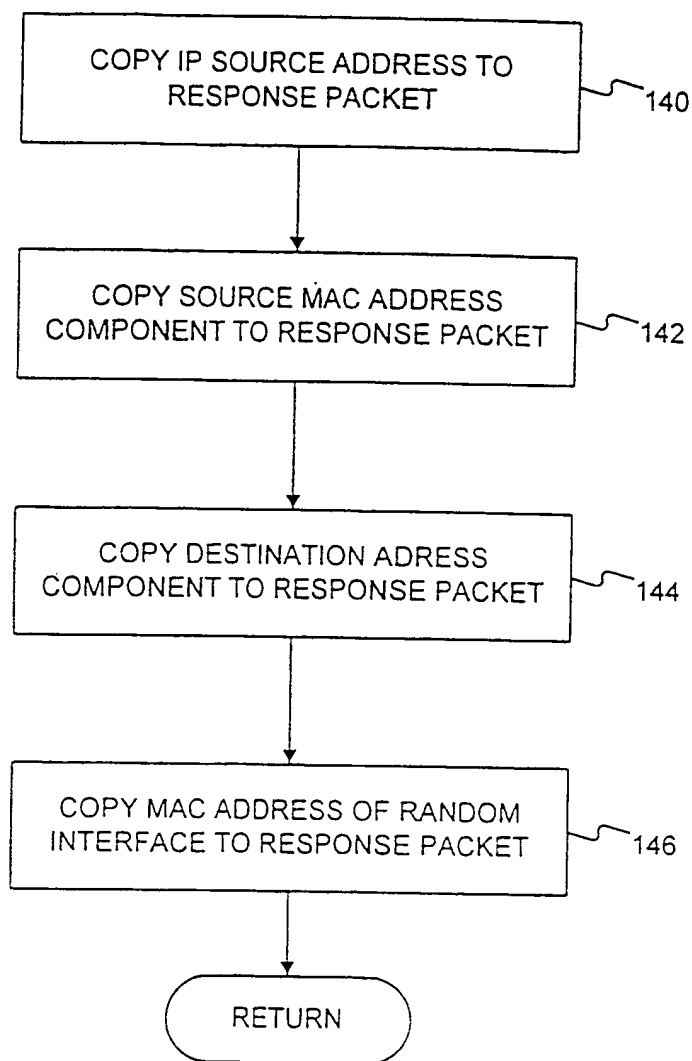


Figure 12

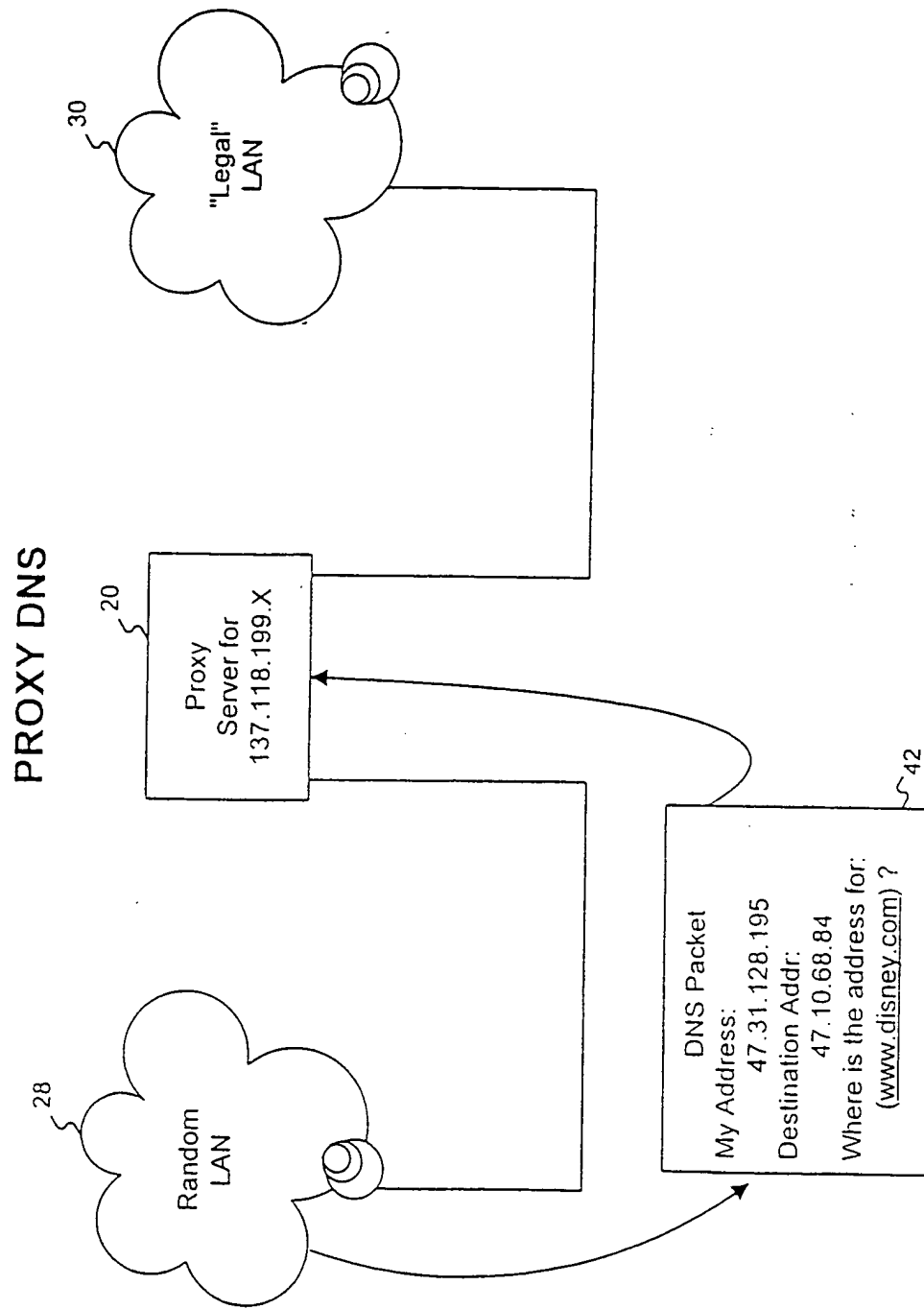


Figure 13

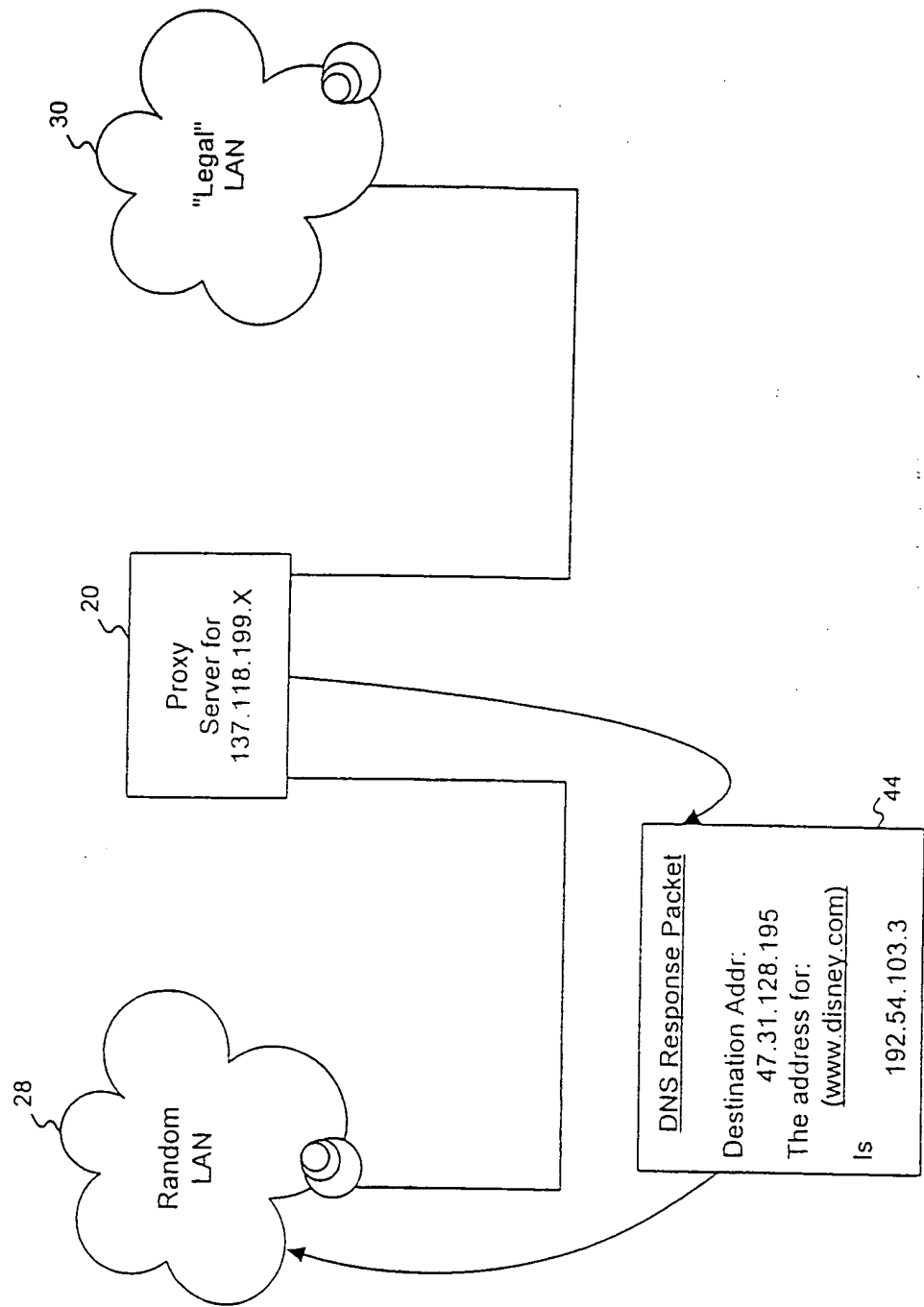


Figure 14

PROXY DNS

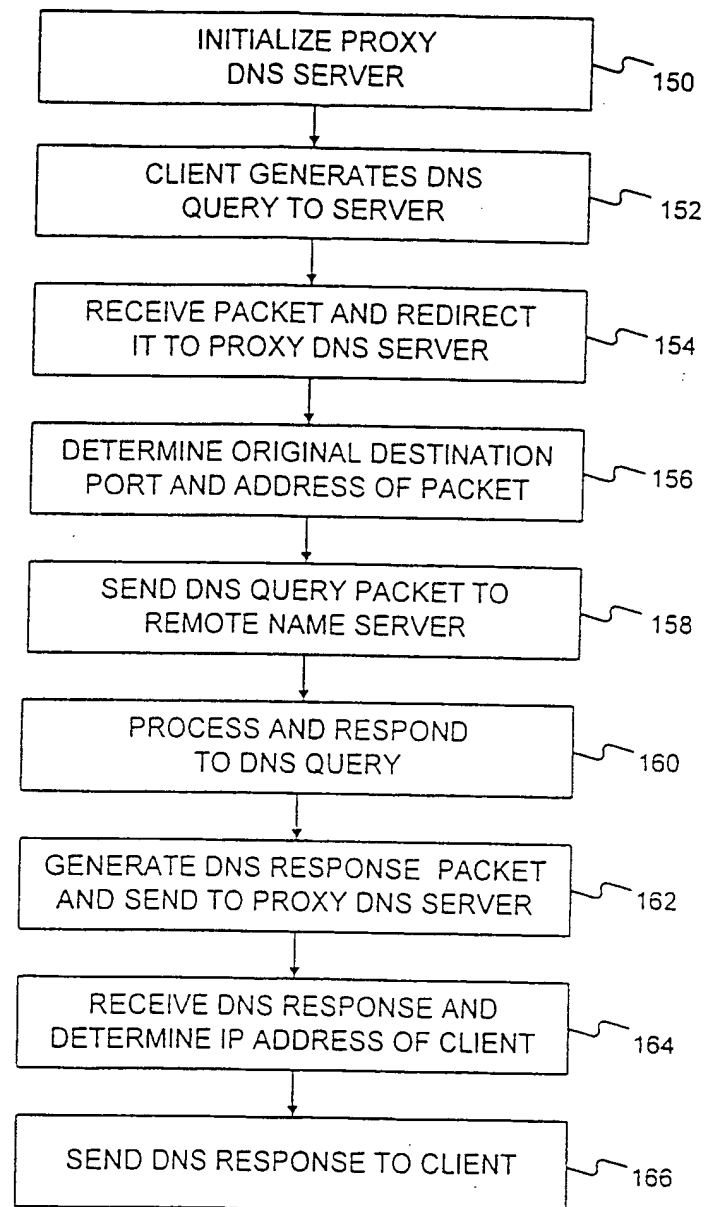


Figure 15

PROXY DNS INITIALIZATION

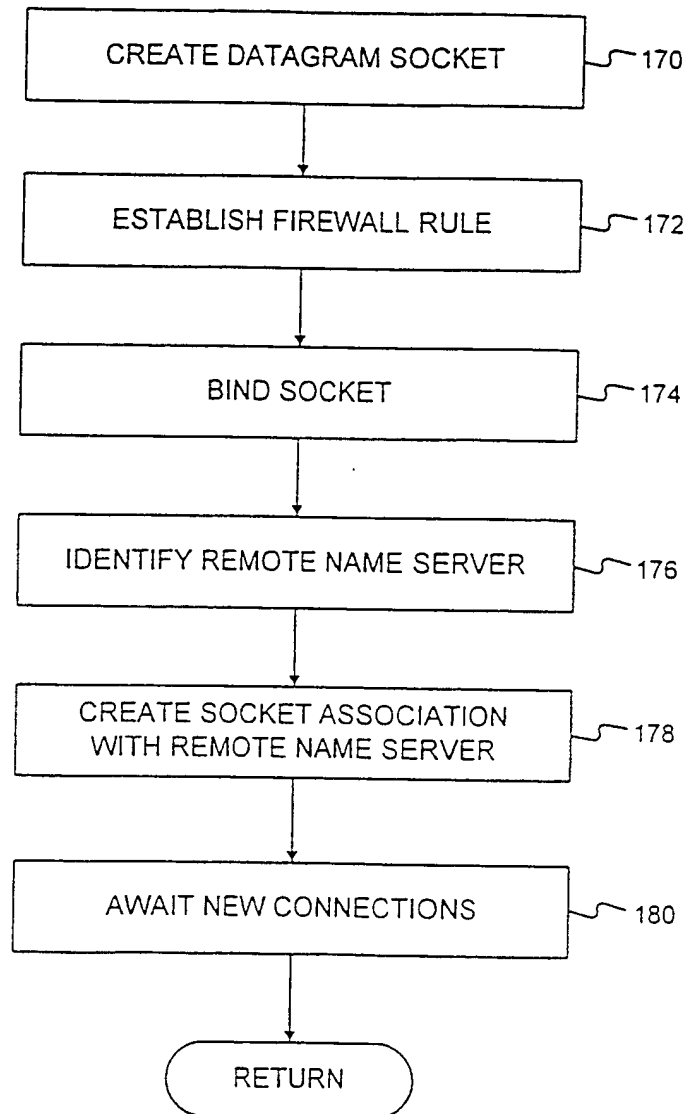


Figure 16

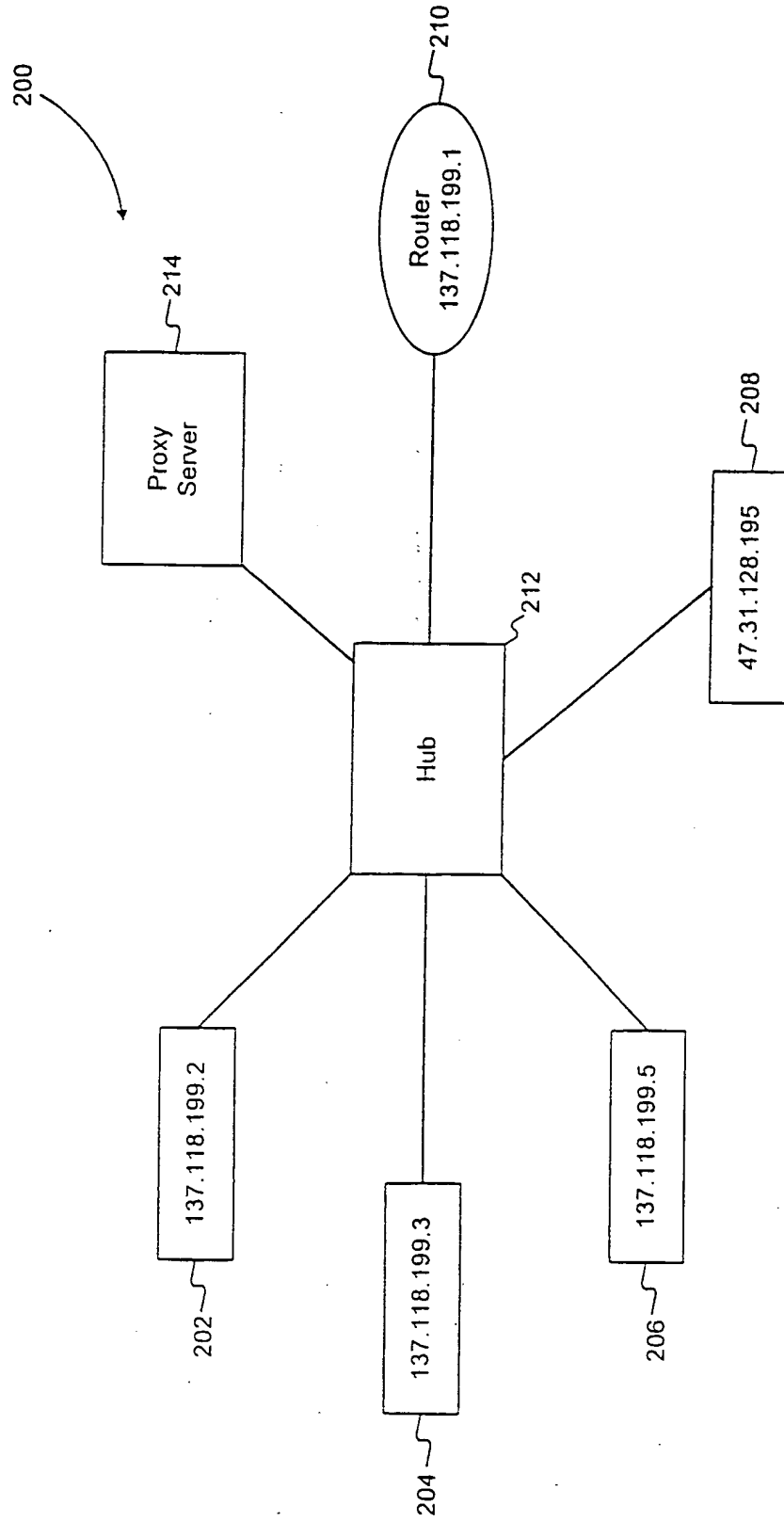


Figure 17

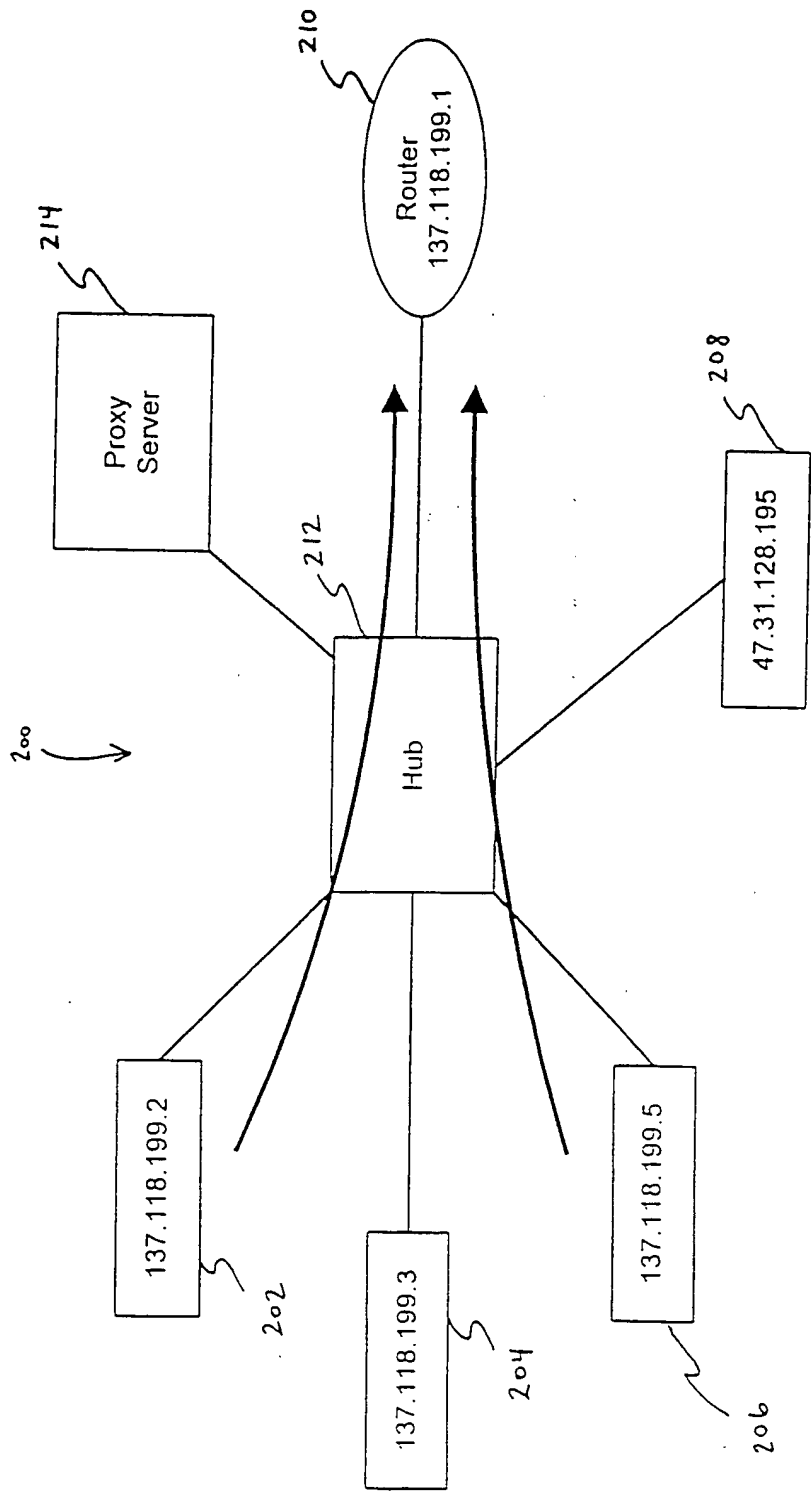


Figure 18

NETWORK ADDRESS TRANSLATION

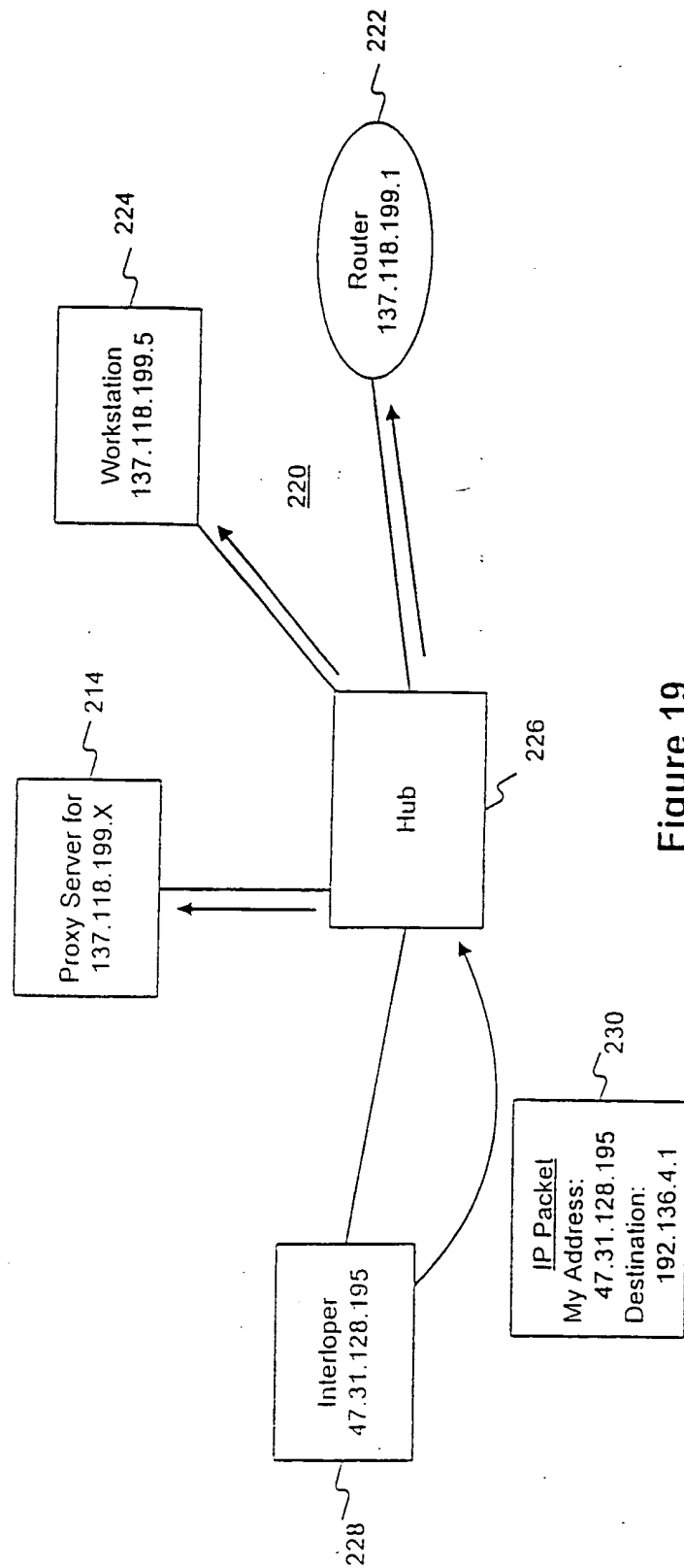


Figure 19

NETWORK ADDRESS TRANSLATION

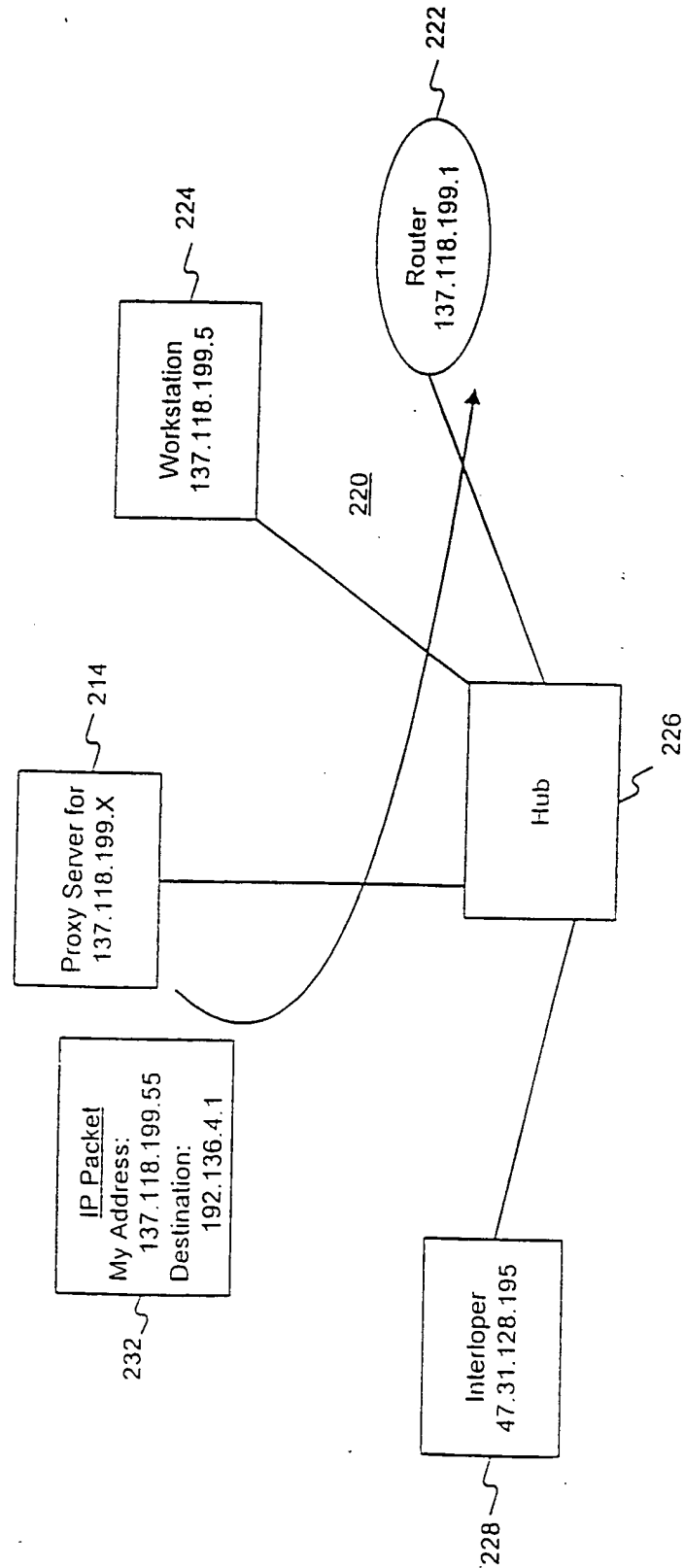


Figure 20

NETWORK ADDRESS TRANSLATION

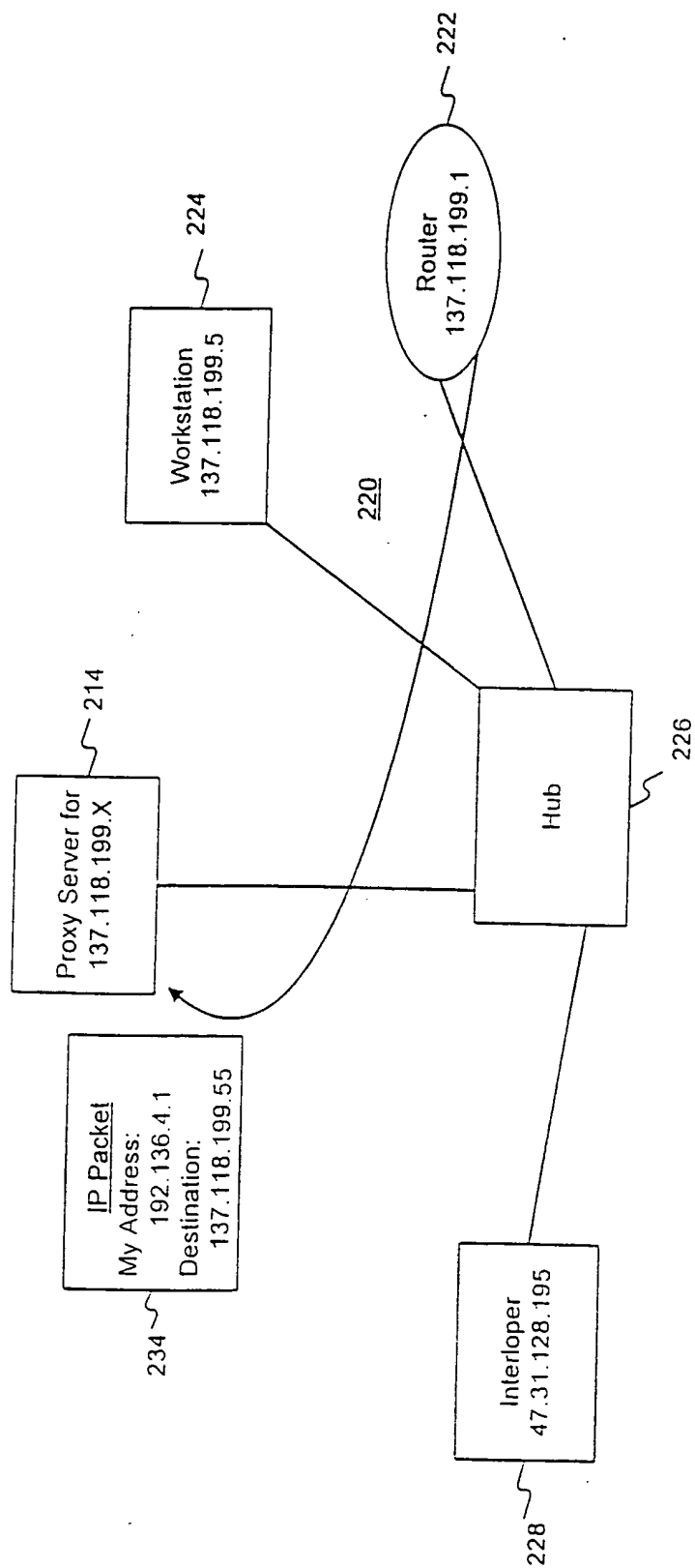


Figure 21

NETWORK ADDRESS TRANSLATION

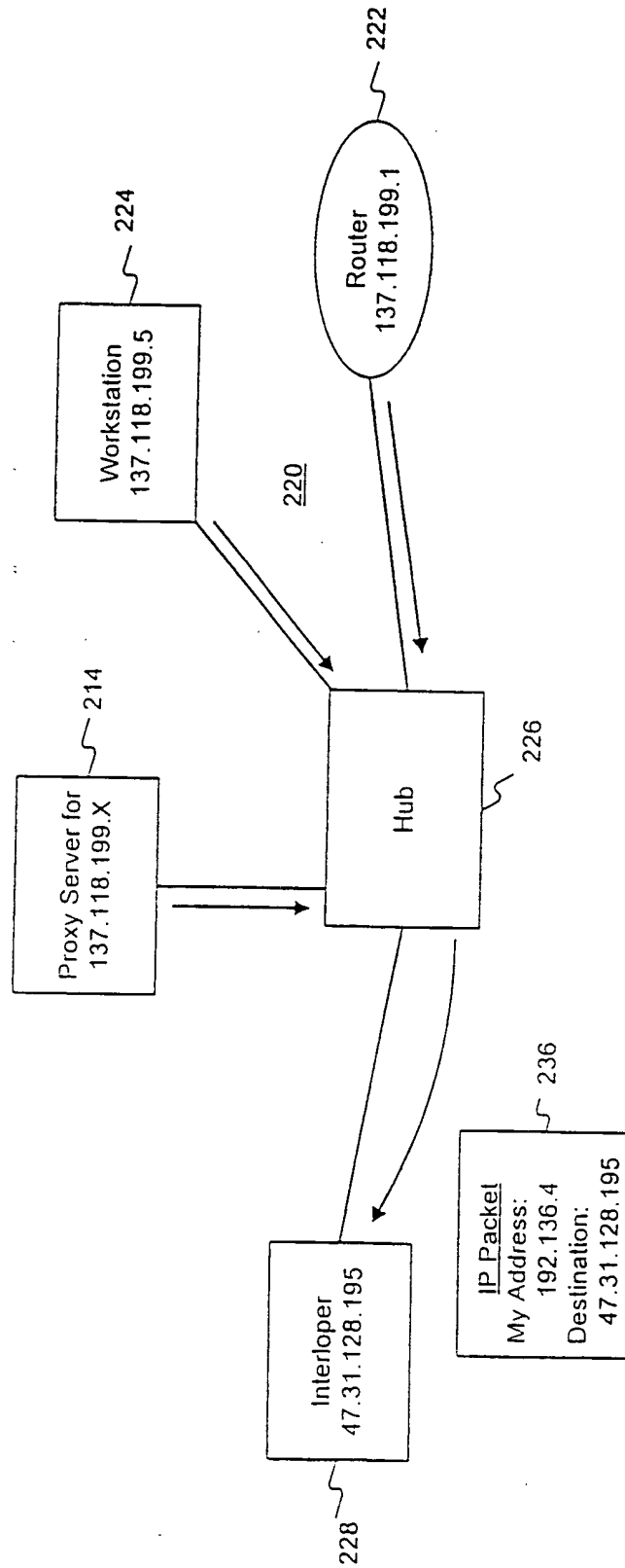


Figure 22

PROXY ARP

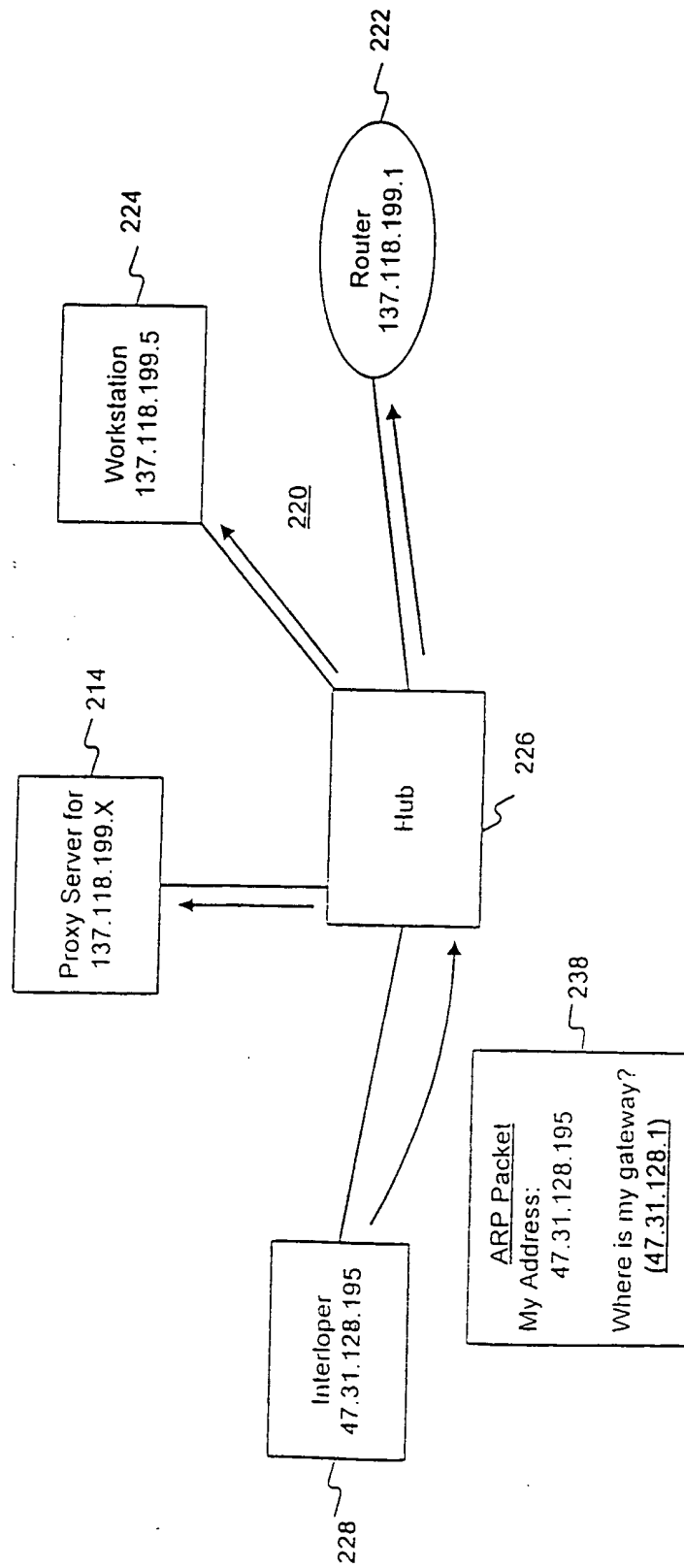


Figure 23

PROXY ARP

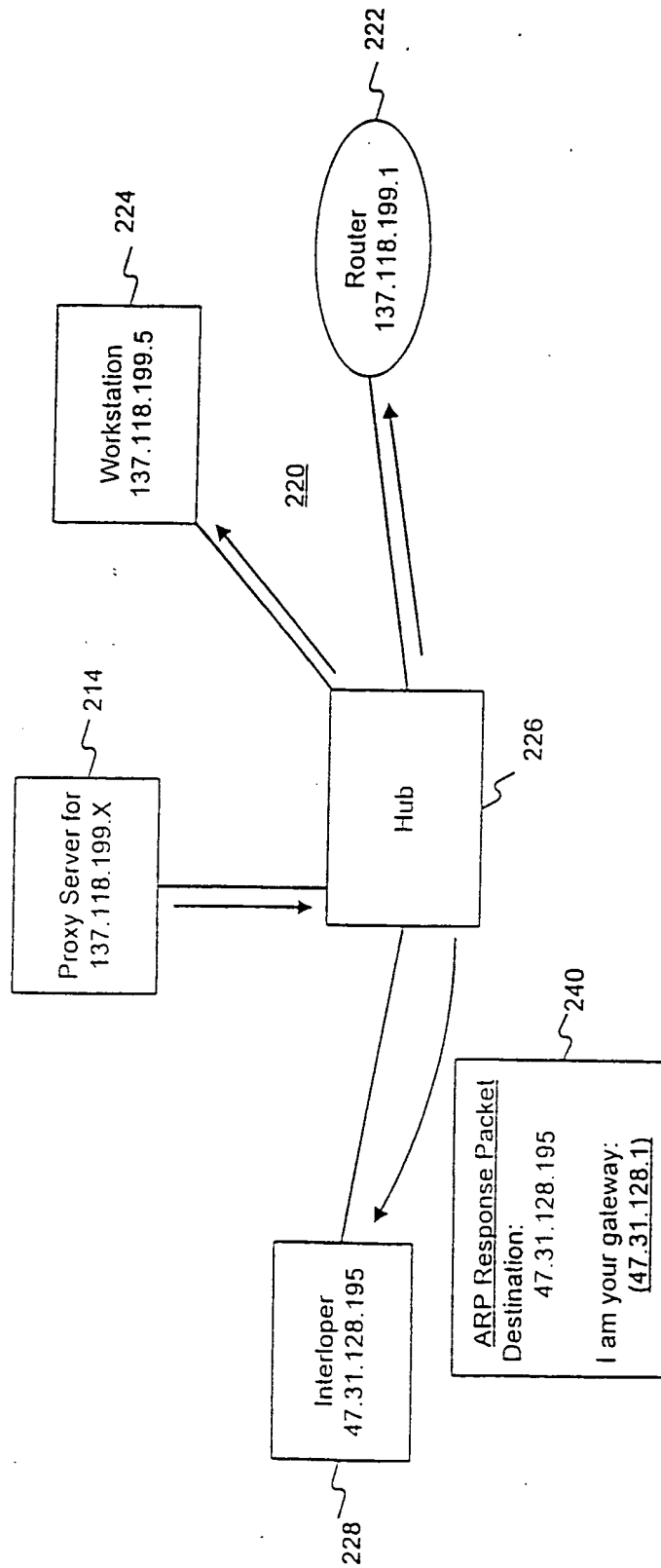


Figure 24

INTERNATIONAL SEARCH REPORT

International Application No
PCT/US 99/01195

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 H04L29/06 H04L29/12

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	GIOVANARDI A ET AL: "TRANSPARENT MOBILE IP: AN APPROACH AND IMPLEMENTATION" IEEE GLOBAL TELECOMMUNICATIONS CONFERENCE, PHOENIX, ARIZONA, NOV. 3 - 8, 1997, vol. 3, 3 November 1997, pages 1861-1865, XP000737840 INSTITUTE OF ELECTRICAL AND ELECTRONICS ENGINEERS	1-16
Y	see page 1861, left-hand column, line 1 - page 1862, right-hand column, line 25 ---	17-23
P,X	WO 98 40990 A (KLEINROCK LEONARD ;NOMADIX LLC (US); SHORT JOEL E (US)) 17 September 1998 see page 5, line 25 - page 6, line 21 see page 12, line 1 - page 13, line 26 see page 14, line 28 - page 16, line 12 --- -/--	1-23

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

* Special categories of cited documents:

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier document but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance: the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance: the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

"&" document member of the same patent family

Date of the actual completion of the international search

4 June 1999

Date of mailing of the international search report

15/06/1999

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2
NL - 2280 HV Rijswijk
Tel. (+31-70) 340-2040, Tx. 31 651 600 nl,
Fax: (+31-70) 340-3016

Authorized officer

RAMIREZ DE AREL... F

INTERNATIONAL SEARCH REPORT

Inter- national Application No
PCT/US 99/01195

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT		
Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	GROOM F M: "THE STRUCTURE AND SOFTWARE OF THE INTERNET" ANNUAL REVIEW OF COMMUNICATIONS, vol. 50, 1 January 1997, pages 695-707, XP000720937 see page 698, left-hand column, line 4 - right-hand column, line 7 ---	17-21
Y	YEOM H Y ET AL: "IP MULTIPLEXING BY TRANSPARENT PORT-ADDRESS TRANSLATOR" PROCEEDINGS OF THE SYSTEMS ADMINISTRATION CONFERENCE. LISA, 29 September 1996, pages 113-121, XP002046289 see page 114, left-hand column, line 3 - page 116, left-hand column, line 24 -----	22,23

Information on patent family members

PCT/US 99/01195

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
WO 9840990 A	17-09-1998	AU 6698498 A	29-09-1998

This Page Blank (uspto)



INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification ⁶ : H04L 29/06, 29/12	A1	(11) International Publication Number: WO 99/38303 (43) International Publication Date: 29 July 1999 (29.07.99)
<p>(21) International Application Number: PCT/US99/01195</p> <p>(22) International Filing Date: 21 January 1999 (21.01.99)</p> <p>(30) Priority Data: 60/072,175 22 January 1998 (22.01.98) US 09/057,578 9 April 1998 (09.04.98) US</p> <p>(71) Applicant (for all designated States except US): NORTHERN TELECOM LIMITED [CA/CA]; Station A, P.O. Box 6123, Montreal, Quebec H3C 3J5 (CA).</p> <p>(72) Inventors; and (75) Inventors/Applicants (for US only): BROTHERS, John, David, West [US/US]; 1620 Grand Junction, Alpharetta, GA 30004 (US). SMITH, Jeffrey, G. [US/US]; 12250 Stevens Creek Drive, Alpharetta, GA 30005 (US).</p> <p>(74) Agents: GARRETT, Arthur, S. et al.; Finnegan, Henderson, Farabow, Garrett & Dunner, L.L.P., 1300 I Street, Washington, DC 20005-3315 (US).</p>	<p>(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GD, GE, GH, GM, HR, HU, ID, IL, IN, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, US, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).</p> <p>Published With international search report. Before the expiration of the time limit for amending the claims and to be republished in the event of the receipt of amendments.</p>	
<p>(54) Title: PROXY SERVER FOR TCP/IP NETWORK ADDRESS PORTABILITY</p> <div data-bbox="389 1155 1380 1533"> <pre> graph LR 12[192.5.144.8] --- 18[Hub] 14[10.3.28.5] --- 18 16[47.31.128.195] --- 18 18 --- 20[Proxy Server] 20 --- 22([Router]) </pre> </div> <p>(57) Abstract</p> <p>System and method for fully transparent IP mobility services for clients in a dynamic LAN Ethernet environment. The functionality within a proxy server, a combination of network address translation, proxy address resolution protocol, and proxy domain name service, allow the proxy server to support and provide full IP client functionality to any IP-enabled network device in any proxy server enabled LAN. The proxy server may be added to an existing Ethernet (or Ethernet-emulated) LAN. Once configured with the necessary subnet range, DNS, and IP pools, the proxy server provides support for any mobile device that enters the LAN.</p>		

FOR THE PURPOSES OF INFORMATION ONLY

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece			TR	Turkey
BG	Bulgaria	HU	Hungary	ML	Mali	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MN	Mongolia	UA	Ukraine
BR	Brazil	IL	Israel	MR	Mauritania	UG	Uganda
BY	Belarus	IS	Iceland	MW	Malawi	US	United States of America
CA	Canada	IT	Italy	MX	Mexico	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NE	Niger	VN	Viet Nam
CG	Congo	KE	Kenya	NL	Netherlands	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NO	Norway	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	NZ	New Zealand		
CM	Cameroon			PL	Poland		
CN	China	KR	Republic of Korea	PT	Portugal		
CU	Cuba	KZ	Kazakstan	RO	Romania		
CZ	Czech Republic	LC	Saint Lucia	RU	Russian Federation		
DE	Germany	LI	Liechtenstein	SD	Sudan		
DK	Denmark	LK	Sri Lanka	SE	Sweden		
EE	Estonia	LR	Liberia	SG	Singapore		

PROXY SERVER FOR TCP/IP NETWORK ADDRESS PORTABILITY

Background of the Invention

The present invention relates generally to address portability and, more particularly,
5 to a method and apparatus for address portability to provide fully transparent internet
protocol (IP) mobility services to IP-enabled network devices in any Ethernet local area
network (LAN).

The transmission control protocol/internet protocol (TCP/IP protocol), a suite of
communications protocols used by host computers to exchange information between
10 application processes over LANs or wide area networks (WANs), was designed when
laptops and other mobile IP devices were essentially nonexistent. As a result, there was
no issue with mobility, since each IP network device was typically a workstation,
minicomputer, or the like. The movement of devices from place to place in such a static
environment was expected to be a very rare occurrence, and one that could be adequately
15 handled by manual intervention. This assumption, in conjunction with various resource
constraints, influenced the development of the IP protocol such that each LAN only
operated with a limited range (a subnet) of IP addresses. Any device with an IP address
outside of that range was simply ignored by the LAN's router, rendering it unable to
communicate with any device within that network.

20 Over the last several years, the IP protocol has become the primary data
communications protocol on virtually every computer in the world. This includes a
substantial number of laptops and other portable computer devices. As the prevalence of
laptops increases, IP mobility issues have substantially increased. For example, it is now
common for customers, vendors, and even business associates that have laptops or other
25 mobile IP devices to attempt to hook into a "foreign" LAN and attempt to use its facilities.
Typically, this results in significant frustration since the amount and complexity of
reconfiguration to permit the connection is not insubstantial.

One attempted solution to this problem, dynamic host configuration protocol
(DHCP), evolved over the last couple of years. Under DHCP, a computer configured to
30 use that protocol may retrieve local IP configuration data automatically when the mobile
IP device is connected to the network. While this is a reasonable solution to mobility
problems, its scope is somewhat limited. For example, the mobile network device must

- 2 -

be configured to use DHCP, and the LAN must have a DHCP server enabled. Moreover, the duration of DHCP "timeouts" within the mobile network device must be short enough to allow the device to request a new address at the new location. As a result of at least these limitations, DHCP has not sufficiently solved the problem. In some cases, DHCP
5 has proven unacceptable to the network clients who may not have DHCP pre-configured or to network administrators who wish to have more knowledge of and control over the mobile IP devices that enter and leave the network.

It is, therefore, desirable to provide fully transparent IP mobility services for clients in a dynamic network environment.

10 Summary of the Invention

Systems and methods consistent with the present invention satisfy this and other needs by supporting and providing full IP client functionality to any IP-enabled network device in any mobility-enabled LAN. The present invention provides full functionality regardless of both the IP address of the mobile device and subnet restrictions of the LAN.

15 A method for use with a proxy server consistent with the present invention establishes communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network. The method includes the step of generating an address resolution protocol packet to identify the arbitrary address for the destination device. The proxy server receives the address
20 resolution protocol packet and generates an address resolution protocol response packet including the arbitrary address of the destination device. The method also includes the step of transmitting the address resolution protocol response packet from the proxy server to the device in the first network.

Another method consistent with the present invention is for use with a proxy server
25 and establishes communications between a random device and a destination device having an arbitrary address on a network. The method includes the steps of generating an address resolution protocol packet to identify the arbitrary address for the destination device and receiving, by the proxy server, the address resolution protocol packet. The method also includes the steps of generating an address resolution protocol response packet including
30 the arbitrary address of the destination device and transmitting the address resolution protocol response packet from the proxy server to the random device.

- 3 -

Yet another method consistent with the present invention is for use with a proxy server which is in communication with a mobile device and remote name server. The method permits obtaining an internet protocol address from the remote name server for the mobile device and includes the steps of generating a query packet including a request for
5 an address associated with a domain name and receiving the query packet from the mobile device in the proxy server. The method also includes the steps of forwarding the query packet to the remote name server and generating a response packet including the requested address. The method also includes transmitting the response packet to the proxy server and transmitting the response packet to the mobile device.

10 Another method consistent with the present invention is for use with a proxy server and provides for communications between a random device and a destination device in a network. The method includes the steps of performing a proxy address resolution protocol to initiate communications between the random device and the destination device, performing a proxy domain name service to identify a destination address in the second
15 network associated with a domain name, and performing a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the destination device in the network. Use of this combination allows a system to support and provide full client functionality to mobile network devices.

20 Systems are also provided for carrying out these and other methods consistent with the present invention.

Several advantages accrue to methods and systems consistent with the present invention. For example, these systems and methods provide a secure and complete mobility solution, including the various cases where prior art solutions were inadequate.

25 Such systems and methods are completely transparent to the end-user, who may or many not use DHCP, but will still be able to communicate with a LAN or even with a WAN. They are also more "administrator-friendly", especially when the acceptance protocol involves e-mail notification to the network administrator that a new device has joined the network. Security is enhanced by reducing the network's exposure to foreign snooping.

30 The above and additional features and advantages of the present invention will be readily appreciated by one of ordinary skill in the art from the following detailed description.

Brief Description of the Drawings

Figure 1 is a diagram of a random Ethernet LAN environment consistent with the present invention;

Figure 2 is a high level system diagram of a proxy server consistent with the present
5 invention;

Figures 3 and 4 illustrate network address translation associated with the routing of traffic from a random LAN to a legal LAN consistent with the present invention;

Figures 5 and 6 illustrate network address translation associated with the routing of traffic from a legal LAN to a random LAN consistent with the present invention;

10 Figures 7 and 8 illustrate generation of a proxy address resolution protocol (ARP) packet and generation of a proxy ARP response packet consistent with the present invention;

Figures 9-12 are flowcharts depicting steps for proxy ARP consistent with the present invention;

15 Figures 13 and 14 illustrate generation of a proxy domain name service (DNS) query packet and generation of a proxy DNS response packet consistent with the present invention;

Figures 15 and 16 are flowcharts depicting steps for proxy DNS consistent with the present invention;

Figure 17 illustrates an alternative proxy server implementation consistent with the present invention;

Figure 18 illustrates normal traffic flow in the alternative proxy server implementation of Figure 16;

Figures 19-22 illustrate network address translation for use with the alternative
25 proxy server implementation consistent with the present invention; and

Figures 23-24 illustrate proxy ARP for use with the alternative proxy server implementation consistent with the present invention.

Detailed Description of the Preferred Embodiments

Reference will now be made in detail to embodiments consistent with this invention
30 that are illustrated in the accompanying drawings. The same reference numbers in
different drawings generally refer to the same or like parts.

Two-Armed Proxy Server

- 5 -

Figure 1 shows a "random" LAN 10. For purposes of this discussion, "random" simply means there are a random number of mobile IP network devices, and those devices each use a random IP address. One example of a "random" LAN would be an IP Ethernet network in a hotel. LAN 10 includes a plurality of interconnected mobile IP network devices, including laptops 12, 14, and 16, having IP addresses scattered across the range of known addresses. As shown, the plurality of network devices communicate through a hub 18 and a proxy server 20 with network router 22. The links used to interconnect the various network elements shown may, for example, be Ethernet links. In the LAN 10, proxy server 20 may be referred to as a "two-armed" (TA) proxy server since it possesses two network interfaces, *e.g.*, two Ethernet links.

Normally, this type of network would be extremely difficult to manage, since a standard router expects all IP addresses that it serves to fall within a limited range. Consistent with the present invention, however, traffic from each of these devices may be modified so that the information presented to the network router is acceptable. The modification may be accomplished by proxy server 20 using a combination of network address translation (NAT), a proxy address resolution protocol (ARP) service, and a proxy domain name system (DNS) service, as discussed below. NAT is a well-known process by which traffic received by and transmitted from a particular device with an arbitrary IP address is modified to present the correct IP address to a network router. NAT service may be specifically configured to translate particular IP addresses. A proxy mobility server consistent with the present invention may translate random IP addresses dynamically.

Figure 2 illustrates a high level diagram of proxy server 20. As shown, proxy server 20 includes a processor in communication with a hard drive, a system memory, and a user memory. The system and user memories may include read-only and/or random access types of memories. These memories are useful for storing packet contents, which may include addresses and the like, as well as data content and packet length, to name a few. Proxy server 20 also includes interfaces, which may take the form of cards, through which proxy server communicates with networks. In the example shown, proxy server 20 is interfaced with a legal/public network and a random network through Ethernet interface 0 and Ethernet interface 1, respectively.

Figures 3 and 4 depict the routing of information, such as an IP packet 26, from a device within a random LAN 28 to a device within a "legal" LAN 30. For purposes of this

- 6 -

discussion, a "legal" LAN is simply a public LAN, *i.e.*, one with a legal set of IP addresses. In Figure 3, packet 26 from the device, which has an IP address of 47.31.128.195, is routed from random LAN 28 to proxy server 20. As shown, proxy server 20 is a server for a subnet (here denoted 137.118.199.X) which, as is known, is a set of machines that are
5 physically connected together in an Ethernet LAN, *e.g.*, the legal LAN 30 in Figure 3. Proxy server 20 performs a network address translation and, as shown in Figure 4, the translated packet 32 is transmitted to the legal LAN. The translation may be performed using known techniques, such as those specified in Internet Engineering Task Force Request for Comments (IETF RFC) 1631. A packet following the opposite path, *i.e.*,
10 packet 34 routed from a device within legal LAN 30 to a device within random LAN 28, would also undergo network address translation (to translated packet 36) in proxy server 20 (see Figures 5 and 6).

Proxy ARP

In addition to NAT, proxy server 20 may employ a proxy address resolution
15 protocol (ARP) service to provide mobile functionality consistent with the present invention. ARP is a known protocol which may be used by a network device to discover what other devices are connected to the local network. Proxy ARP service allows TA proxy server 20 to "spoof" mobile network IP devices having random IP addresses into thinking that server 20 is the device with which those mobile IP devices wish to
20 communicate. This is necessary when the mobile IP device first boots and attempts to determine its gateway. As is known, existing proxy ARP implementations are limited in their use since only traffic from certain select specific addresses can be handled. Proxy ARP consistent with the present invention is not so limited and may be used to identify any arbitrary address.

25 Figure 7 illustrates an ARP packet 38 being transmitted by a mobile device in random LAN 28 to proxy server 20. As shown, ARP packet 38 includes the address of the sending device, as well as a query from the device regarding the whereabouts of its gateway, which has the address of the gateway sought. The gateway, which may be a network router, connects the smaller LAN (*e.g.*, a random LAN) to a larger WAN (*e.g.*,
30 a public LAN, such as the "legal" LAN 30 of Figure 7) and passes traffic from the LAN to the WAN. When a mobile device in the random LAN wishes to send traffic to a device having an arbitrary address in a second network (*e.g.*, the WAN) outside of the LAN, the

- 7 -

sending device needs to know to which gateway device the traffic should be sent. While normally the gateway is on the same network as the mobile device that is searching for it (using the ARP), this is not possible in a random LAN. Accordingly, the proxy server pretends that it is the gateway, and the mobile device will use it to reach the WAN. In
5 response to receiving ARP packet 38, proxy server 20 generates an ARP response packet 40 destined for the sending device in random LAN 28. As shown in Figure 8, response packet 40 includes the sending device's IP source address as the destination address and informs the sending device that proxy server 20 is the device's gateway.

Figure 9 depicts steps for proxy ARP consistent with the present invention. During
10 initialization (step 80), IP addresses and network masks are determined, as shown in greater detail in Figure 10. First, a raw socket is created to examine ARP packets (step 100). This socket is a communications programming interface created between proxy server 20 and the random and public LANs. In the Redhat Linux operating system, only one such socket is needed, whereas other operating systems, such as Sun Microsystems's
15 Solaris operating system, must create one socket per Ethernet interface. Next, IP network masks for both the public LAN interface and the random LAN interface are identified (steps 102, 104). As is known, these interfaces may be cards in proxy server 20. This information, used by the proxy server in combination with other information to determine what set of devices are part of the networks and, therefore, whether it can send packets
20 directly to any mobile device or whether the packets must be sent through a router, is typically maintained in a stable storage device, such as hard disk, flash memory, and the like, in proxy server 20. Similarly, IP addresses of the public and random LAN interfaces are identified (steps 106, 108), as is the medium access control (MAC) address of the random interface (step 110). Like the IP network masks, these addresses are typically
25 stored or built into proxy server 20.

With continuing reference to Figure 9, when a new ARP packet from a random IP device, such as ARP packet 38 of Figure 7, arrives at the random network interface from a mobile IP device in the random LAN (step 82), proxy server 20 retrieves the packet contents from the operating system (step 84). The packet, as is known, has a header and
30 data, which includes *inter alia* the packet source address (*i.e.*, the address of the mobile device that sent the ARP packet) and the packet destination address (*i.e.*, the address of that device's gateway). The server then applies the ARP data format to the IP packet (step 86).

- 8 -

The ARP data format is defined in IETF Standard 37, and the application of the format to the data may be done using the standard method of casting.

Next, the proxy server performs a proxy ARP network determination to determine the network to which the IP packet is destined for (step 88). Figure 11 shows a flowchart
5 detailing steps for this determination consistent with the present invention. A public network ID (PubNetID) is determined first (step 120). In one embodiment, the public network ID is derived from the public IP interface address and the IP network mask of the public LAN, *e.g.*, logically "anding" the address with the mask. The proxy server uses the PubNetID to discover what network it is part of, *i.e.*, what IP devices are local and which
10 are not local. Devices that not local are reached through a router. Next, the proxy server determines if the IP destination address of the incoming packet is for a "local" network, *i.e.*, the public LAN. In one embodiment, a network ID associated with the incoming packet (NewNetID) is determined (step 122) based on the destination IP address sub-component of the ARP data structure associated with the ARP packet and the IP network
15 mask of the public LAN interface, *e.g.*, logically "anding" the destination address with the public IP network mask. If PubNetID is equal to NewNetID (step 124), the incoming packet is destined for a device that has an IP address on the public LAN but is physically part of or on the random LAN. The proxy server discards this incoming packet (step 126) because another device in the random LAN will receive the packet by Ethernet. The proxy
20 server simply ignores the packet because it does not need to create a response packet; the intended device physically on the random LAN should respond.

If the incoming packet is not destined for the public LAN, a random network ID (RandNetID) may be determined based on the random IP address and the IP network mask of the random LAN, *e.g.*, logically "anding" the address with the mask (step 128). The
25 proxy server uses RandNetID to discover what devices are local to the random LAN. Devices that not local are reached through a router. A network ID associated with the packet (NewNetID2) is also determined (step 130). This may be determined based on the destination IP address sub-component of the ARP data structure associated with the ARP packet and the IP network mask of the random LAN interface, *e.g.*, logically "anding" the
30 destination address with the random IP network mask. The proxy server uses this information to determine if the IP destination address of the ARP packet is for a "local" network, *i.e.*, the random LAN. RandNetID may then be compared to NewNetID2 (step

- 9 -

132). If RandNetId is equal to NewNetID2, the incoming packet is destined for a device that has an IP address on the random LAN and physically part of or on the random LAN. Again, the proxy server discards the incoming packet (step 134), since another device in the random LAN should respond and will receive the packet by Ethernet.

5 If the incoming packet has not been discarded by the proxy server based on these comparisons, the packet is destined for a device outside of a local network, *i.e.*, from a device in a random LAN, such as a hotel, to a device outside of that LAN and physically part of, for example, a public LAN. Consistent with the present invention, the proxy server prepares an ARP response packet (step 136) to the incoming packet to convince the
10 sending device that the proxy server is the device's gateway. The response ARP packet format is defined in IETF Standard 37.

Referring once again to Figure 9, once an ARP response packet has been created, a proxy ARP address exchange is performed (step 90), as shown in Figure 12. Consistent with the present invention, the IP source address component from the incoming packet is
15 copied into the destination address component of the response packet (step 140). This directs the response packet to the appropriate mobile IP device. Similarly, the source MAC address component of the incoming packet is copied to the destination MAC address of the response packet (step 142). The destination address component of the incoming packet is copied into the source address component of the response packet (step 144). Address
20 exchange consistent with the present invention also contemplates filling in the source MAC address component of the response packet with the MAC address of the random interface (step 146). By inserting the operation component of the response packet with the appropriate value in network-byte order (*e.g.*, the value "2" for RedHat Linux 5.0), the packet is considered a response packet for purposes of the ARP protocol.

25 Once the address exchange is completed, the response packet may be written to the random Ethernet interface using a standard system call (Figure 9, step 92). When the mobile IP device receives the response packet, it will believe proxy server 20 is its gateway from the random LAN to outside public LANs. Subsequent traffic destined for public LANs will be routed there by the proxy server.

30 Proxy DNS

Normally, a mobile network device communicates with a nearby network element commonly referred to as a DNS server. The DNS server functions to translate an IP name

- 10 -

input by a user, such as "undefined.etherloop.com," into a corresponding IP address, such as 137.118.199.33. Thus, when a mobile IP device user inputs an IP name as an intended destination, the device communicates with the DNS server, which then performs a translation, *e.g.*, a name lookup. However, with a random LAN, this is not possible.

5 Figure 13 shows the transmission of a DNS packet 42 from a mobile IP device within random LAN 28 to proxy server 20. Consistent with the present invention, the proxy server pretends that it is the correct DNS server and handles the DNS translation activities. In general, DNS packet 42 includes the IP address of the sending mobile IP device, a destination address, *i.e.*, the address of the DNS server, which is typically
10 preconfigured in the device, and the IP name the device is seeking an IP address for. Proxy server 20 redirects the request to a local process, *i.e.*, a process within the proxy server which, in turn, performs the required translation. After the translation is performed, proxy server 20 generates a DNS response packet 44 (see Figure 14) which includes the address of the sending network device as a destination address and the IP address corresponding
15 to the IP name.

Figure 15 depicts steps for proxy DNS consistent with the present invention. First, the proxy DNS server is initialized (step 150). As shown in Figure 16, during initialization, proxy server 20 creates an unreliable datagram protocol (UDP) datagram socket (step 170) using known methods. This socket is a communications interface
20 between the local process (proxy DNS) and the operating system. Proxy server 20 also establishes a firewall rule such that any packet headed for any destination in any LAN or WAN (*i.e.*, any IP address) from any source in any LAN or WAN with a destination port of "53" is delivered to the proxy server 20 at the same port, *i.e.*, port "53" of the server (step 172). Port "53" is the port for the DNS server of the host machine (a machine with
25 an IP address). The socket can then be bound to port "53" using standard methods (step 174), such that any DNS query with a destination of port "53" is routed to the local process (proxy DNS). The identity of a remote name server physically located outside the random LAN on the Internet is also identified by the proxy server using, for example, a configuration file or some other known method (step 176). After creating a UDP socket
30 association between the proxy DNS server and the identified remote name server (step 178), proxy server 20 enters a loop waiting for new connections, such as new DNS queries

from the random LAN or DNS responses from the remote name server, using standard methods (step 180).

After proxy DNS initialization, a "random" client (*i.e.*, a device in the random LAN having an IP address, such as 1.1.1.1) makes a DNS query to the identified remote name
5 server (*e.g.*, having an IP address, such as 2.2.2.2) at port "53" (step 152) using standard DNS protocol, *e.g.*, IETF standard 13. Proxy server 20 receives the packet and, based on the port address used as the destination port (*i.e.*, port "53"), redirects the packet to the proxy DNS server (the local process within the proxy server) (step 154) using, for example, firewall redirection code built into Linux Redhat 5.0. The proxy DNS server in turn
10 receives the packet and determines the original destination address (*i.e.*, 2.2.2.2) and port (*i.e.*, port "53") of the intended destination, storing them in appropriate variables. In Redhat Linux 5.0, this may be accomplished using an appropriate system call to collect the packet data from the kernel.

The proxy DNS server may then send the DNS query packet to the remote name
15 server identified during initialization (step 158). The proxy DNS server creates a new UDP socket, a communications interface between it and the remote name server. Typically, the proxy DNS server uses a separate socket and port (such as port "2001," in this example) which may be arbitrarily assigned) for each mobile device IP address so as to be able to identify to which device in the random LAN the response should be sent. The remote
20 name server then processes and responds to the DNS query from the proxy DNS server, as defined in IETF Standard 13 (step 160). A DNS response packet, which includes the requested address, is generated by the remote name server and sent to the proxy server using the port defined in step 158 (*i.e.*, port "2001" in this example). Proxy server (such as proxy server 20) receives the DNS response from the remote name server on the
25 specified port (*i.e.*, port "2001") using, for example, standard Unix system calls and determines the IP address of the client (*i.e.*, 1.1.1.1 in this example) using that port (step 164). The DNS response is then sent to the client by the proxy server (step 166; see also Figure 14), which performs a source address and port modification. In one embodiment, the proxy server modifies the source address (the address of the proxy server, *e.g.*, 3.3.3.3.
30 in this example) and source port of the response packet (the port of the link back to the client, which may be arbitrarily assigned) to the original destination address (*i.e.*, 2.2.2.2.

- 12 -

in this example) and the original destination port (*i.e.*, port "53" in this example) of the DNS query packet, respectively.

One-Armed Proxy Server

The TA proxy server described above is extremely well-suited for operation in any
5 environment where a fully random assortment of users may attempt to connect to the LAN. As previously noted, one such environment is a hotel environment. Today, hotel guests frequently have mobile network devices, such as laptops, and wish to connect via Ethernet, EtherLoop, or the like, into the hotel's network and from there to the Internet to retrieve electronic mail and conduct other business.

10 A TA proxy server, however, may not fit well into a LAN networking community since most end-user network devices will have a stable IP address, correctly configured and assigned for the LAN, unlike the hotel environment. Moreover, a TA proxy server increases latency (delay) on a network. For at least these reasons, there is a need for proxy routers that do not interfere with the normal operation of a LAN. These services will exist
15 on a normally configured LAN and only begin operation only when a random interloper, *i.e.*, an individual device, connects to the network. The device that supports this service will be lower in cost and will not create any performance problems for the standard network traffic.

Figure 17 illustrates an alternative proxy server implementation consistent with the
20 present invention that satisfies this need. In this implementation, LAN 200 includes a plurality of mobile IP network devices 202, 204, and 206 which communicate with a router 210 through hub 212. The links used to interconnect the various network elements shown may, for example, be Ethernet links. Proxy server 214 communicates with the various network elements through a single link, *e.g.*, a single Ethernet interface. As such, proxy
25 server 214 may be referred to as one-armed proxy server. Generally, proxy server 214 includes the same hardware as proxy server 20 (Figure 2). The same Ethernet interface receives data from LAN 200 and delivers translated data back to the LAN. In Figure 17, the majority of the devices on the network are on the same network as the router. These devices operate normally without any interference from proxy server 214, which is always
30 listening to the traffic on this network. Device 208, an interloper with an address of 47.31.128.195 will not be able to directly communicate with the router, which ignores all traffic from an address that is not in its network domain (*i.e.*, 137.118.199.X). One of the

- 13 -

benefits of this implementation is that the standard network traffic is not interrupted by the proxy server, as shown in Figure 18. As with the TA proxy server of Figure 1, one-armed proxy server 214 performs network address translation, proxy ARP, and proxy DNS services to provide similar fully transparent client functionality to any IP-enabled network device.

Network Address Translation

Figure 19 depicts a network environment useful for discussing NAT performed by one-armed proxy server 64. Network 220, an Ethernet type of network, includes a plurality of interconnected network elements including router 222, workstation 224, hub 226, and proxy server 214. Interloper 228 is also connected to LAN 220, although it is a foreign IP-enabled network device relative to the network, *i.e.*, it is not part of LAN 220.

As shown in Figure 19, interloper 228 generates a packet 230 for a LAN/WAN other than the LAN 220. Packet 230 includes the IP address of the transmitting device, interloper 228, as well as an IP address of the destination device, which is not particularly shown. Due to the nature of Ethernet networks, every element on the LAN receives packet 230 from hub 226. The IP protocols within workstation 224 and similar network elements recognize that the destination is not a local network device and consequently ignore the packet. Router 222 may or may not accept packet 230 depending on the source IP address. Even if the router accepts the packet and passes it on to the LAN/WAN, the packet will not return to the router (address 137.118.199.1) since the source address is 47.31.128.195.

Proxy server 214, however, recognizes that it is capable of properly translating packet 230 into translated packet 232 having an acceptable format utilizing known address translation methods (see Figure 20). Since translated packet 232 has an associated IP address recognizable by router 222, the router believes packet 232 originated from within LAN 220 and not by a foreign mobile device, *i.e.*, interloper 228. As shown, packet 232 includes the same destination address as packet 230. Once router 222 receives this packet, it can send the packet on to the Internet as normal.

After the remote device receives the translated packet, it may generate a response packet. If so, the response packet 234 must be translated by proxy server 214 so that it can be delivered to interloper 228 (see Figure 21). Router 222 knows that the .55 IP address is associated with a device on its network. In this case, the device happens to be proxy server 214, but the router is unaware of the presence of the server nor does it matter to the

- 14 -

router that the destination is the proxy server and not a "normal" network device such as a workstation. Instead, router 222 simply forwards response packet 234 to proxy server 214, just like it would forward any other packet to other network devices. At about the same time the proxy server receives response packet 234, interloper 228 also receives the response packet from hub 226. However, since the interloper 228 knows that its IP address is 47.31.128.195, it throws the response packet away. Consistent with the present invention, proxy server 214 receives the response packet, performs a reverse network address translation, and sends a translated packet 236 back out on the LAN through hub 226 (see Figure 22). Packet 236 is broadcast across the entire LAN, but since only one device on the network has IP address 47.31.128.195 (interloper 228), only that device will not discard the packet. Proxy server is able to specifically target the interloper by using the interloper's medium access control (MAC) address as the destination.

Proxy ARP

Figures 23 and 24 show a network environment useful for discussing the proxy ARP capabilities of proxy server 214. Interloper 228 may generate an ARP packet 238 in order to discover the MAC Address of its gateway device, *i.e.*, 47.31.128.1. Normally, since no device in LAN 220 has the appropriate MAC address, packet 238 would be ignored and interloper 228 would be unable to function. Consistent with the present invention, proxy server 214 will however recognize that packet 238 does not belong on the 137.118.199.X network and will automatically generate a response (see Figure 24). Response packet 240 includes, as a destination address, the IP address of interloper 228, as well as a reply to the gateway query. Once interloper 228 receives packet 240, it considers the proxy server 214 to be its gateway device and will use the server for all further communications outside of the local LAN. The steps for proxy ARP performed by TA proxy server 20 discussed above are equally applicable to proxy server 214.

Proxy DNS

Proxy server 214 is also capable of performing proxy DNS. If, from the point of view of the interloper, the DNS server is usually on the same LAN as the interloper, the interloper will generate an ARP request for the DNS server. Consistent with the present invention, the proxy server 214 will respond to the ARP request with its own address. Future DNS queries will be delivered directly to proxy server 214, which can then answer them. Similarly, if, from the point of view of the interloper, the DNS server is outside of

- 15 -

the local LAN, *i.e.*, on a WAN, the proxy server will automatically receive the DNS query, since it is the interloper's gateway. In this case, proxy serve 214 will see the DNS query packets arrive and will be able to response to them locally. The steps for proxy DNS performed by TA proxy server 20 discussed above are equally applicable to the proxy DNS
5 service performed proxy server 214.

In addition to the features described above, proxy servers consistent with the present invention support certain security functions to improve network administration. For example, each time an interloper connects to an proxy server-enabled network, the proxy server will be able to provide connectivity for that user. To improve the security of
10 the network, the proxy server will deliver a message to a specified network administrator e-mail account to the alert the administrator to the presence of this new user. While this is not ironclad security, it is a reasonable first step in network security.

Properly configured, proxy servers consistent with the present invention can provide interlopers with a secondary gateway. This has at least two benefits, including
15 reduced congestion on the standard router and improved control over the interloper's internet access. Reduced congestion is a relatively straightforward concept, *i.e.*, by using a different router than the standard network traffic, it reduces the possibility of excessive demand on router resources, that in turn might affect the performance of the standard network users. Further, by specifying a secondary gateway, the network administrator can
20 funnel interlopers into a less open corporate environment, preventing those users from reaching sensitive material within the standard corporate network. This is one of the major benefits of the present invention over the straight DHCP model, for two reasons. First, if the network is served by a proxy server, a stand-alone DHCP server is not needed. The standard network users will not need to use DHCP on a day-to-day basis. Given the first
25 constraint, everyone who uses DHCP is, by implication, an interloper and can be treated with additional security restrictions.

As for other protocols, one of ordinary skill will appreciate that proxy servers consistent with the present invention are only capable of supporting IP-based translation services. This is primarily because IP is both a LAN and WAN protocol. Other common
30 LAN protocols, such as Apple Talk, or IPX are significantly limited in scope, and are not capable of the "long range" communications that make the proxy server translation services possible. However, it may be possible for "bridges" to be built to allow IPX-based

ENCLOSURE TWO 000000000000

- 17 -

We claim:

1. A method, for use with a proxy server, for establishing communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network, the method comprising the steps of:

5 generating an address resolution protocol packet to identify the arbitrary address for the destination device;

receiving, by the proxy server, the address resolution protocol packet;

generating an address resolution protocol response packet including the arbitrary address of the destination device; and

10 transmitting the address resolution protocol response packet from the proxy server to the device in the first network.

2. The method of claim 1 further comprising the step of:

determining if the proxy server should respond to the address resolution protocol

15 packet.

3. The method of claim 2 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:

determining if the address resolution protocol packet is destined for a device that

20 has an address outside of the first network but that is part of the first network; and

discarding the address resolution protocol packet if it is destined for a device that has an address outside of the first network but that is part of the first network.

4. The method of claim 3 wherein the step of determining if the proxy server

25 should respond to the address resolution protocol packet includes the substeps of:

determining if the address resolution protocol packet is destined for a device that has an address on the first network and that is part of the first network; and

discarding the address resolution protocol packet if it is destined for a device that has an address on the first network and that is part of the first network.

30

5. The method of claim 4 wherein the step of generating an address resolution protocol response packet includes the substep of:

- 18 -

generating a response packet if the address resolution protocol packet is destined for a device that is outside of the first network.

6. The method of claim 5 further comprising the step of:

5 performing an address exchange to direct the response packet to the appropriate device.

7. The method of claim 6 wherein the step of performing an address exchange includes the substeps of:

10 writing an internet protocol source address component of the address resolution protocol packet into a destination address component of the response packet;

writing a source MAC address component of the address resolution protocol packet into a destination MAC address component of the response packet; and

15 writing a MAC address associated with the first network into a source MAC address component of the response packet.

8. A system for establishing communications between a device in a first network and a destination device having an arbitrary address on a second network outside of the first network, the system comprising:

20 a proxy server in communication with the first network and the second network, the proxy server including:

memory for receiving an address resolution protocol packet generated by the device in the first network; and

25 a processor, in communication with the memory, for determining if the proxy server should respond to the address resolution protocol packet with an address resolution protocol response packet, the processor also for generating the response packet and transmitting the response packet to the device in the first network.

9. A method, for use with a proxy server, for establishing communications
30 between a random device and a destination device having an arbitrary address on a network, the method comprising the steps of:

- 19 -

generating an address resolution protocol packet to identify the arbitrary address for the destination device;

receiving, by the proxy server, the address resolution protocol packet;

generating an address resolution protocol response packet including the arbitrary
5 address of the destination device; and

transmitting the address resolution protocol response packet from the proxy server to the random device.

10 10. The method of claim 9 further comprising the step of:
determining if the proxy server should respond to the address resolution protocol packet.

11. The method of claim 10 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:
15 determining if the address resolution protocol packet is destined for a device that has an address associated with the network but that is physically not part of the network;
and

discarding the address resolution protocol packet if it is destined for a device that has an address associated with the network but that is physically not part of the network.
20

12. The method of claim 11 wherein the step of determining if the proxy server should respond to the address resolution protocol packet includes the substeps of:
determining if the address resolution protocol packet is destined for a random device that is not part of the network; and
25 discarding the address resolution protocol packet if it is destined for a random device that is not part of the network.

13. The method of claim 12 wherein the step of generating an address resolution protocol response packet includes the substep of:
30 generating a response packet if the address resolution protocol packet is destined for a device that is part of the network.

- 20 -

14. The method of claim 13 further comprising the step of:
performing an address exchange to direct the response packet to the appropriate device.

5 15. The method of claim 14 wherein the step of performing an address exchange includes the substeps of:

writing an internet protocol source address component of the address resolution protocol packet into a destination address component of the response packet;

writing a source MAC address component of the address resolution protocol packet
10 into a destination MAC address component of the response packet; and

writing a MAC address associated with the random device into a source MAC address component of the response packet.

16. A system for establishing communications between a random device and
15 a destination device having an arbitrary address on a network, the system comprising:

a proxy server in communication with the random device and the destination device, the proxy server including:

memory for receiving an address resolution protocol packet generated by the random device; and

20 a processor, in communication with the memory, for determining if the proxy server should respond to the address resolution protocol packet with an address resolution protocol response packet, the processor also for generating the response packet, and transmitting the response packet to the random device.

25 17. A method, for use with a proxy server, for communicating with a mobile device and remote name server and for obtaining an internet protocol address from the remote name server for the mobile device, the method comprising the steps of:

generating a query packet including a request for an address associated with a domain name;

30 receiving the query packet from the mobile device in the proxy server;

forwarding the query packet to the remote name server;

generating a response packet including the requested address;

- 21 -

transmitting the response packet to the proxy server; and
transmitting the response packet to the mobile device.

18. The method of claim 17 wherein the query packet includes a destination
5 port and wherein the step of forwarding the query packet to the remote name server
includes the substeps of:

directing the query packet to a proxy DNS server based on the destination port;
determining, by the proxy DNS server, the address of the remote name server and
the destination port; and
10 identifying a port associated with the mobile device.

19. The method of claim 18 wherein the step of transmitting the response
packet to the proxy server includes the substeps of:

receiving the response packet on the port associated with the mobile device; and
15 determining, by the proxy server, the address of the mobile device based on the port
associated with the mobile device.

20. The method of claim 19 wherein the response packet includes a source
address and a source port, and wherein the step of transmitting the response packet to the
20 mobile device includes the substeps of:

modifying the source address of the response packet to the address of the remote
name server; and
modifying the source port of the response packet to the destination port.

21. A system for obtaining an internet protocol address from a remote name
25 server for a mobile device, the system comprising:

a proxy server in communication with the mobile device and the remote name
server, the proxy server including:

memory for receiving an address query packet generated by the mobile
30 device; and

- 22 -

a processor, in communication with the memory, for forwarding the query packet to the remote name server and for transmitting a response packet, generated by the remote name server and including the requested address, to the mobile device.

5 22. A method, for use with a proxy server, for providing communications between a random device and a destination device in a network, the method comprising the steps of:

performing a proxy address resolution protocol to initiate communications between the random device and the destination device;

10 performing a proxy domain name service to identify a destination address in the second network associated with a domain name; and

performing a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the destination device in the network.

15

23. A system for providing communications between a random device and a device in a network, the system comprising:

a proxy router including

memory; and

20 a processor programmed to

(a) perform a proxy address resolution protocol to initiate communications between the random device and the device in the network;

(b) perform a proxy domain name service to identify a destination address in the network based on a destination name; and

25 (c) perform a network address translation of an arbitrary random address associated with traffic from the random device to an appropriate address for routing the traffic to the device in the network.

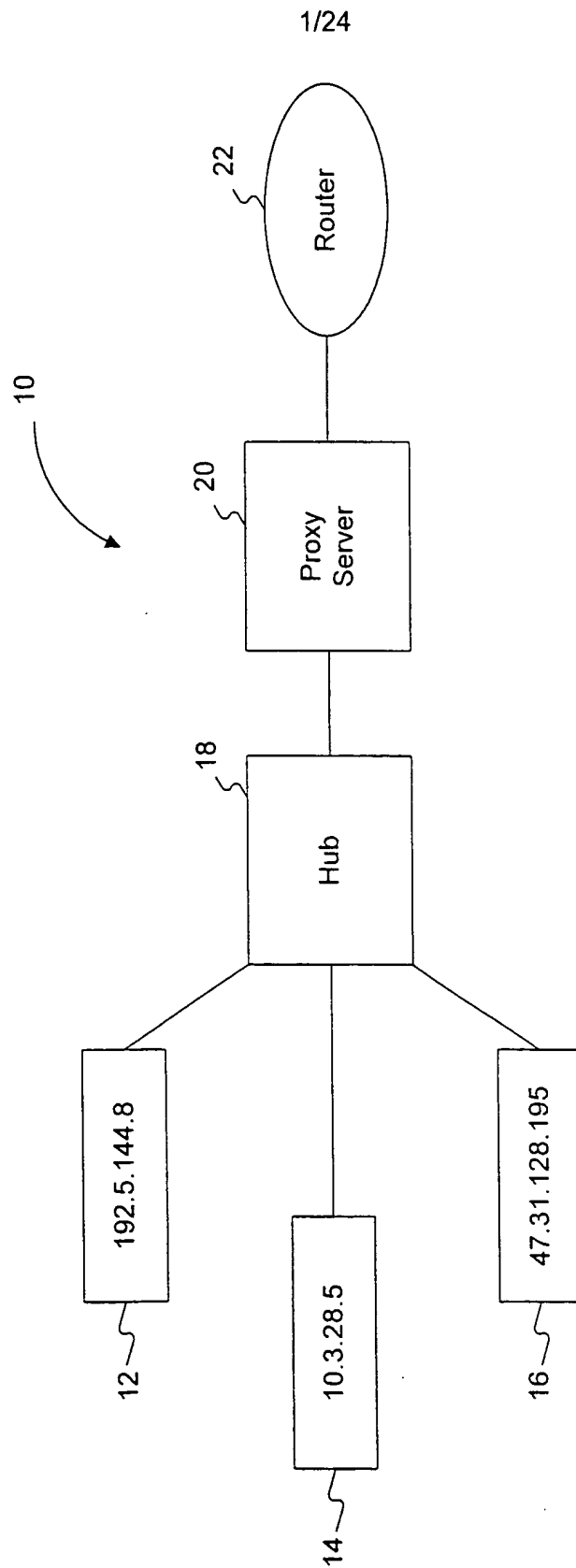


Figure 1

2/24

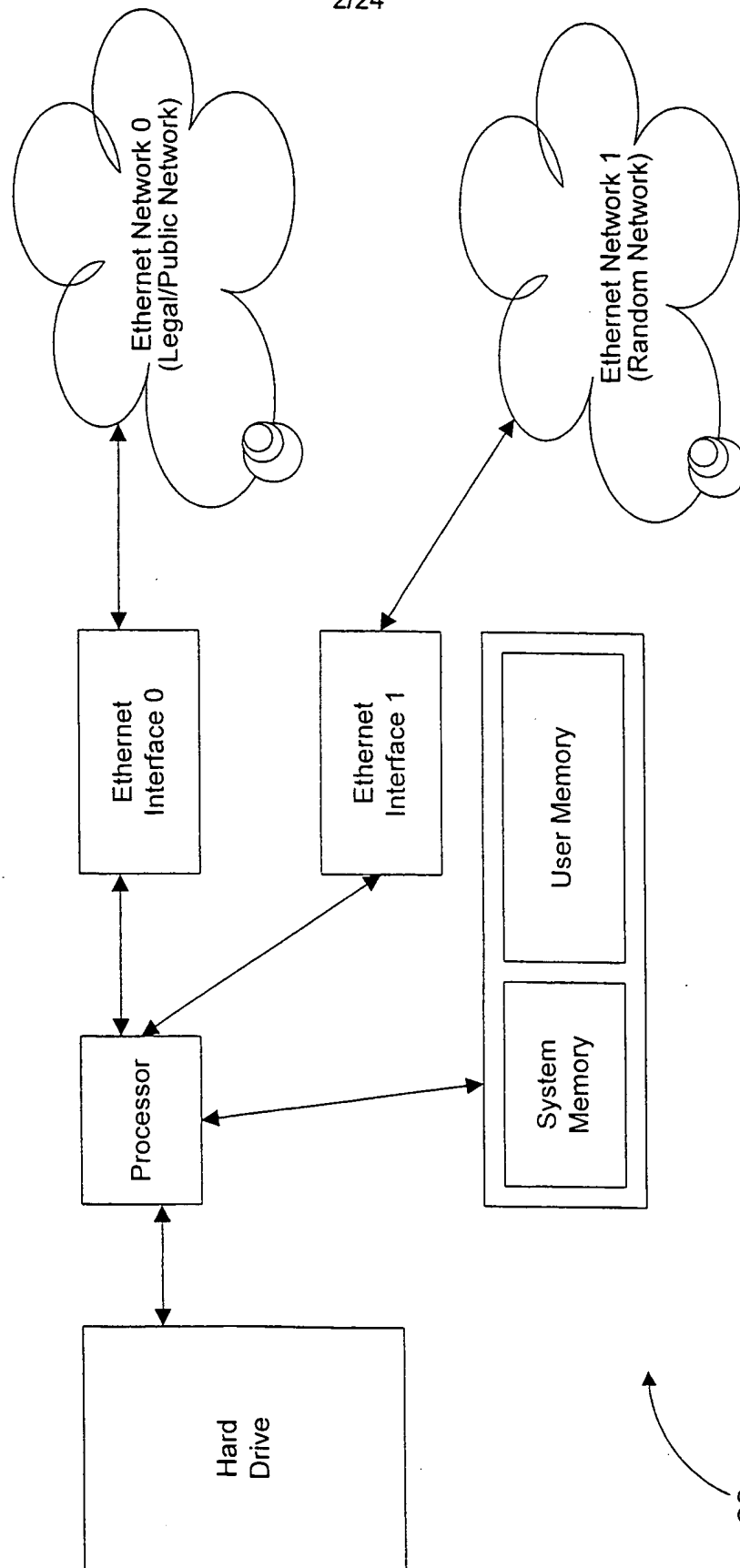


Figure 2

3/24

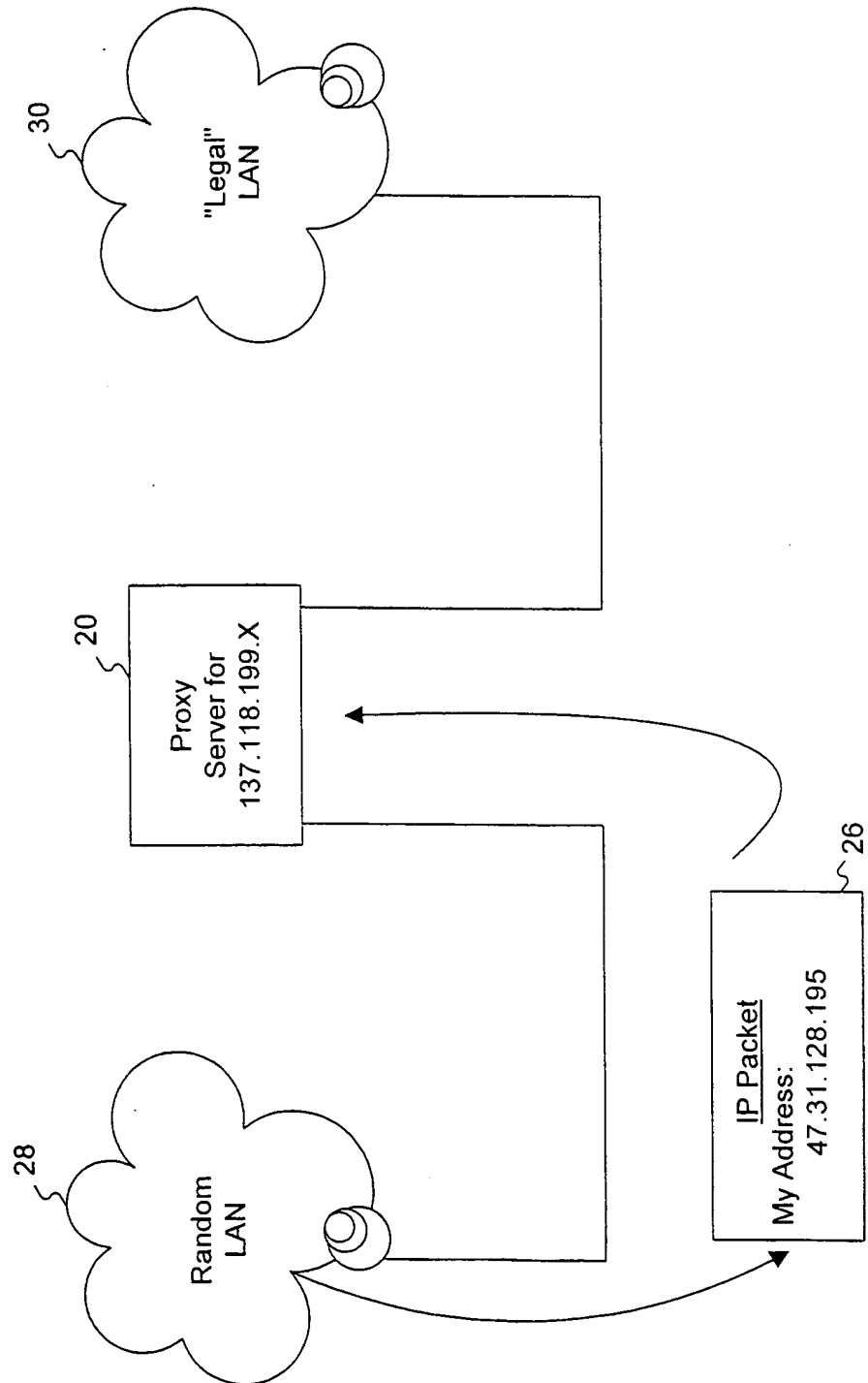


Figure 3

4/24

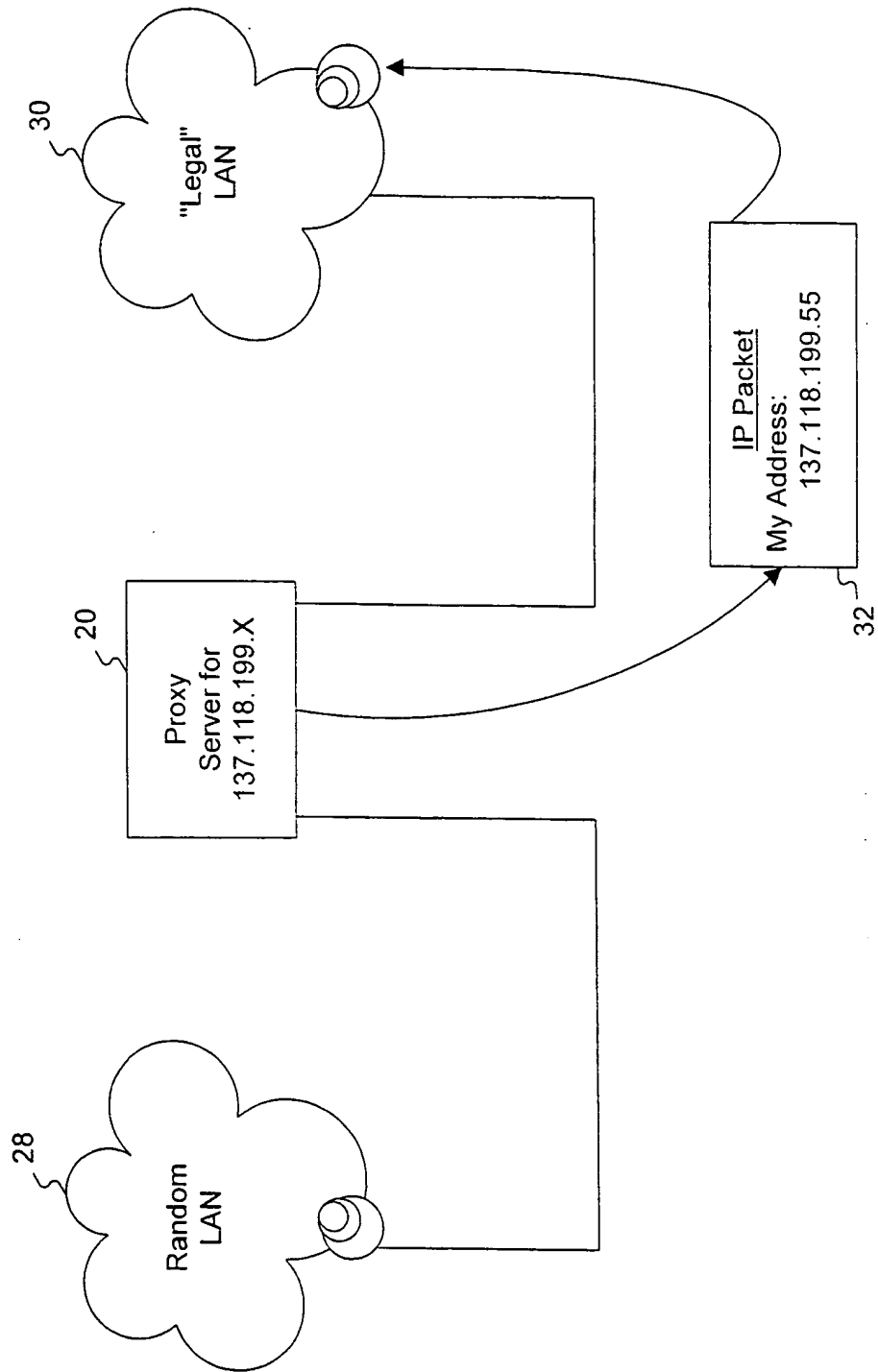


Figure 4

5/24

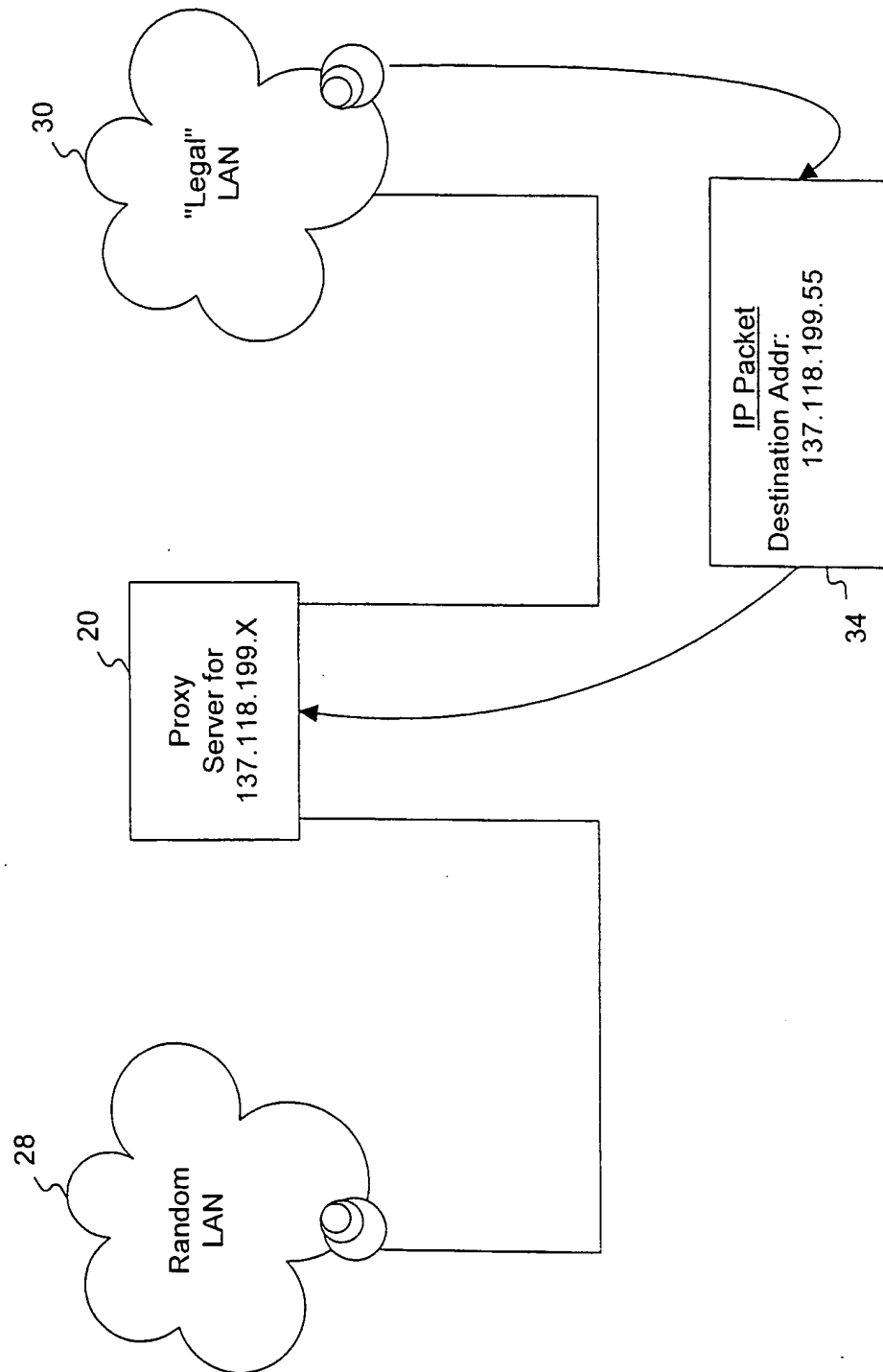


Figure 5

6/24

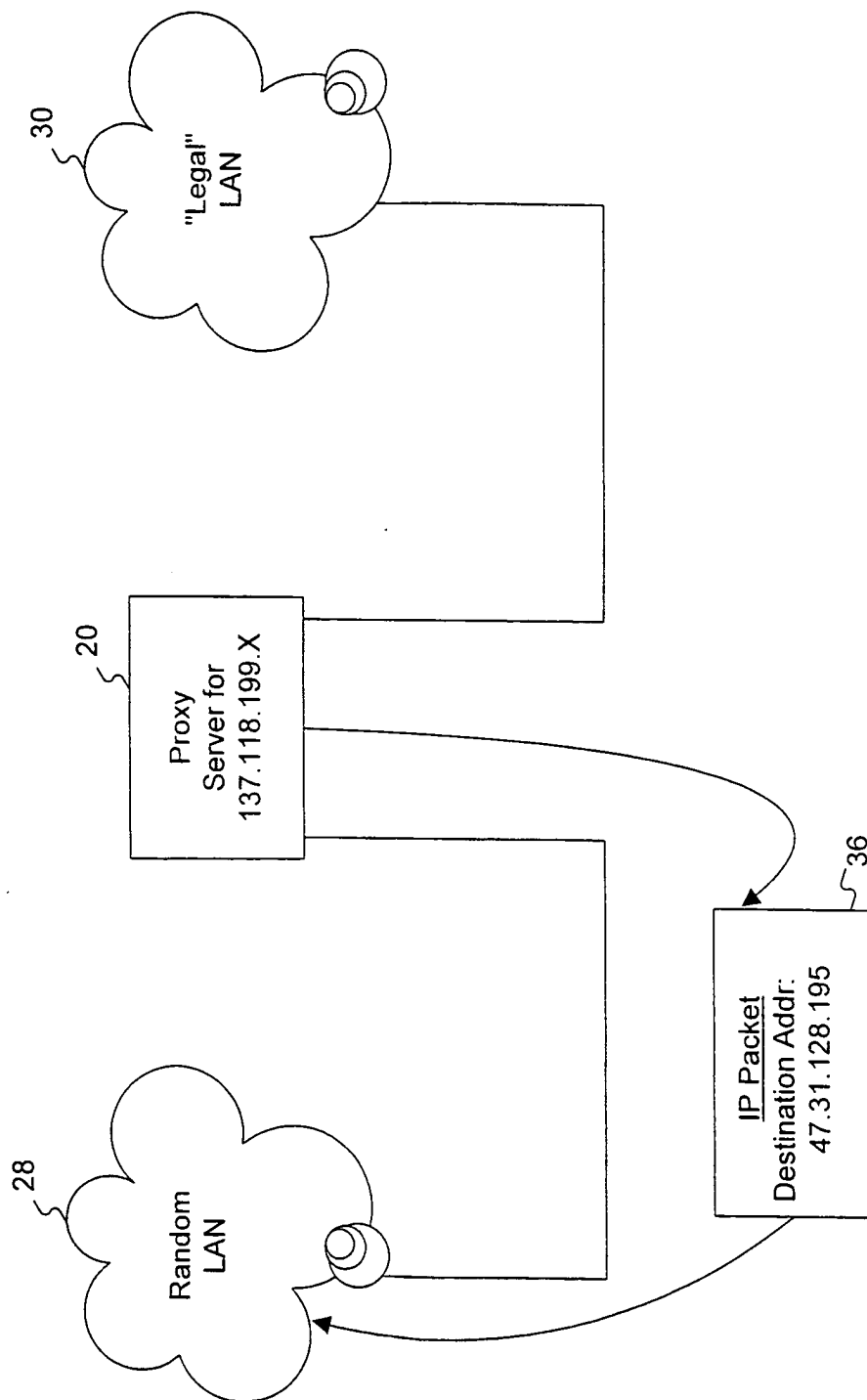


Figure 6

7/24

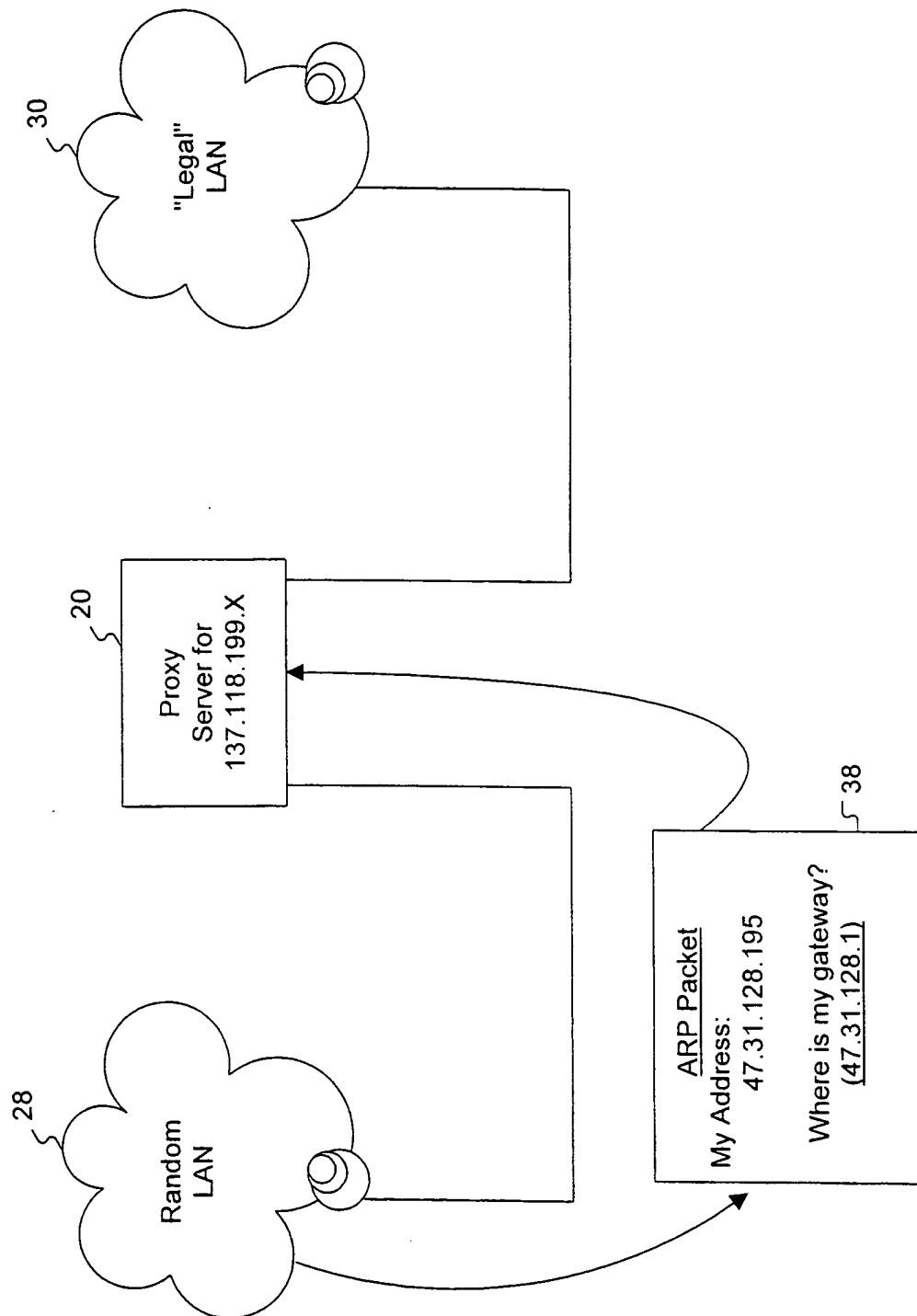


Figure 7

8/24

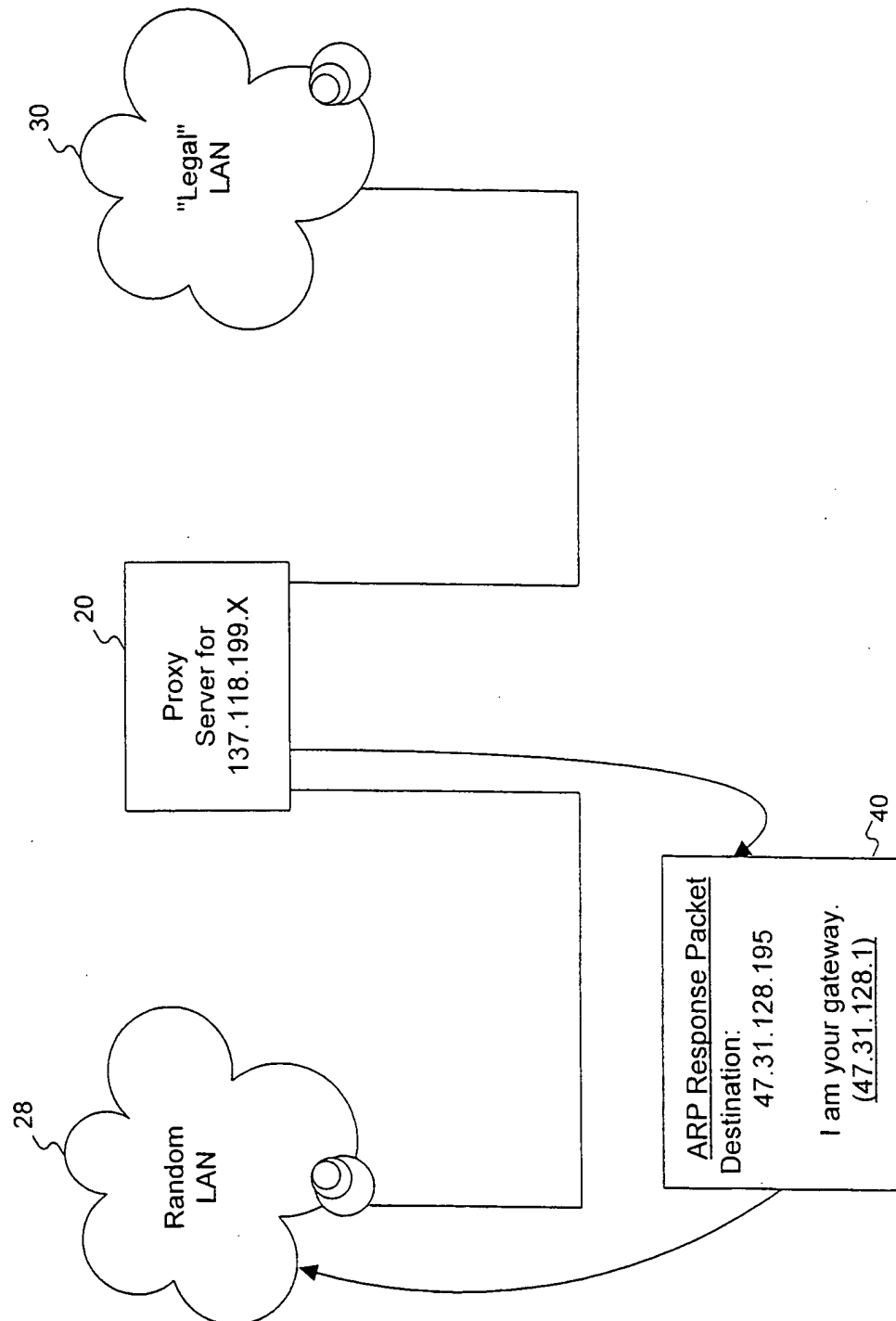
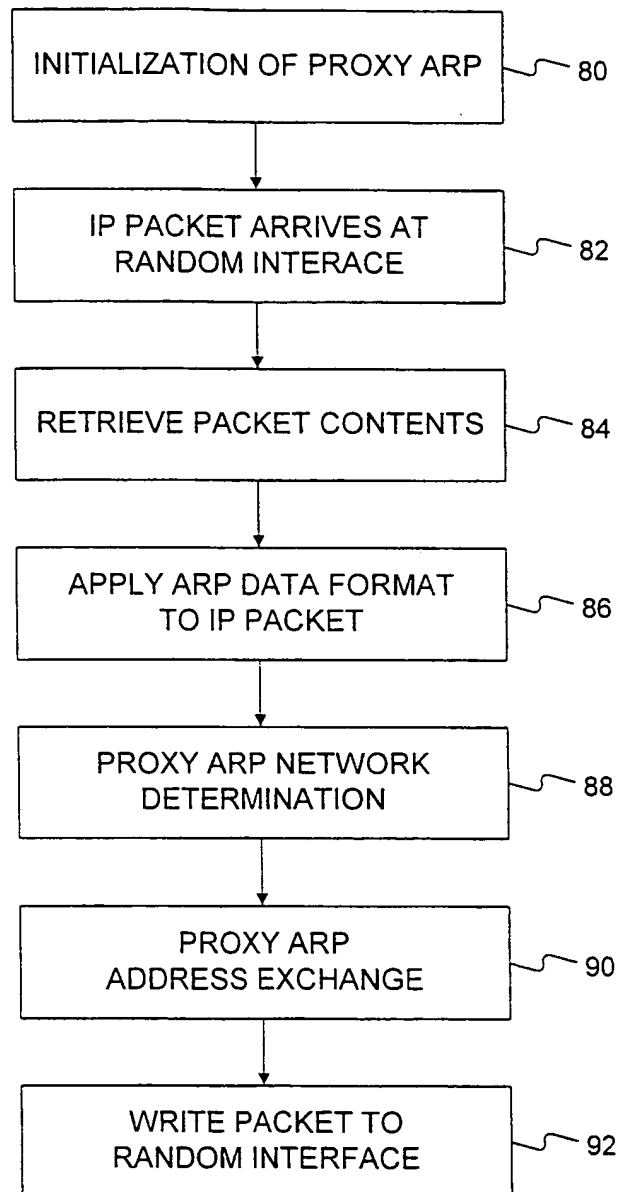
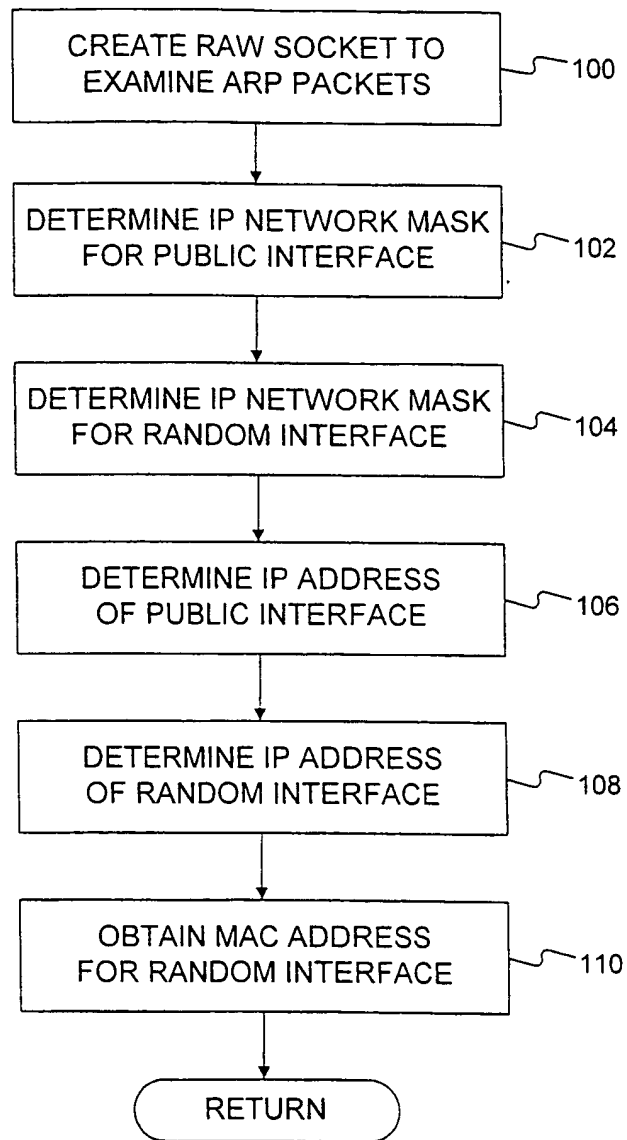


Figure 8

9/24

**Figure 9**

10/24

**Figure 10**

11/24

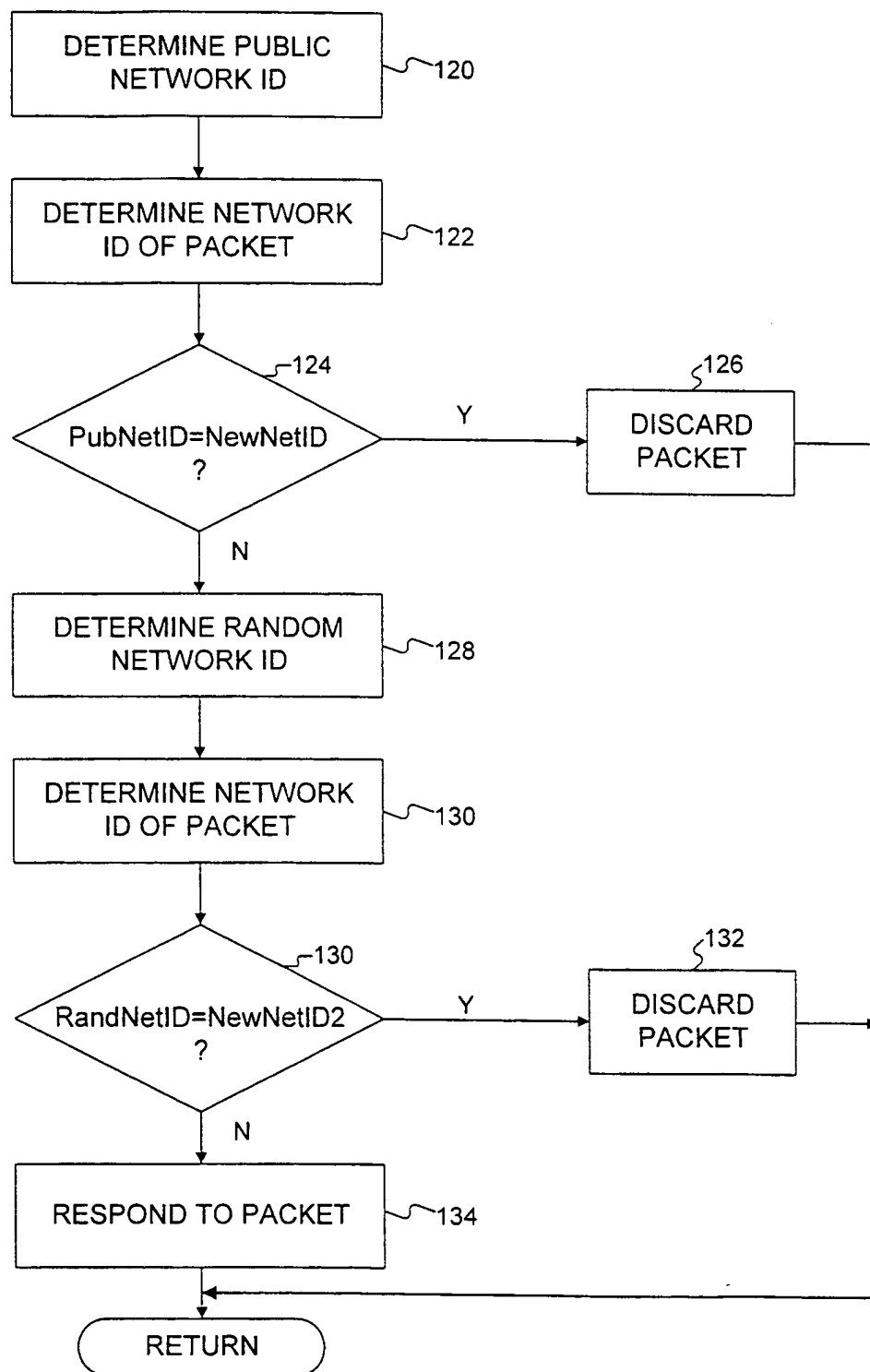
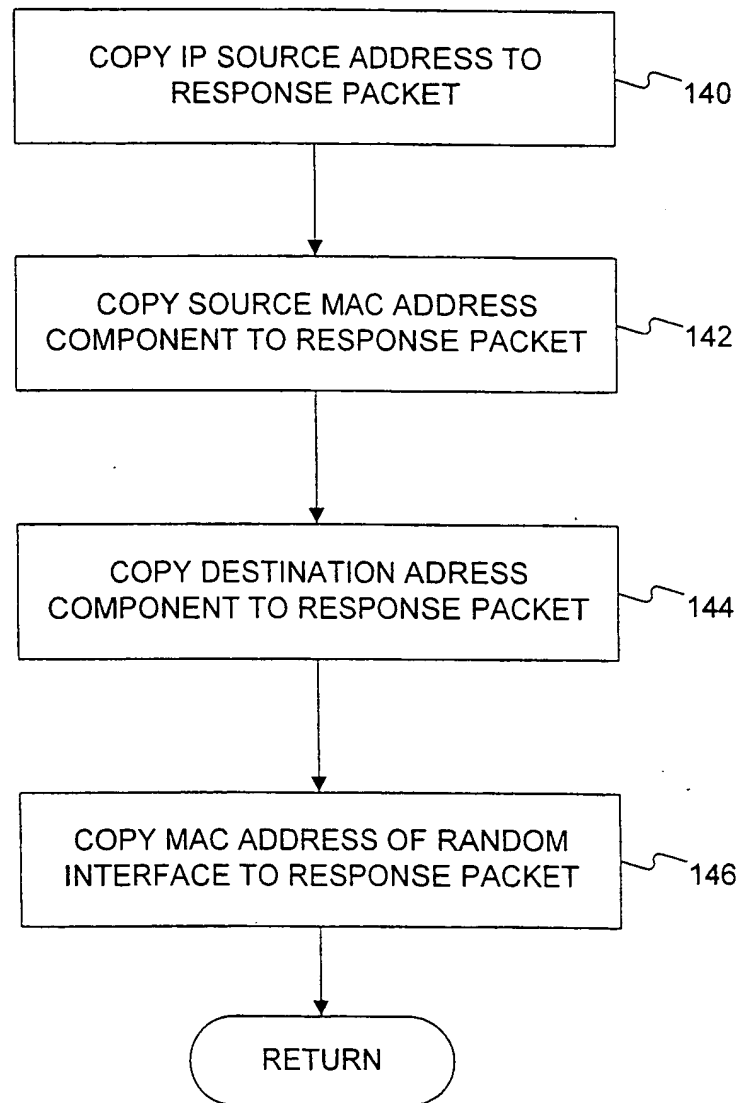


Figure 11

12/24

**Figure 12**

13/24

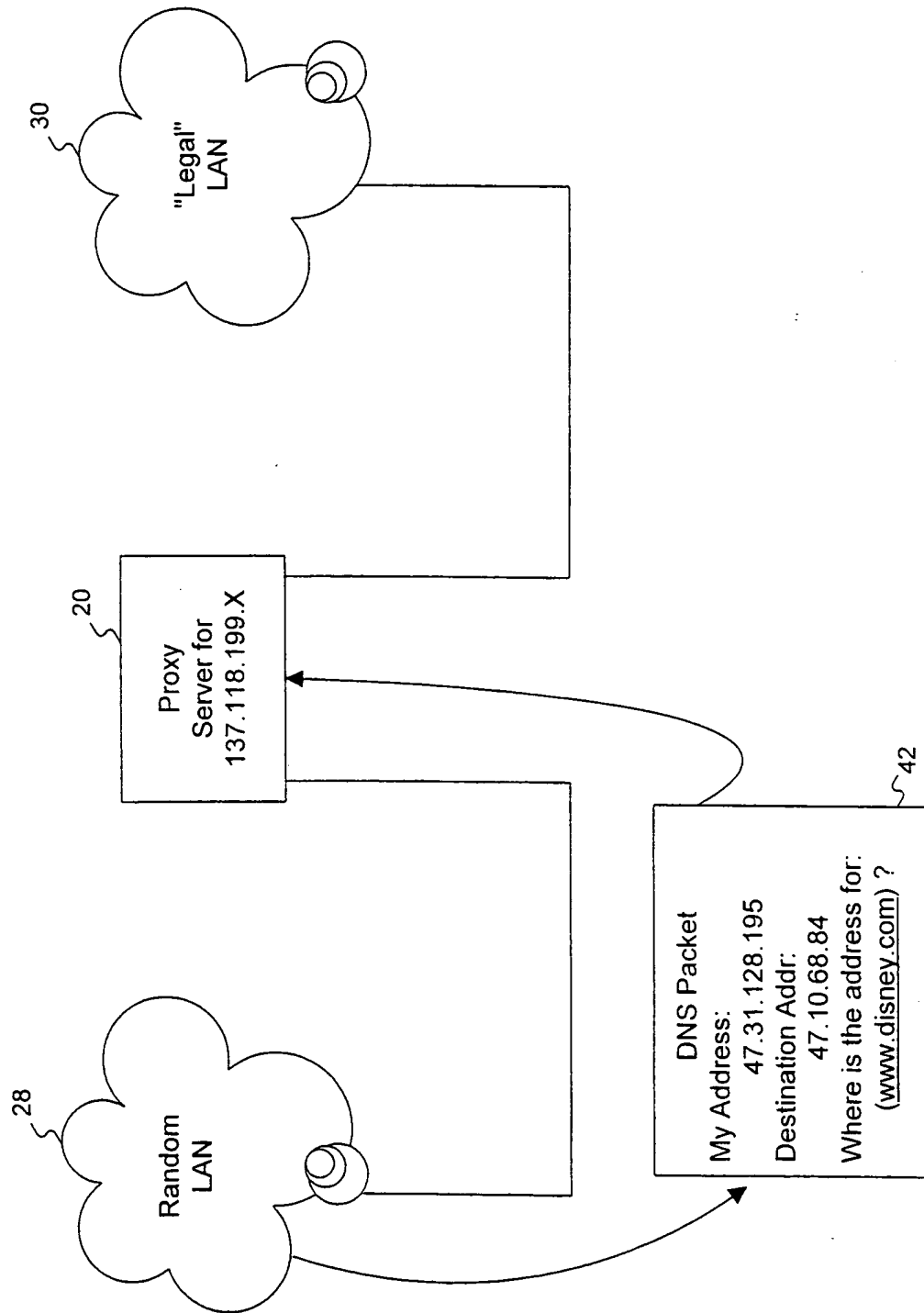


Figure 13

14/24

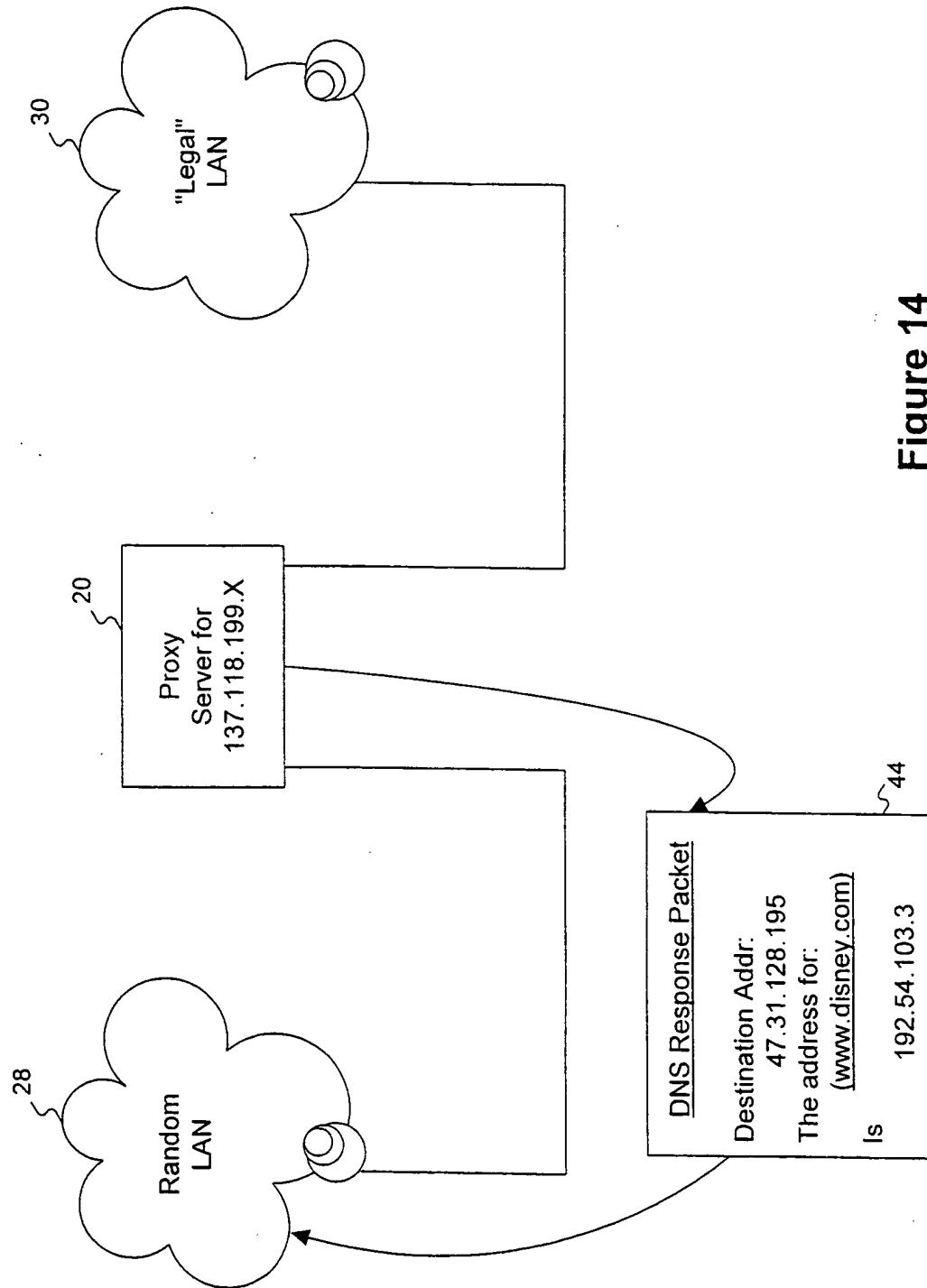
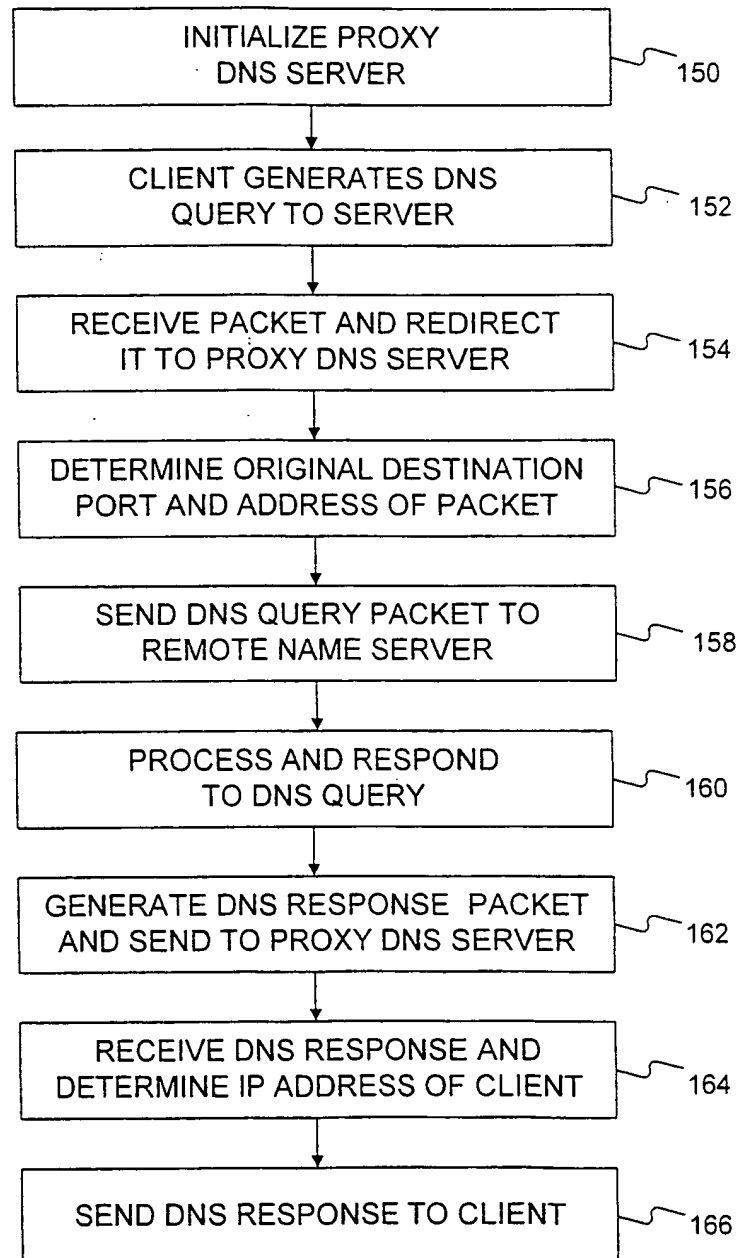
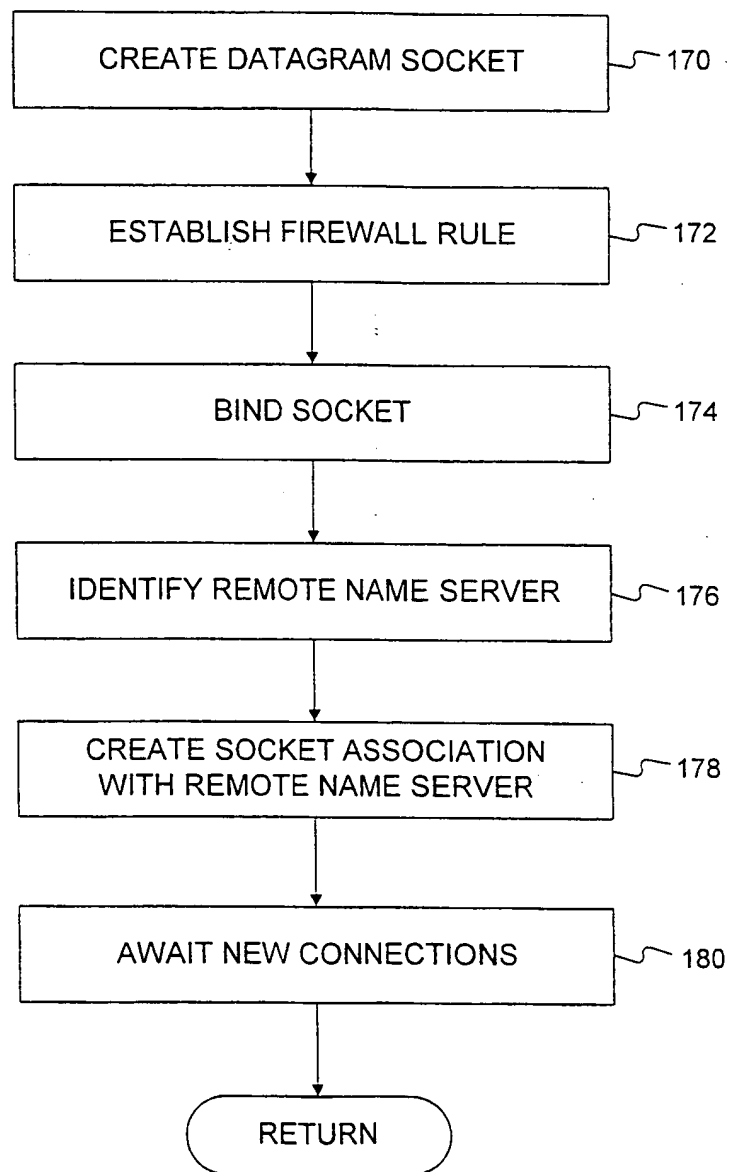


Figure 14

15/24

**Figure 15**

16/24

**Figure 16**

17/24

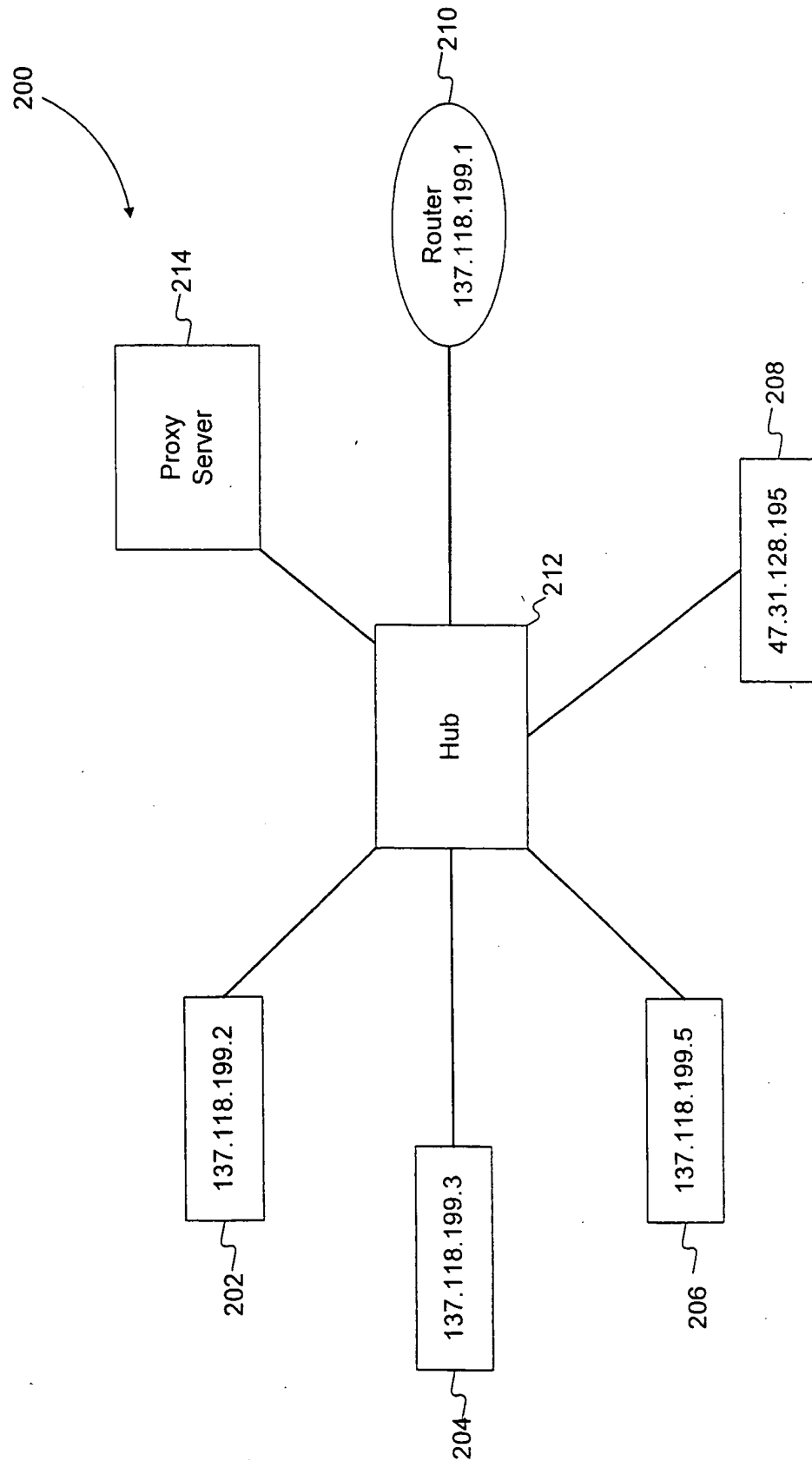


Figure 17

18/24

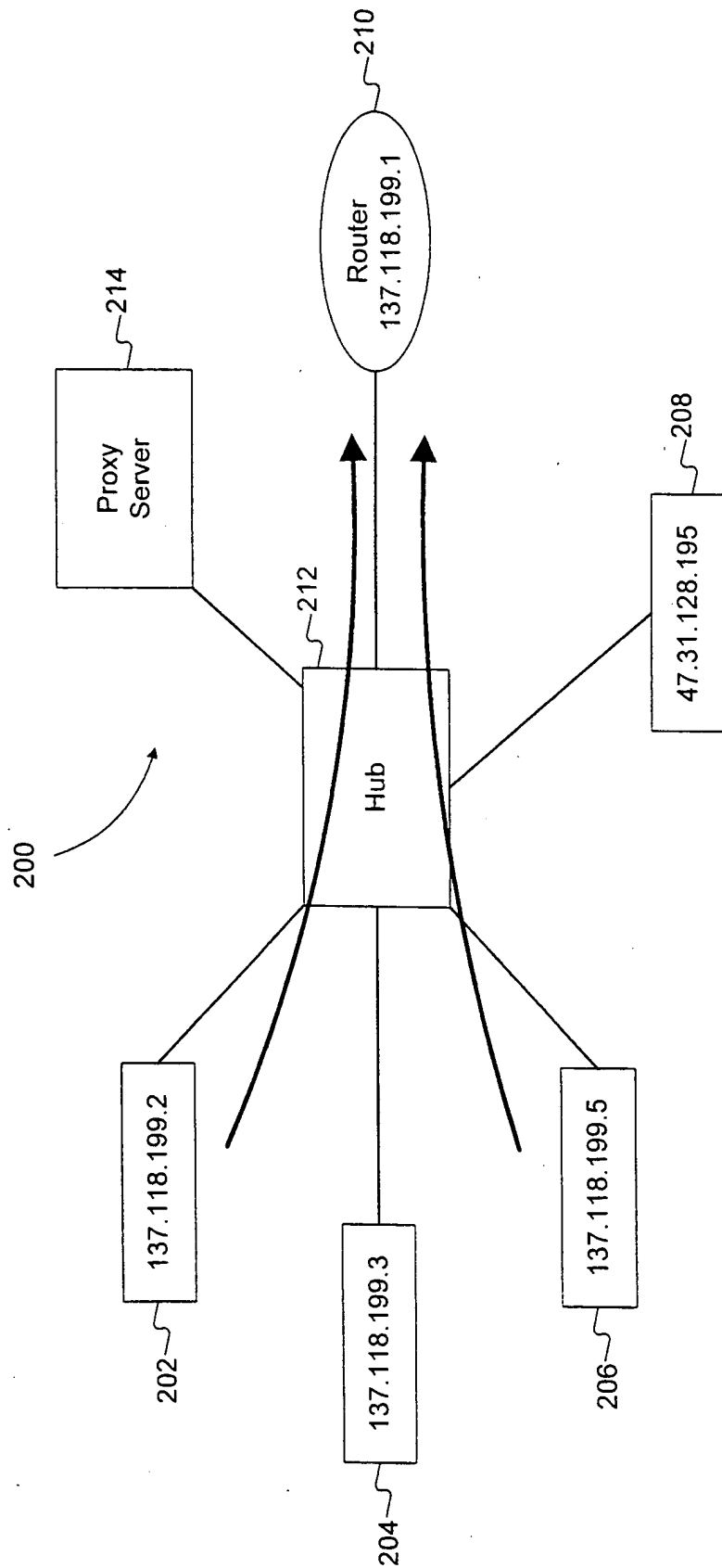


Figure 18

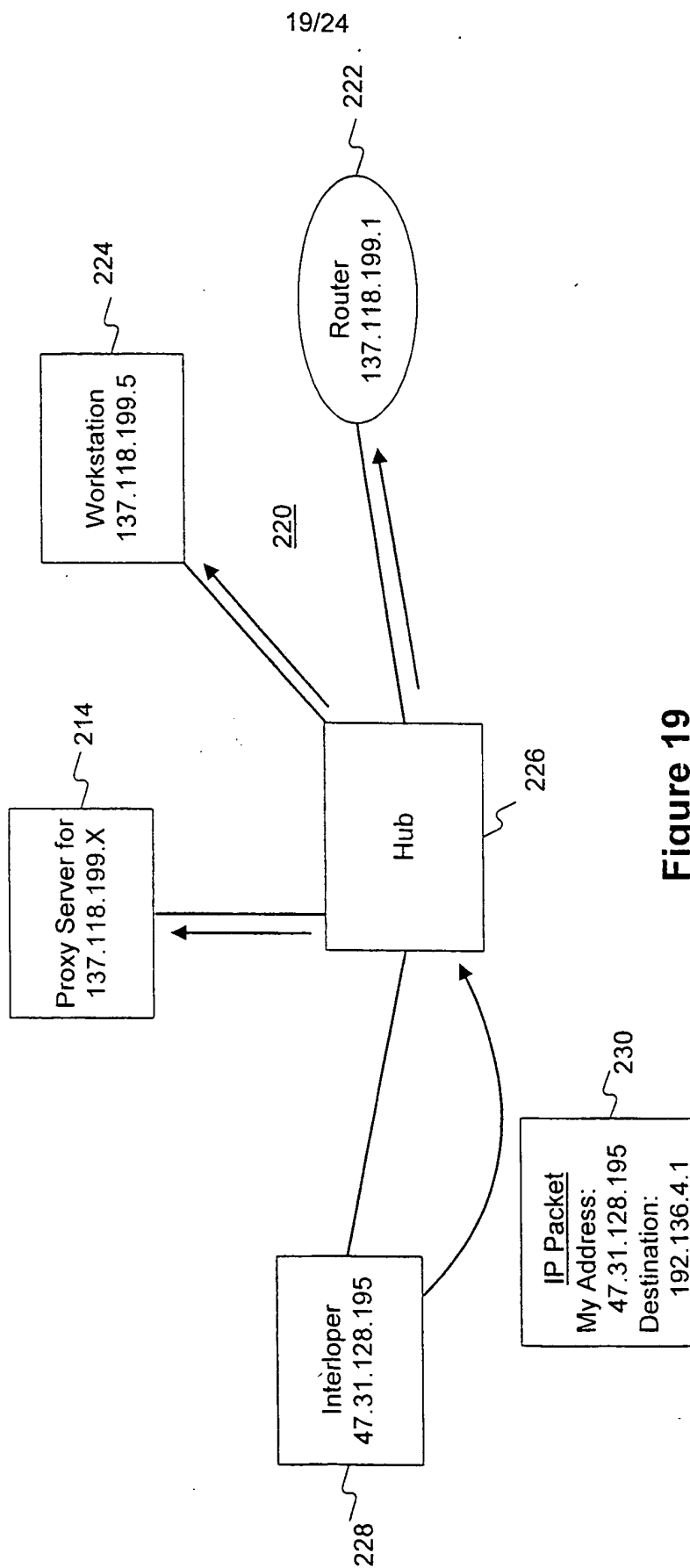


Figure 19

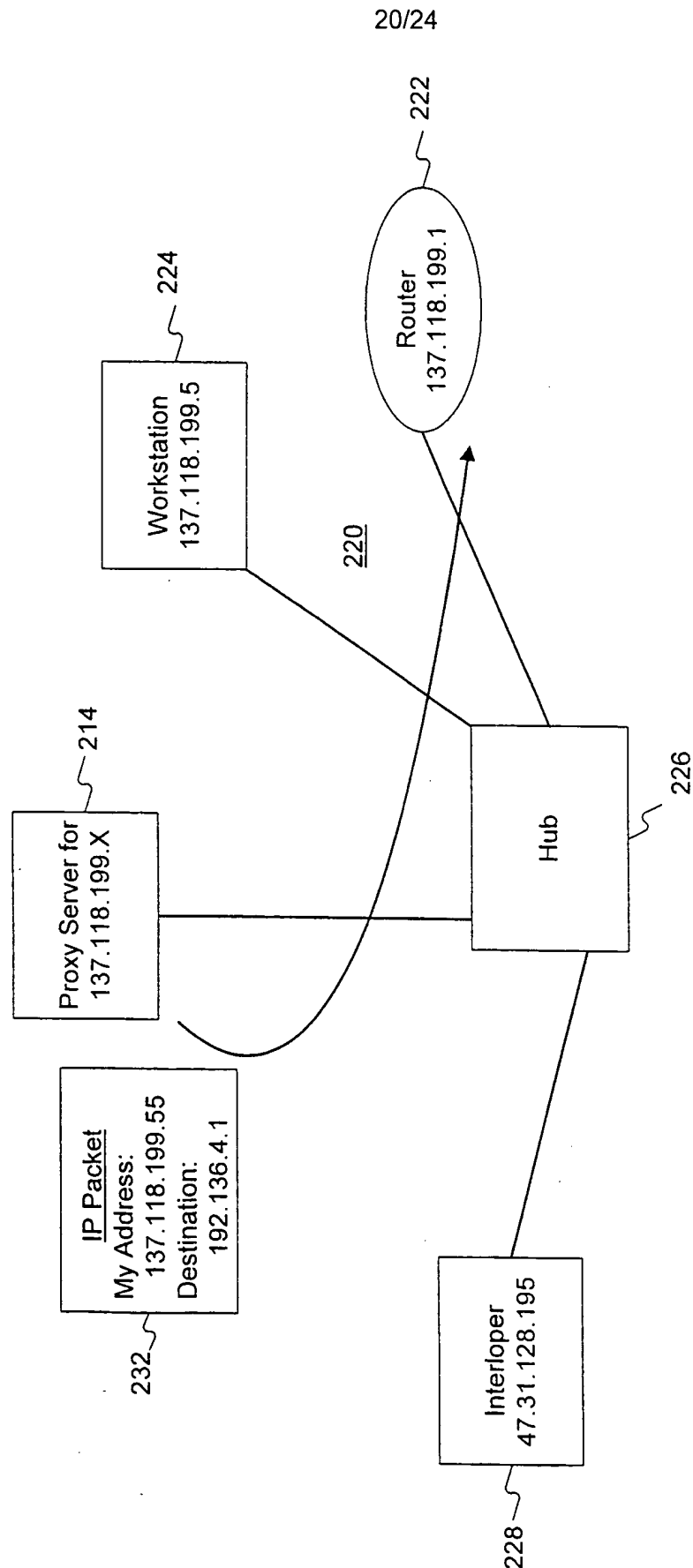


Figure 20

21/24

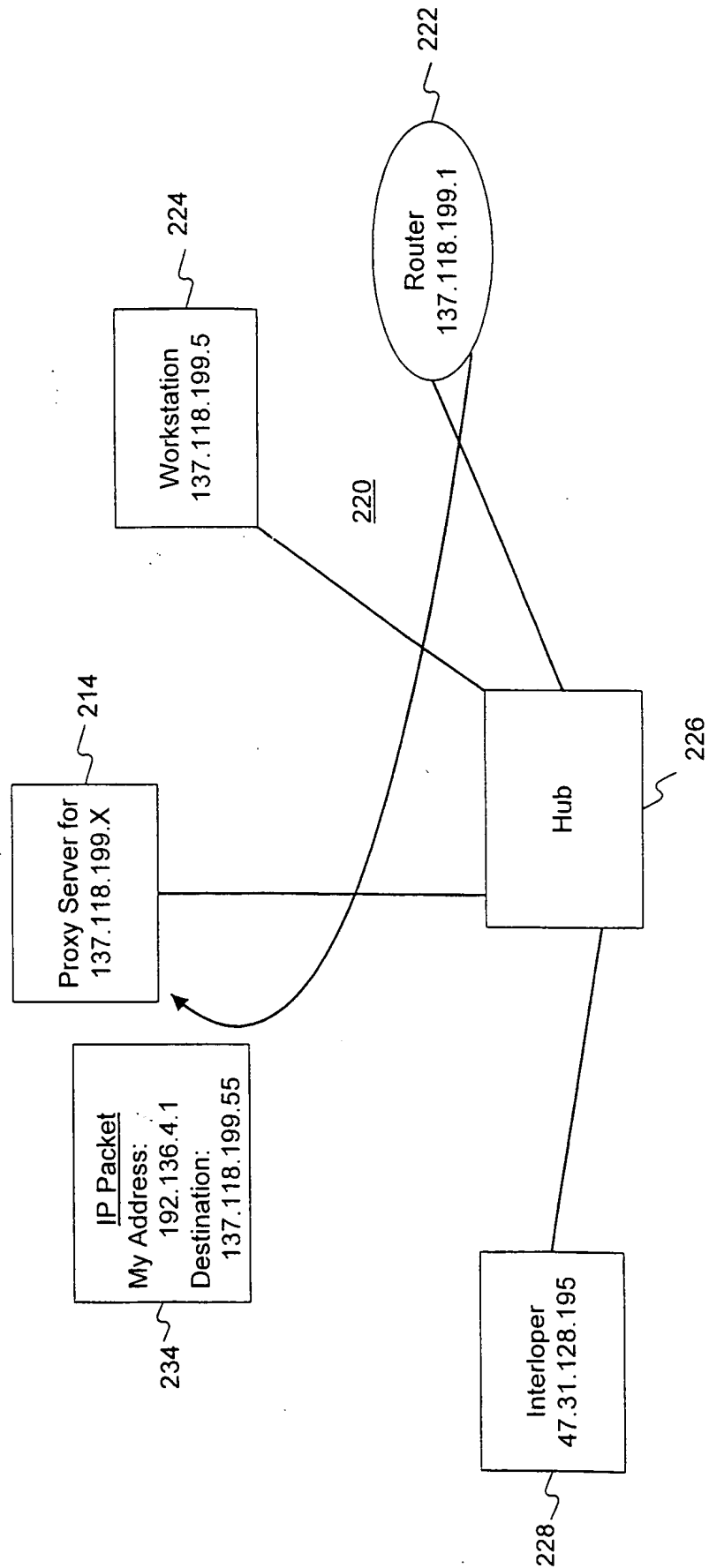


Figure 21

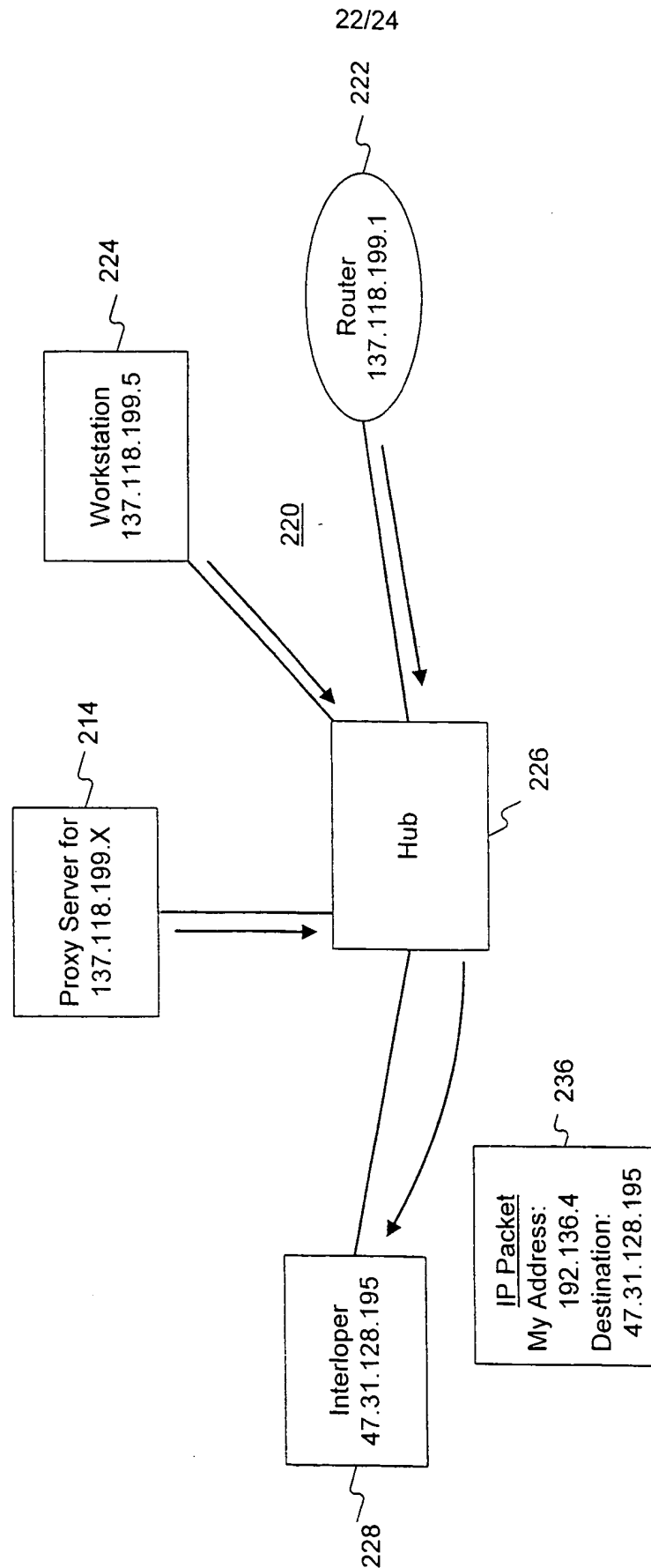


Figure 22

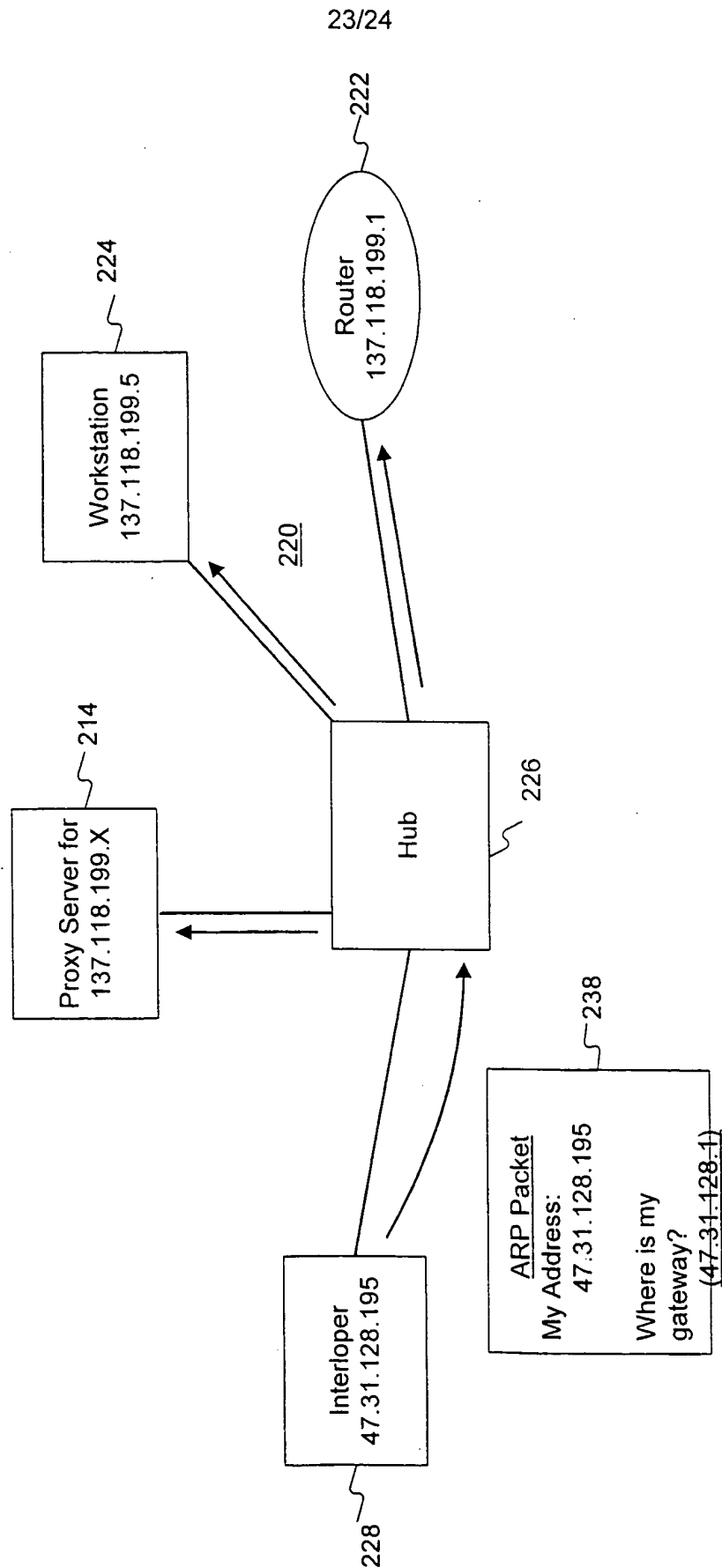


Figure 23

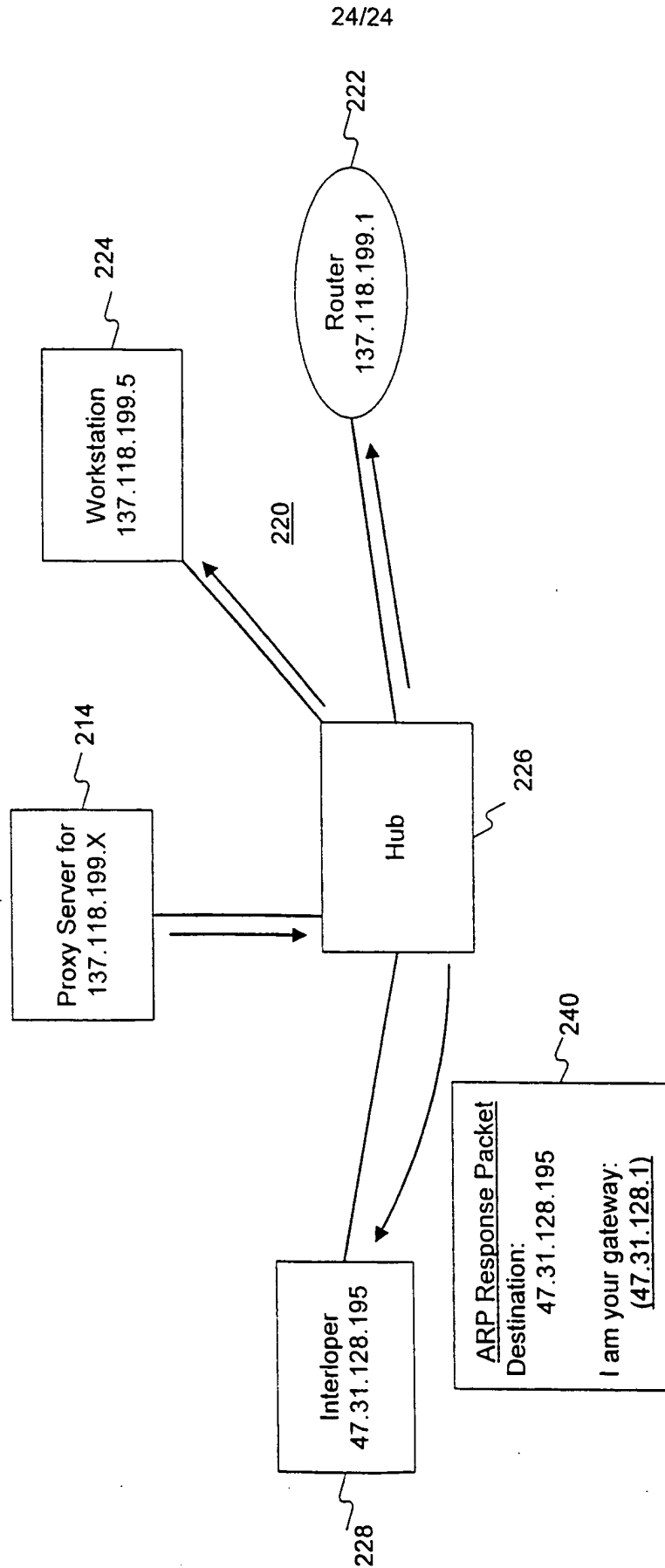


Figure 24

INTERNATIONAL SEARCH REPORT

Inter. Application No
PCT/US 99/01195

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 H04L29/06 H04L29/12

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	GIOVANARDI A ET AL: "TRANSPARENT MOBILE IP: AN APPROACH AND IMPLEMENTATION" IEEE GLOBAL TELECOMMUNICATIONS CONFERENCE, PHOENIX, ARIZONA, NOV. 3 - 8, 1997, vol. 3, 3 November 1997, pages 1861-1865, XP000737840	1-16
Y	INSTITUTE OF ELECTRICAL AND ELECTRONICS ENGINEERS see page 1861, left-hand column, line 1 - page 1862, right-hand column, line 25	17-23
P, X	WO 98 40990 A (KLEINROCK LEONARD ;NOMADIX LLC (US); SHORT JOEL E (US)) 17 September 1998 see page 5, line 25 - page 6, line 21 see page 12, line 1 - page 13, line 26 see page 14, line 28 - page 16, line 12 --- -/--	1-23

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

* Special categories of cited documents:

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier document but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

"Z" document member of the same patent family

Date of the actual completion of the international search

4 June 1999

Date of mailing of the international search report

15/06/1999

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2
NL - 2280 HV Rijswijk
Tel. (+31-70) 340-2040, Tx. 31 651 ep8 nl,
Fax: (+31-70) 340-3016

Authorized officer

RAMIREZ DE AREL..., F

INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 99/01195

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y	GROOM F M: "THE STRUCTURE AND SOFTWARE OF THE INTERNET" ANNUAL REVIEW OF COMMUNICATIONS, vol. 50, 1 January 1997, pages 695-707, XP000720937 see page 698, left-hand column, line 4 - right-hand column, line 7 ----	17-21
Y	YEOM H Y ET AL: "IP MULTIPLEXING BY TRANSPARENT PORT-ADDRESS TRANSLATOR" PROCEEDINGS OF THE SYSTEMS ADMINISTRATION CONFERENCE. LISA, 29 September 1996, pages 113-121, XP002046289 see page 114, left-hand column, line 3 - page 116, left-hand column, line 24 -----	22,23

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/US 99/01195

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
WO 9840990 A	17-09-1998	AU 6698498 A	29-09-1998
<p>This Page Blank (uspto)</p> <p>This Page Blank (uspto)</p>			

This Page Blank (uspto)

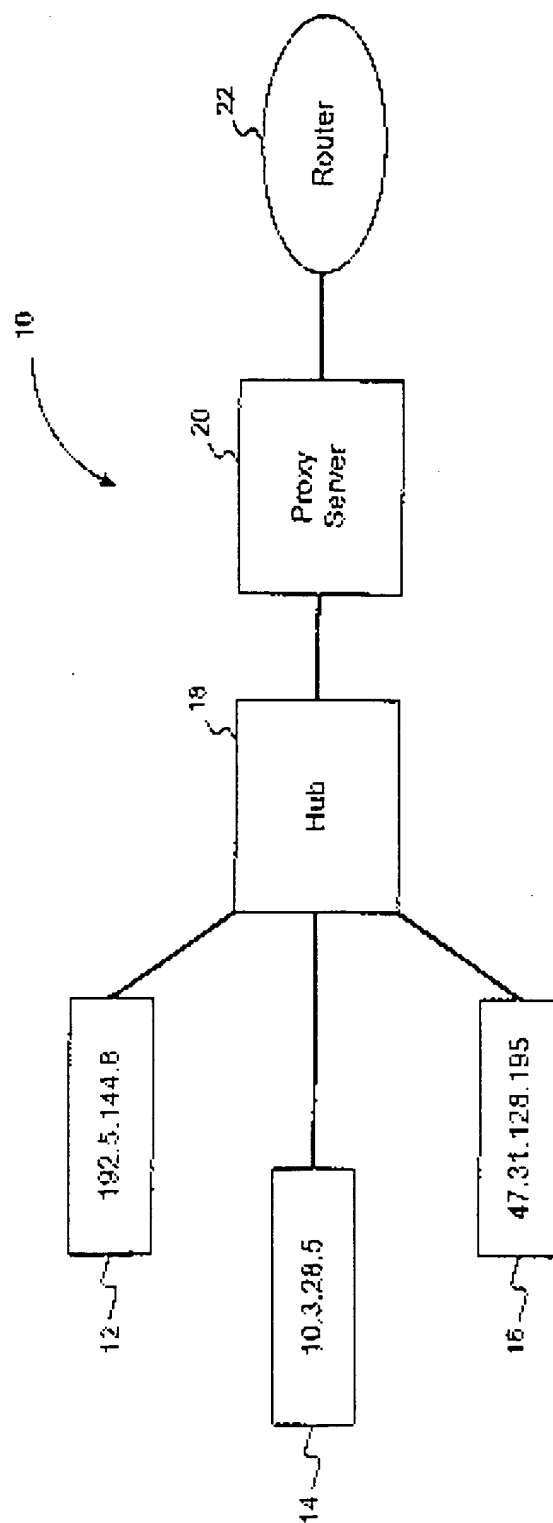


Figure 1

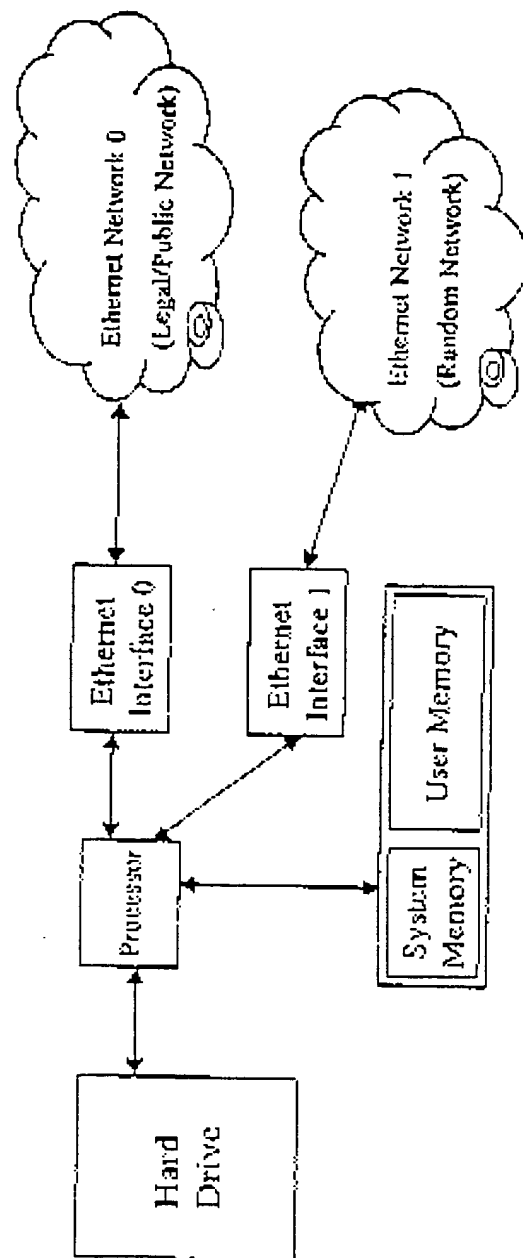


Figure 2

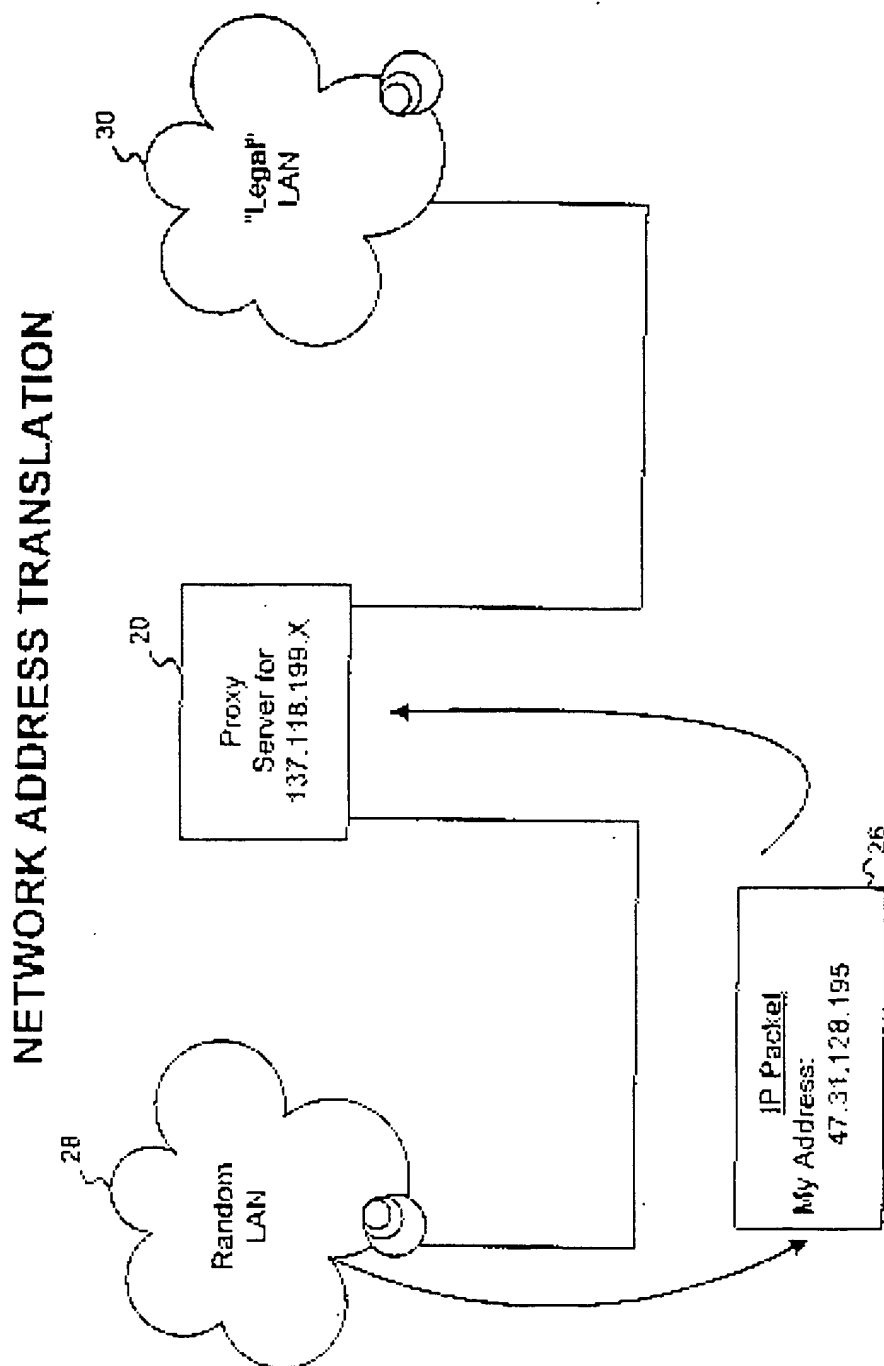


Figure 3

NETWORK ADDRESS TRANSLATION

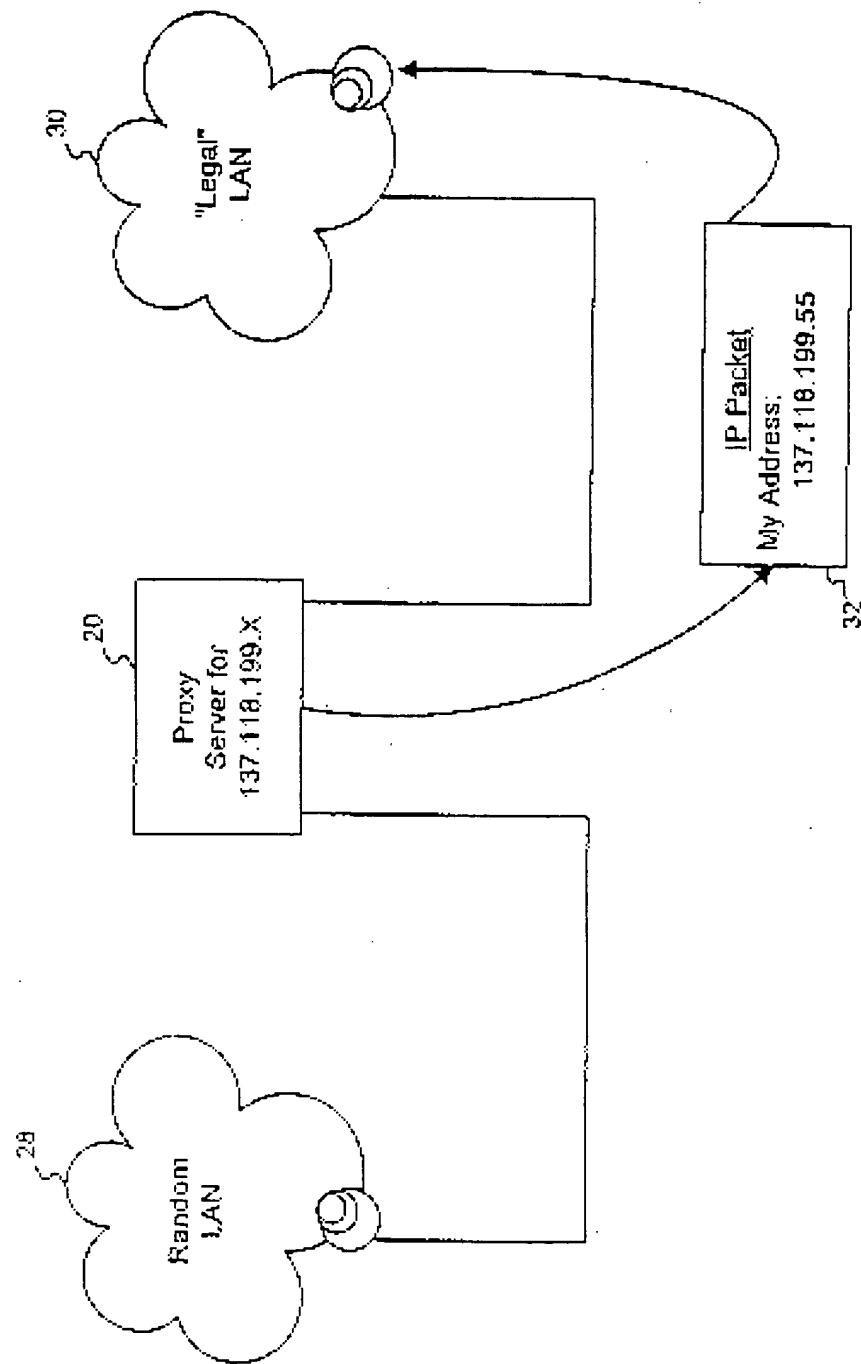


Figure 4

NETWORK ADDRESS TRANSLATION

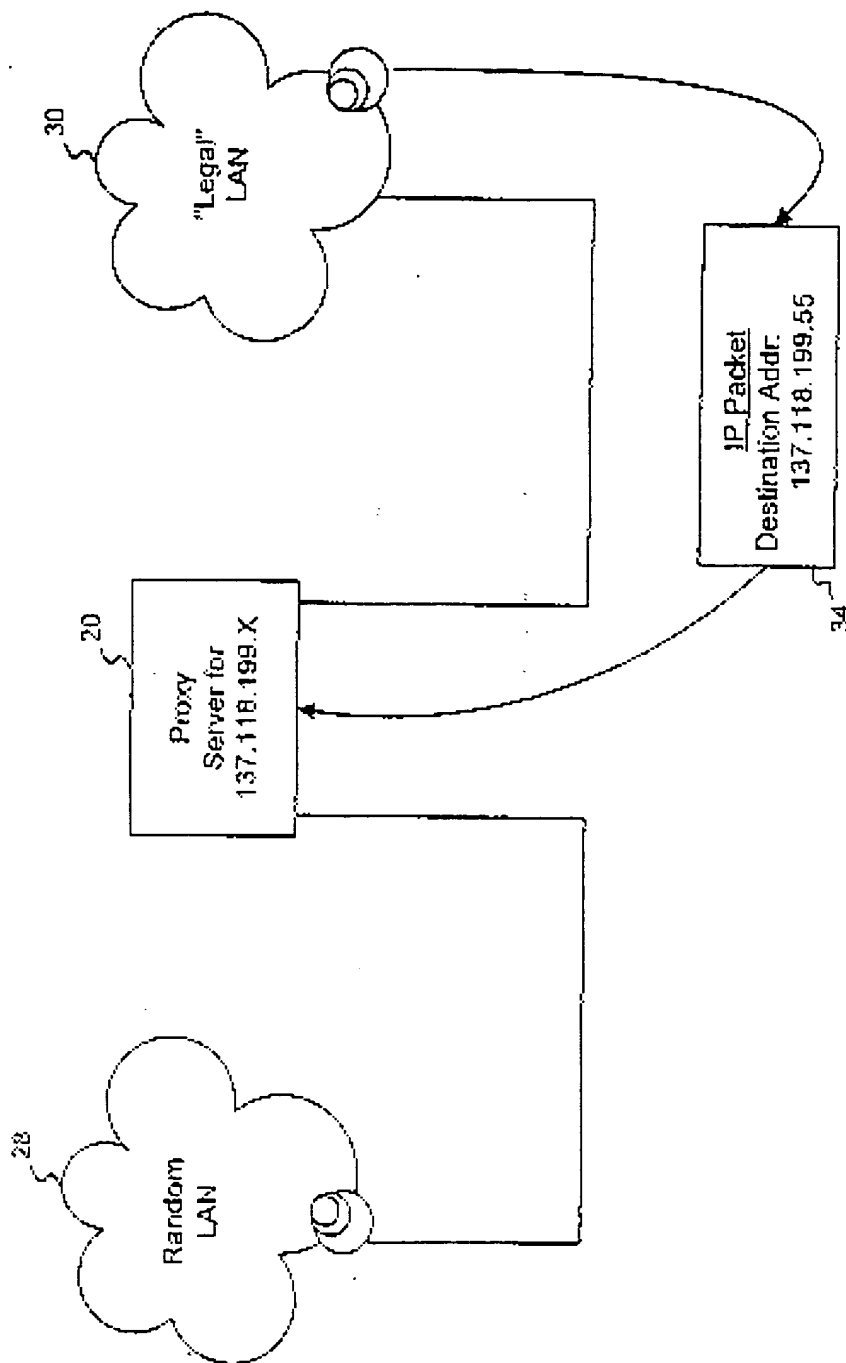


Figure 5

NETWORK ADDRESS TRANSLATION

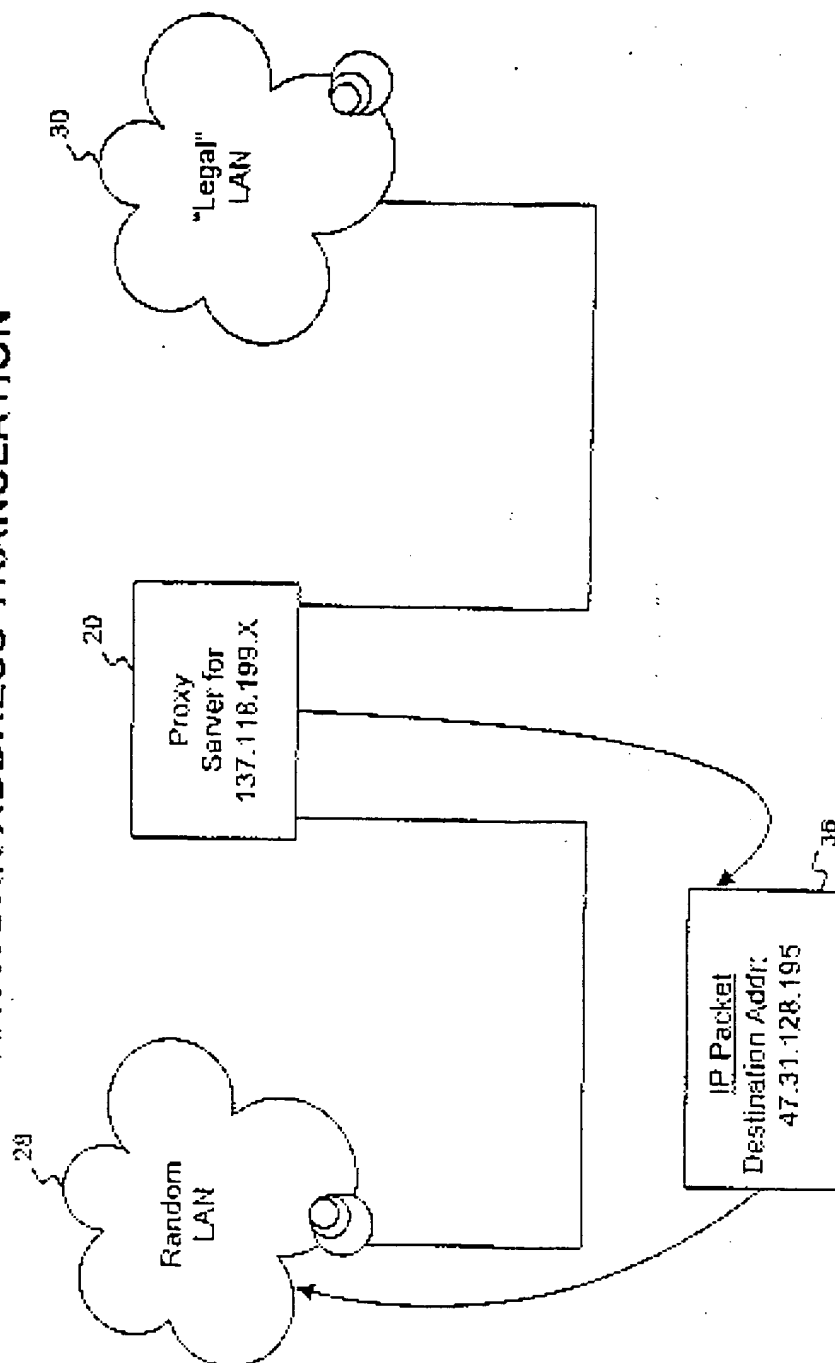


Figure 6

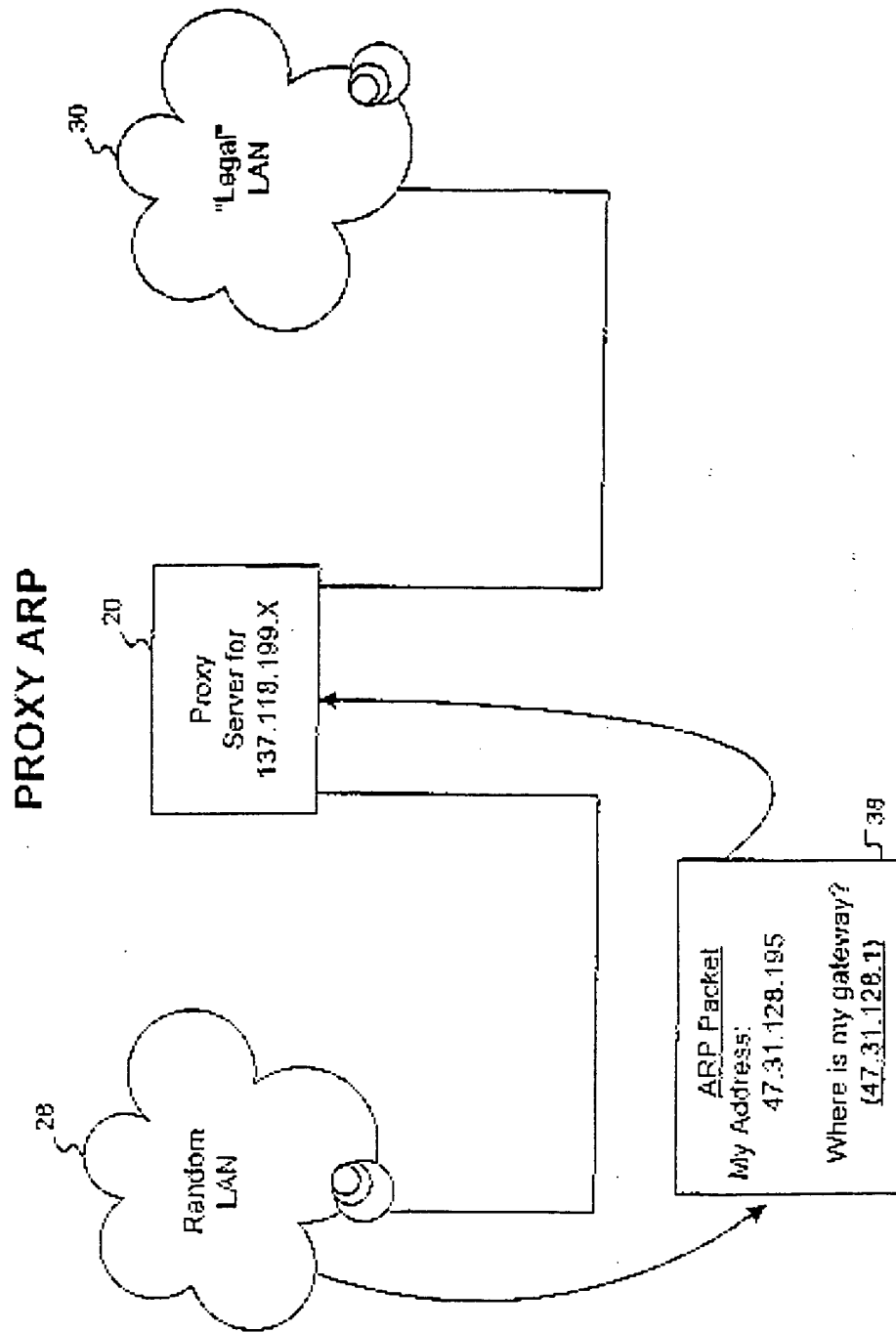


Figure 7

PROXY ARP

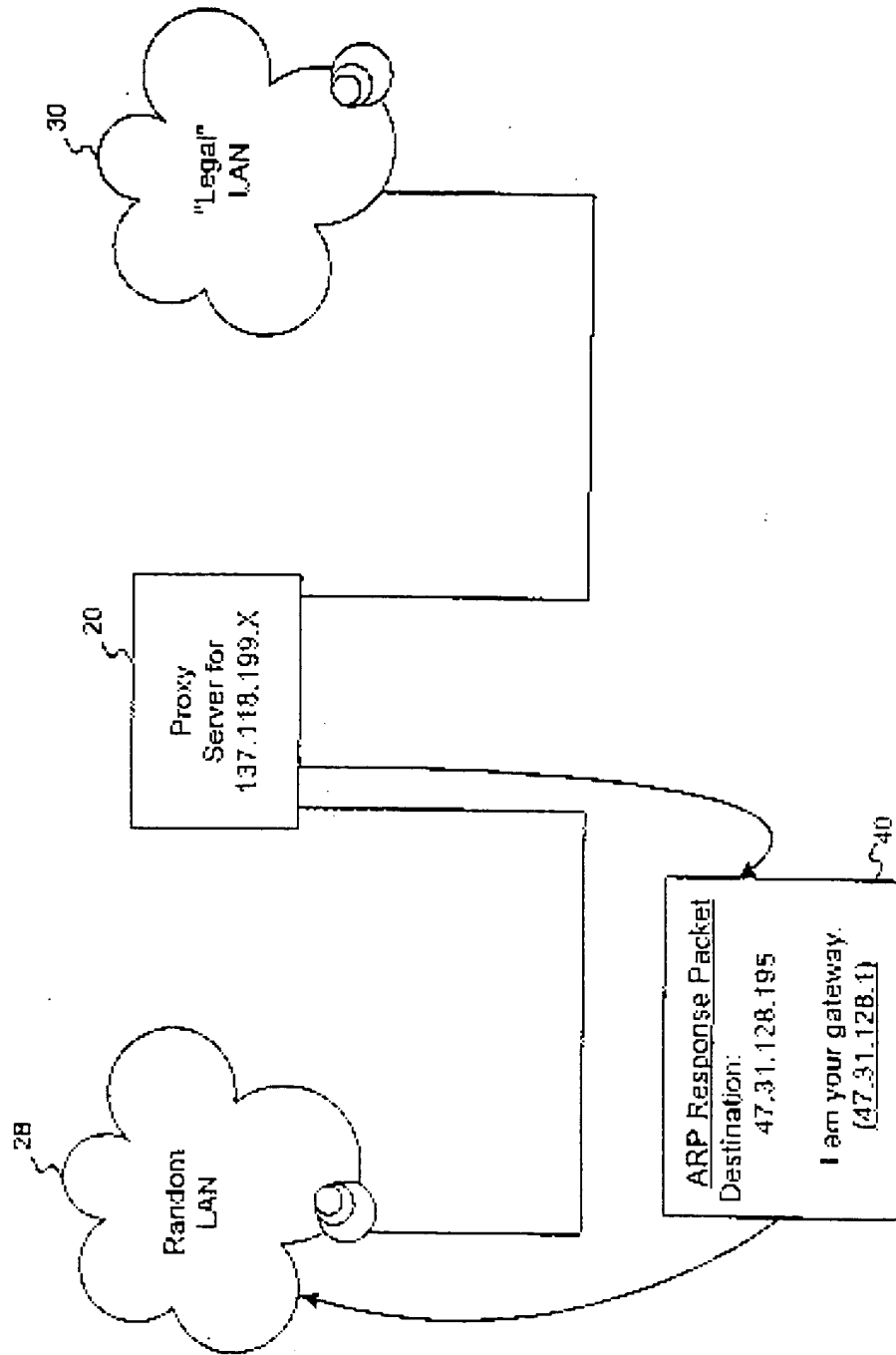


Figure 8

PROXY ARP

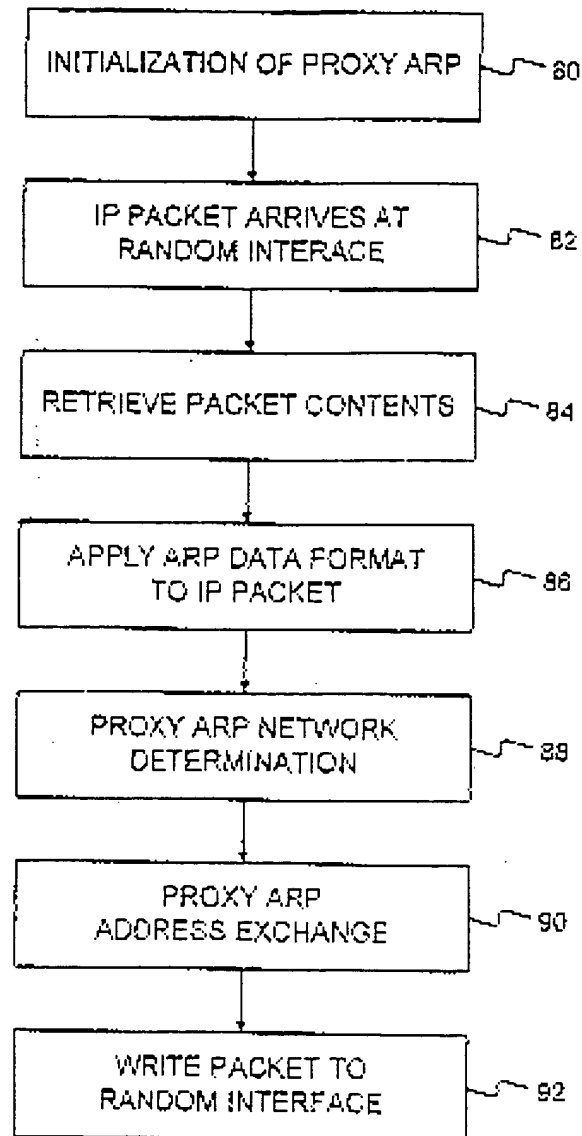


Figure 9

PROXY ARP INITIALIZATION

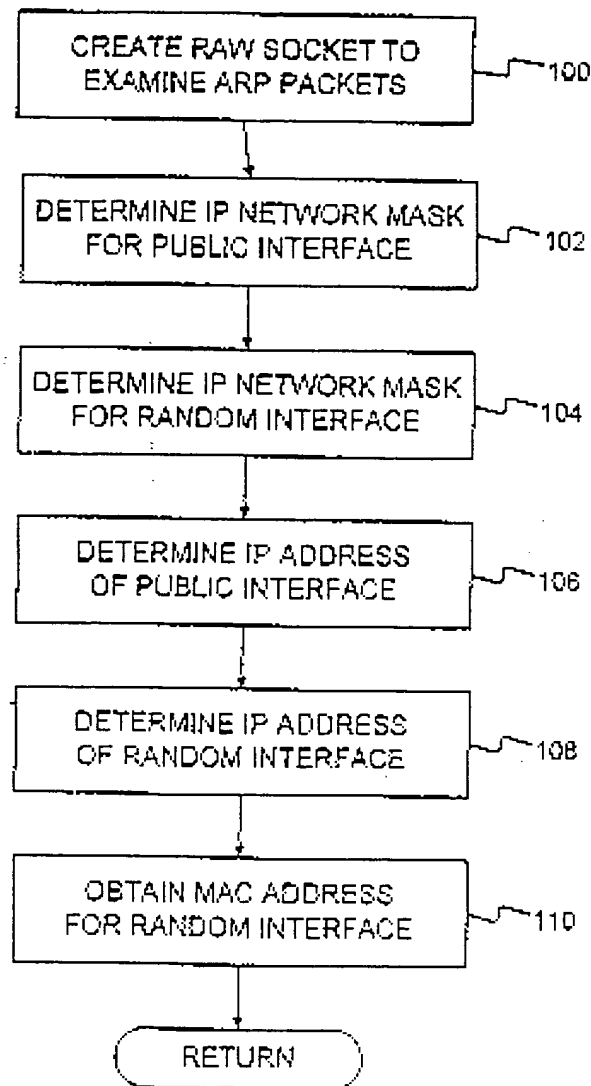


Figure 10

NETWORK DETERMINATION

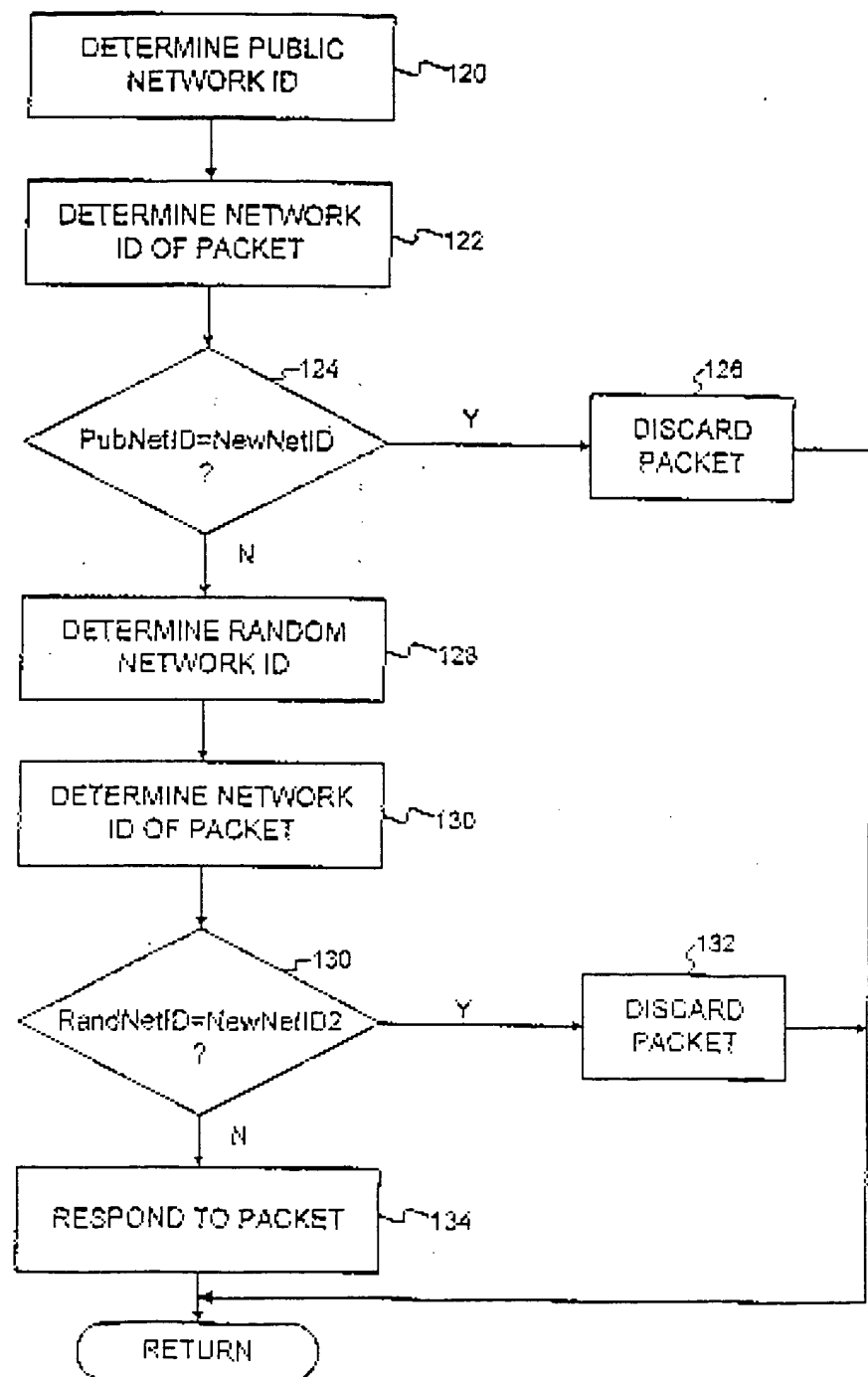


Figure 11

ADDRESS EXCHANGE

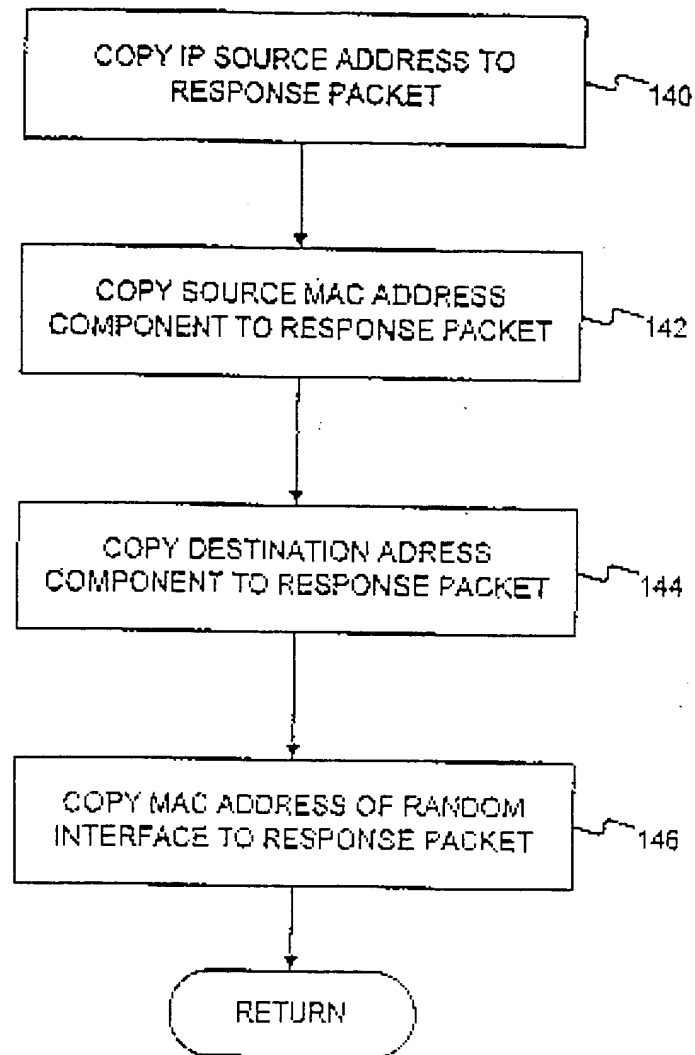


Figure 12

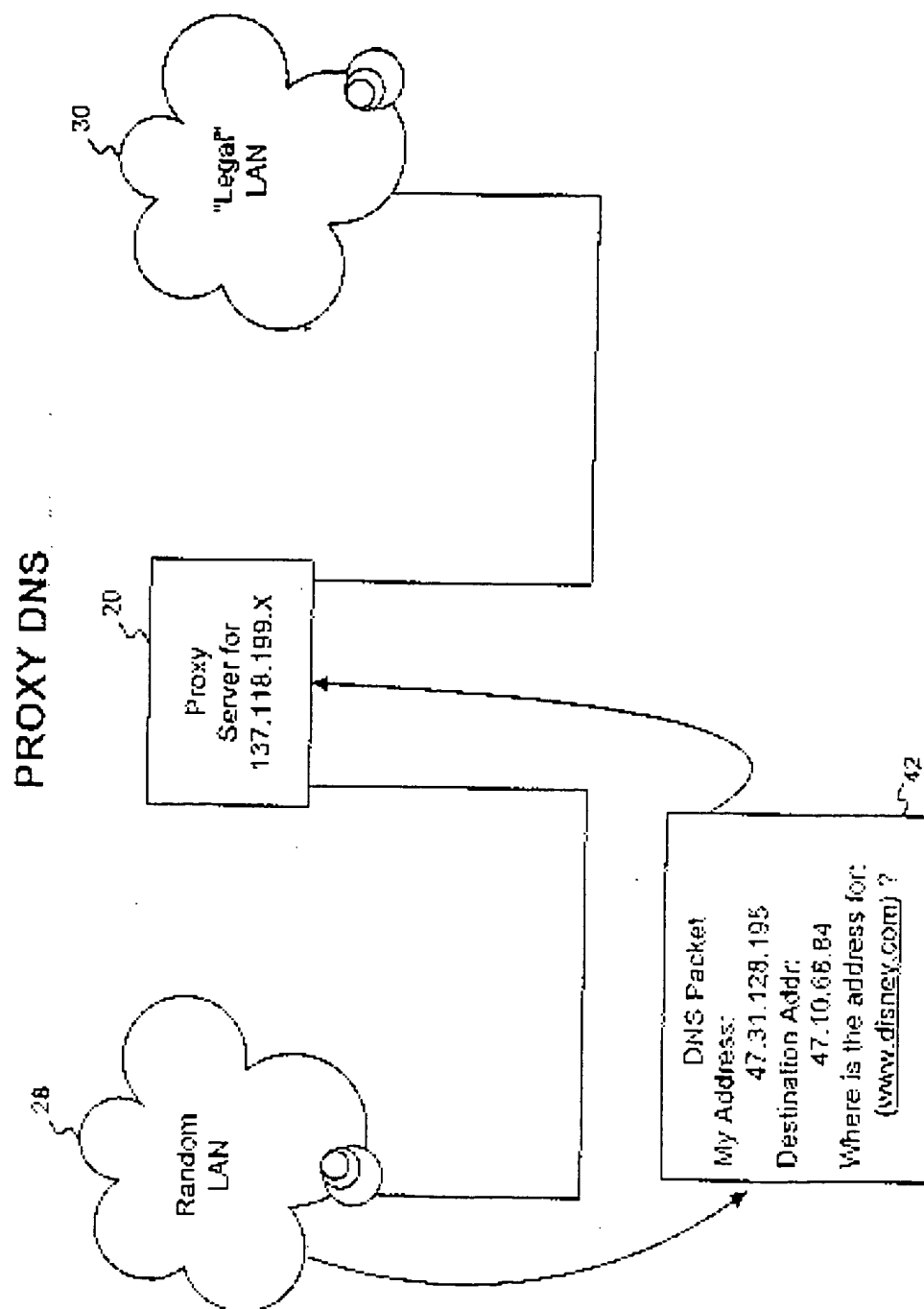


Figure 13

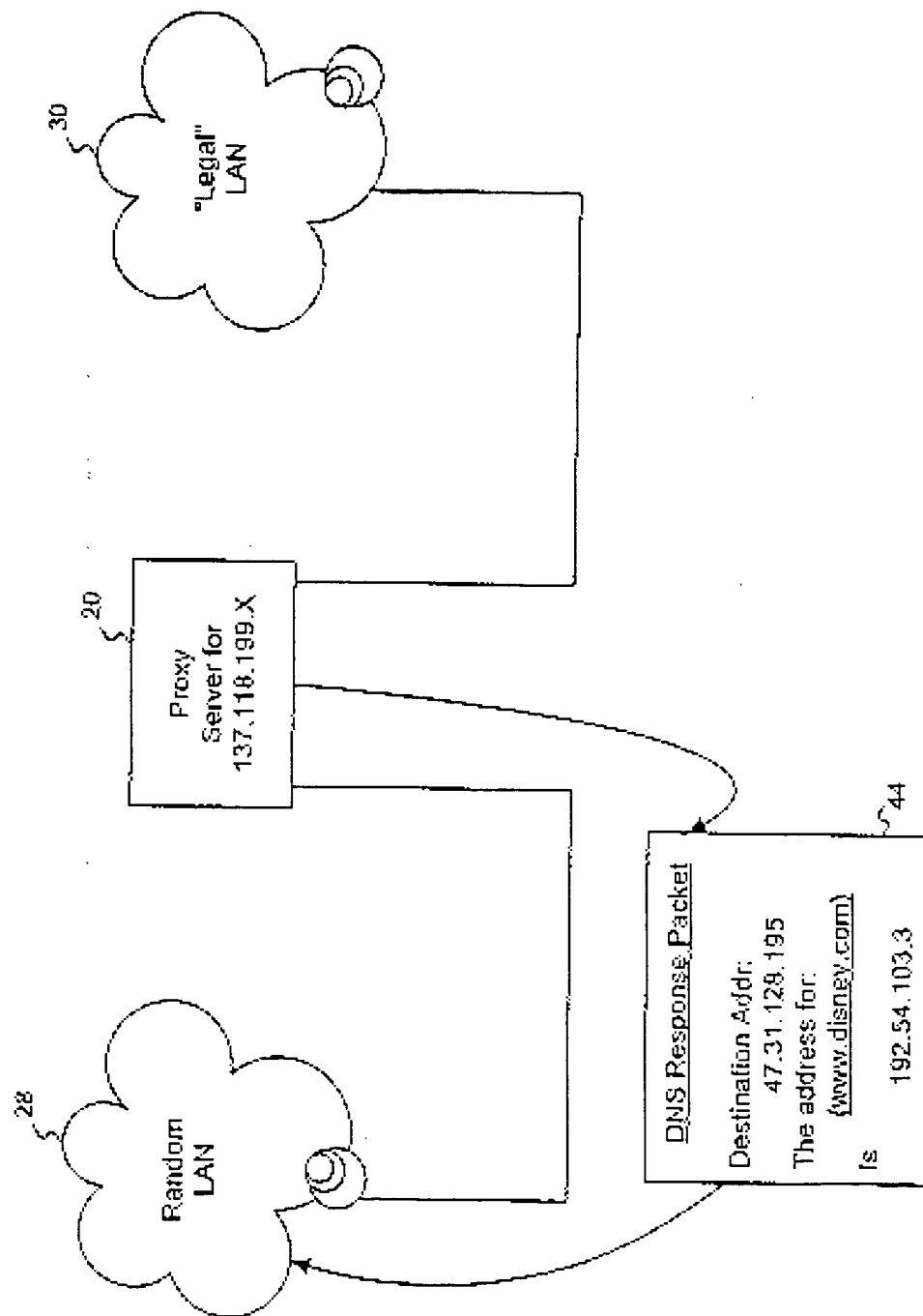


Figure 14

PROXY DNS

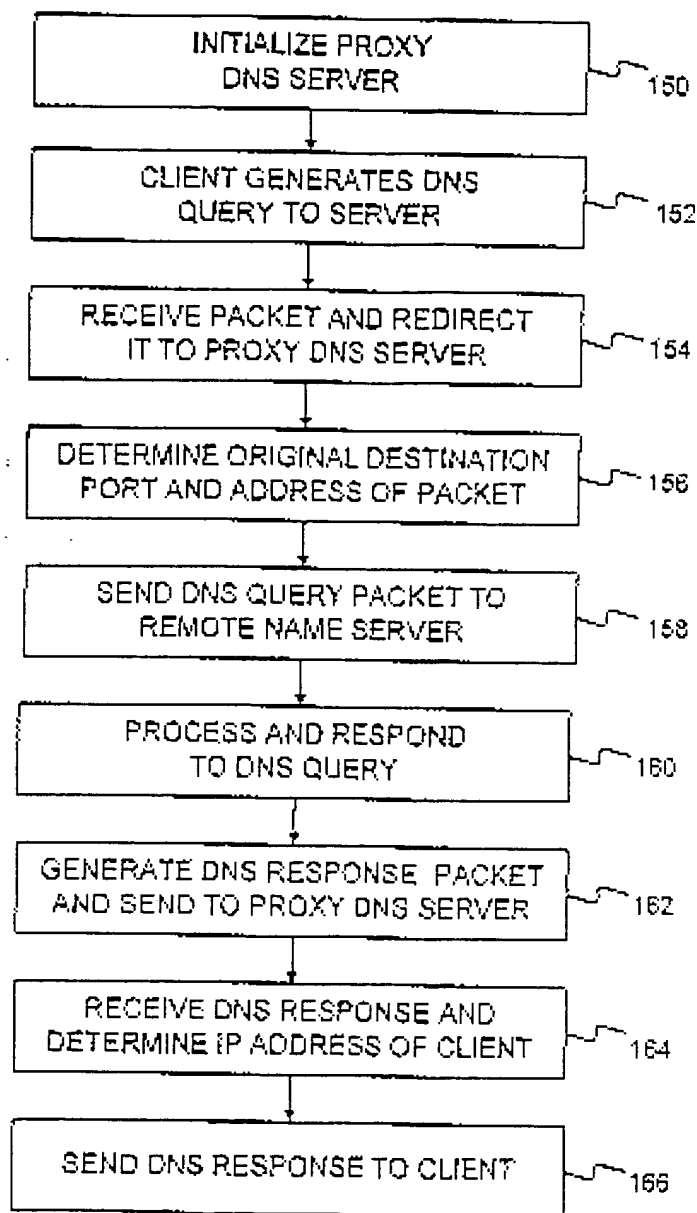


Figure 15

PROXY DNS INITIALIZATION

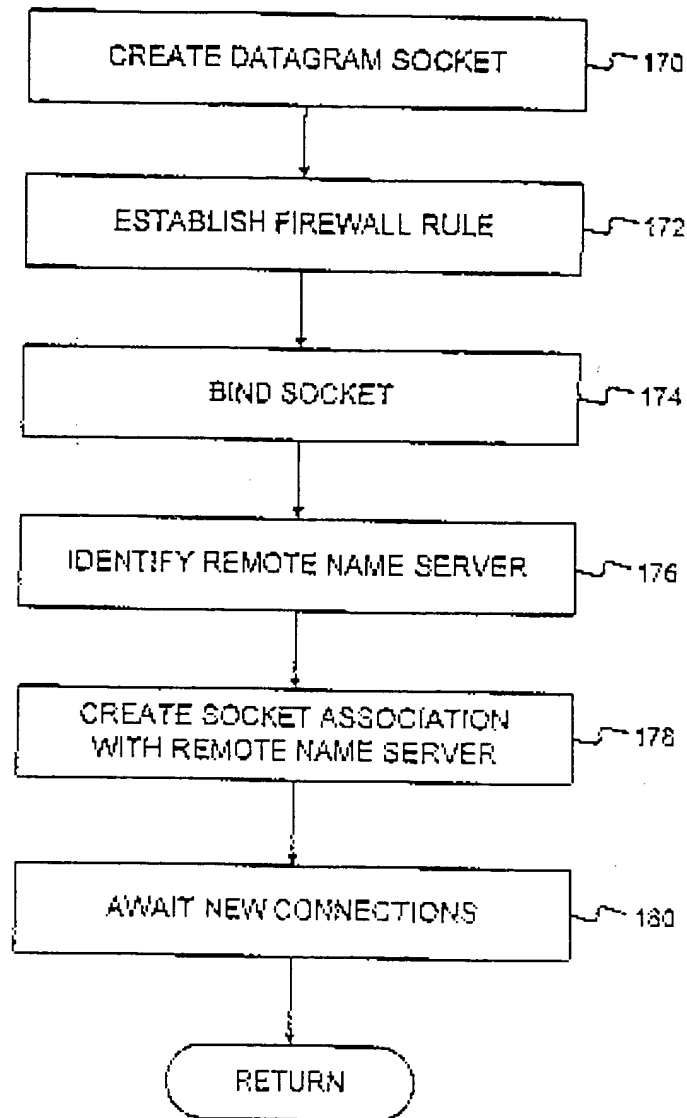


Figure 16

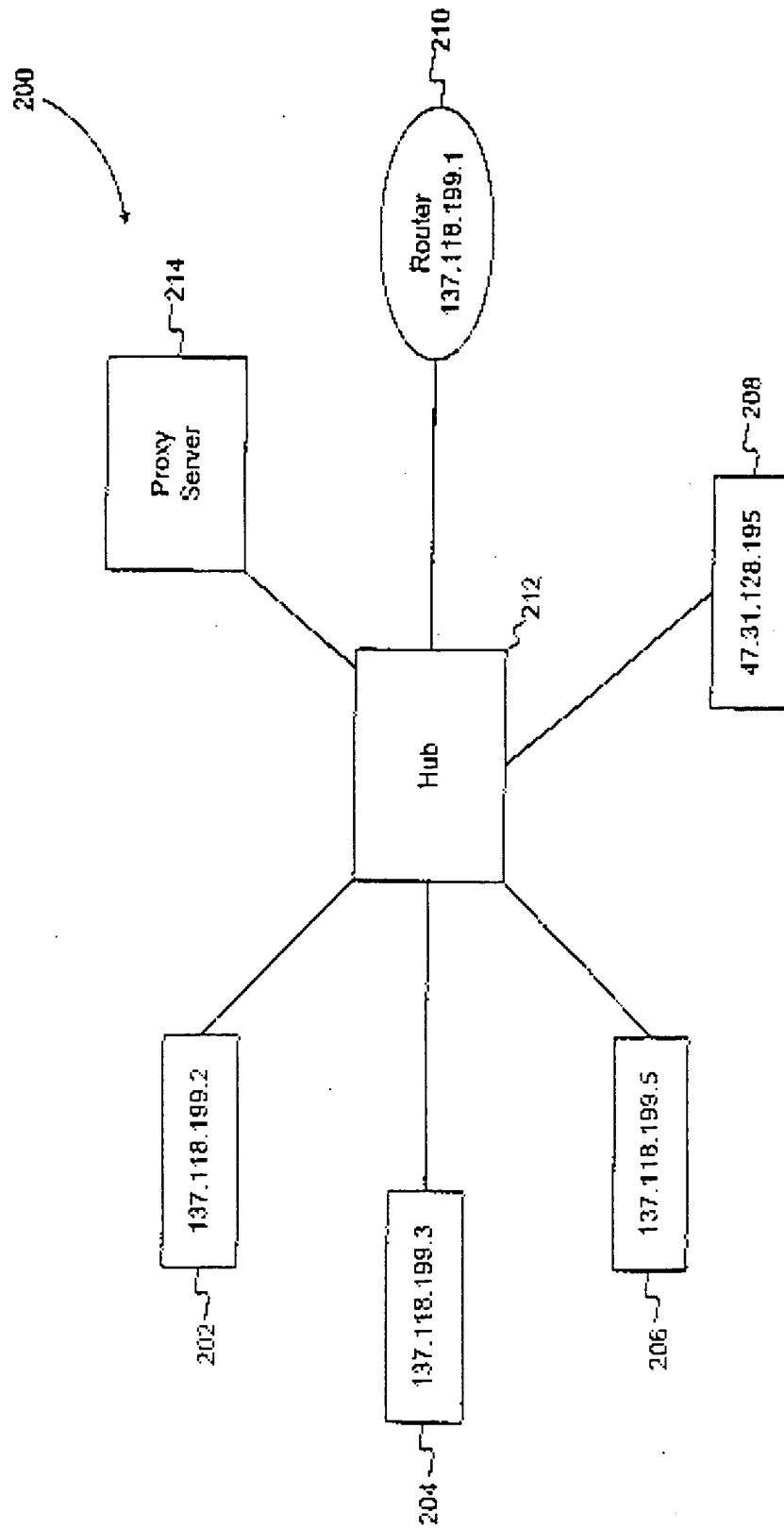


Figure 17

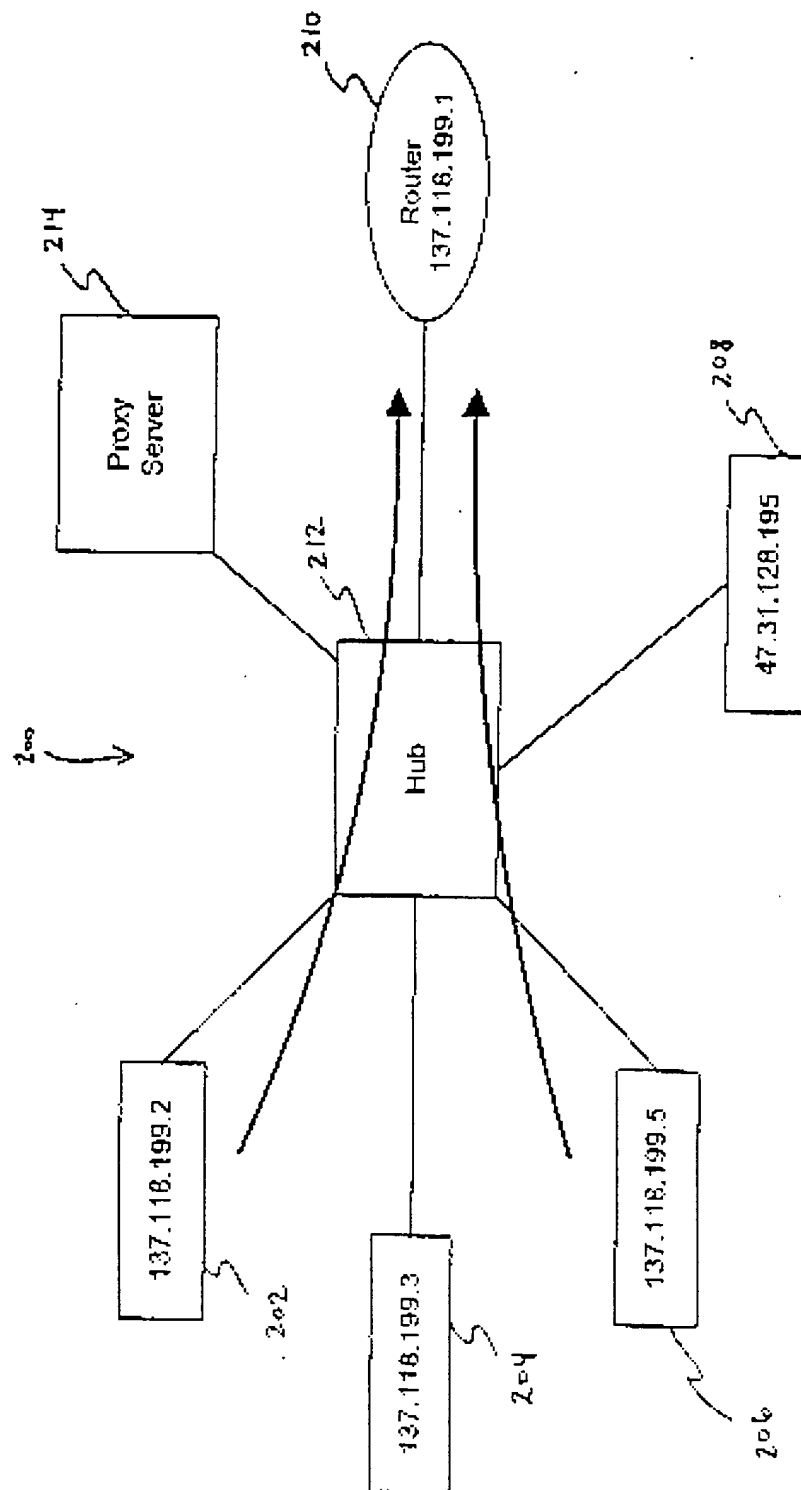


Figure 18

NETWORK ADDRESS TRANSLATION

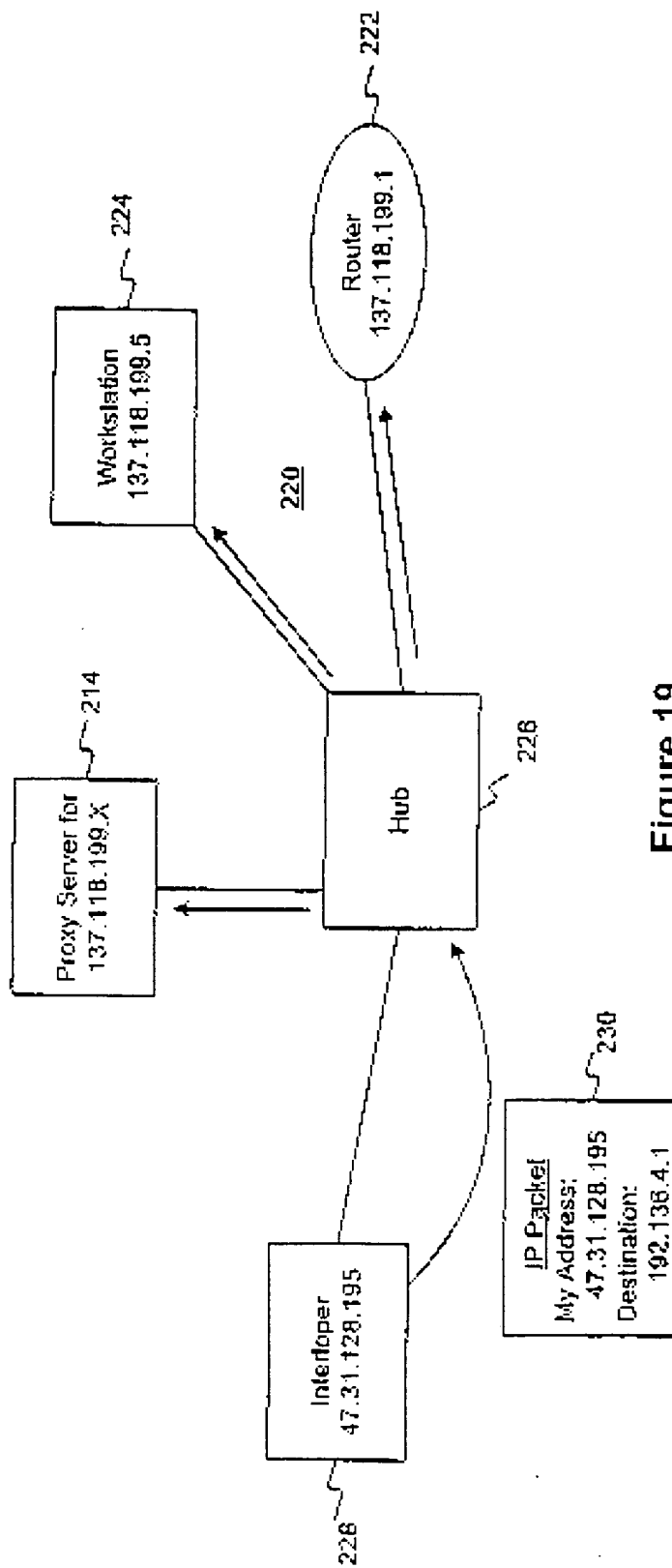


Figure 19

NETWORK ADDRESS TRANSLATION

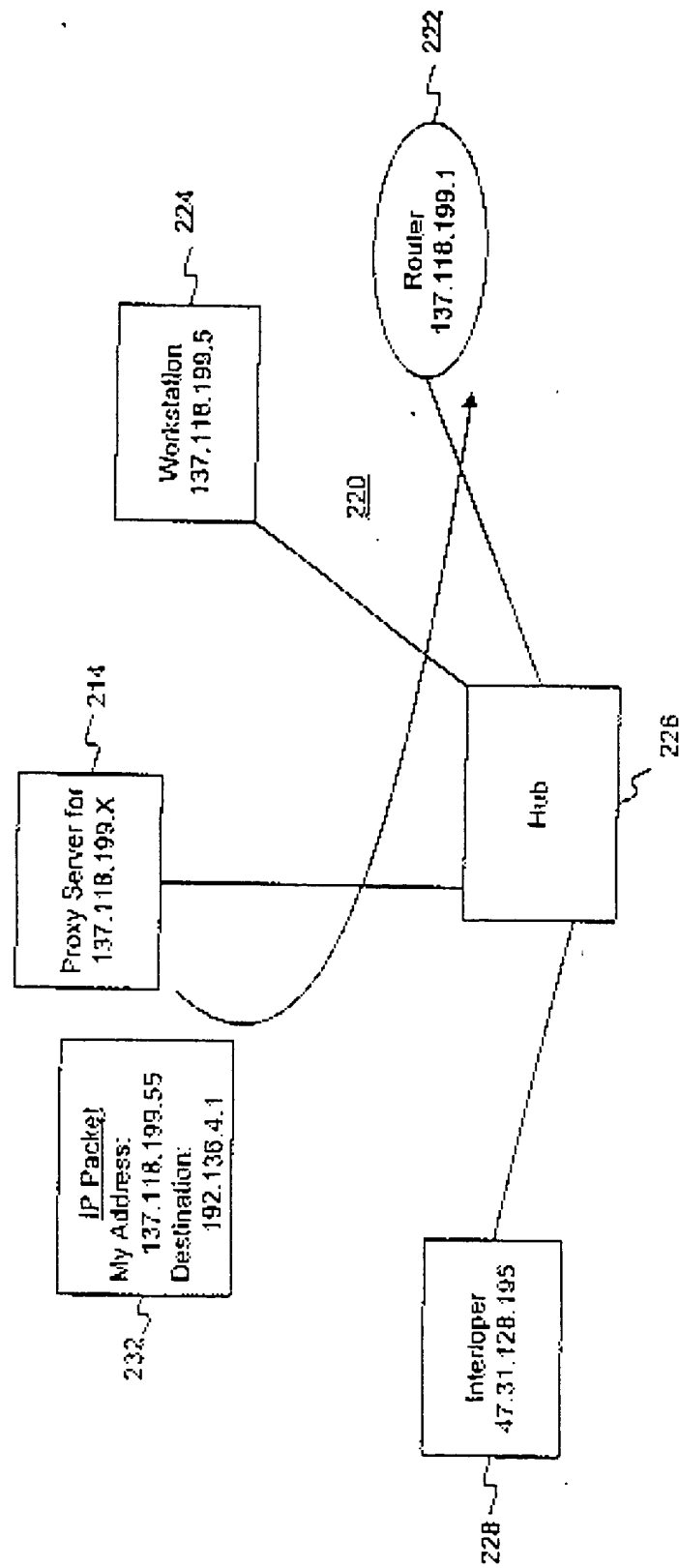


Figure 20

NETWORK ADDRESS TRANSLATION

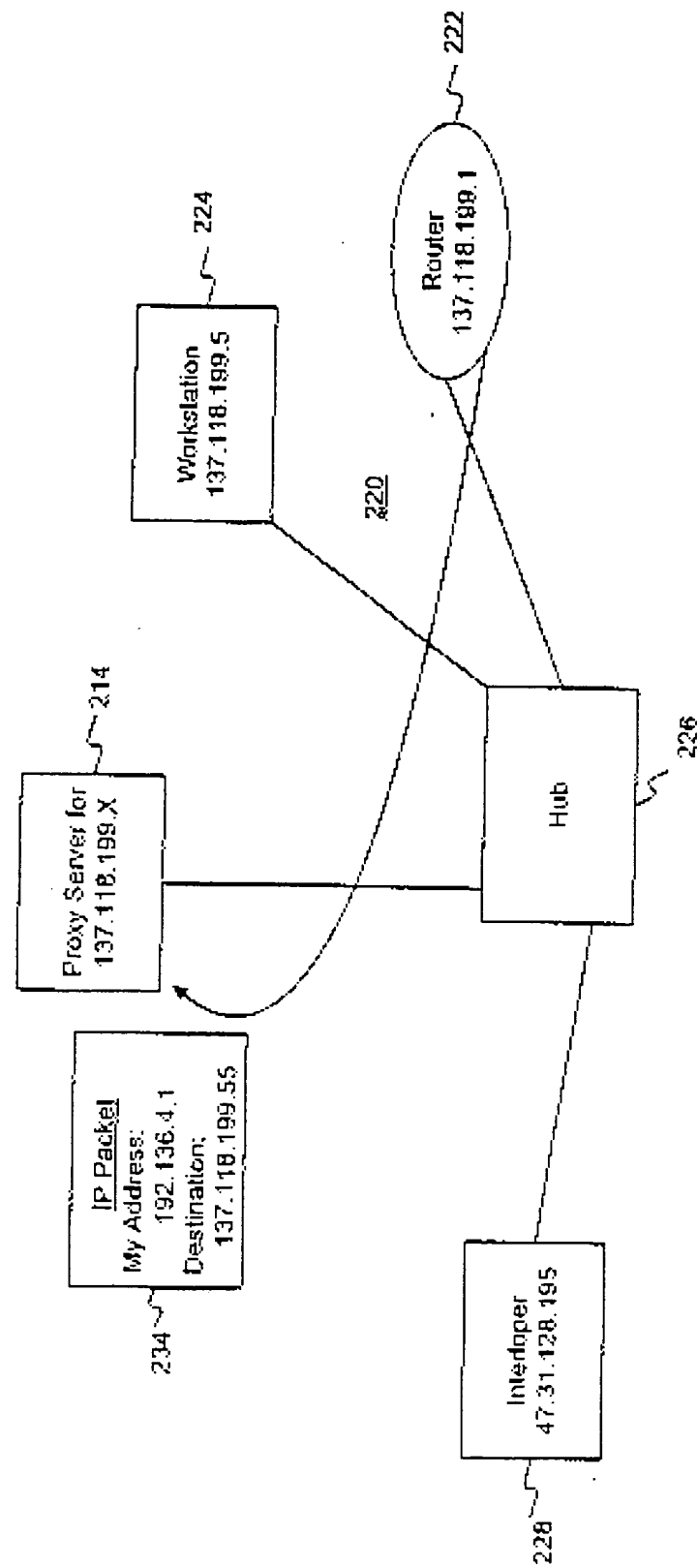


Figure 21

NETWORK ADDRESS TRANSLATION

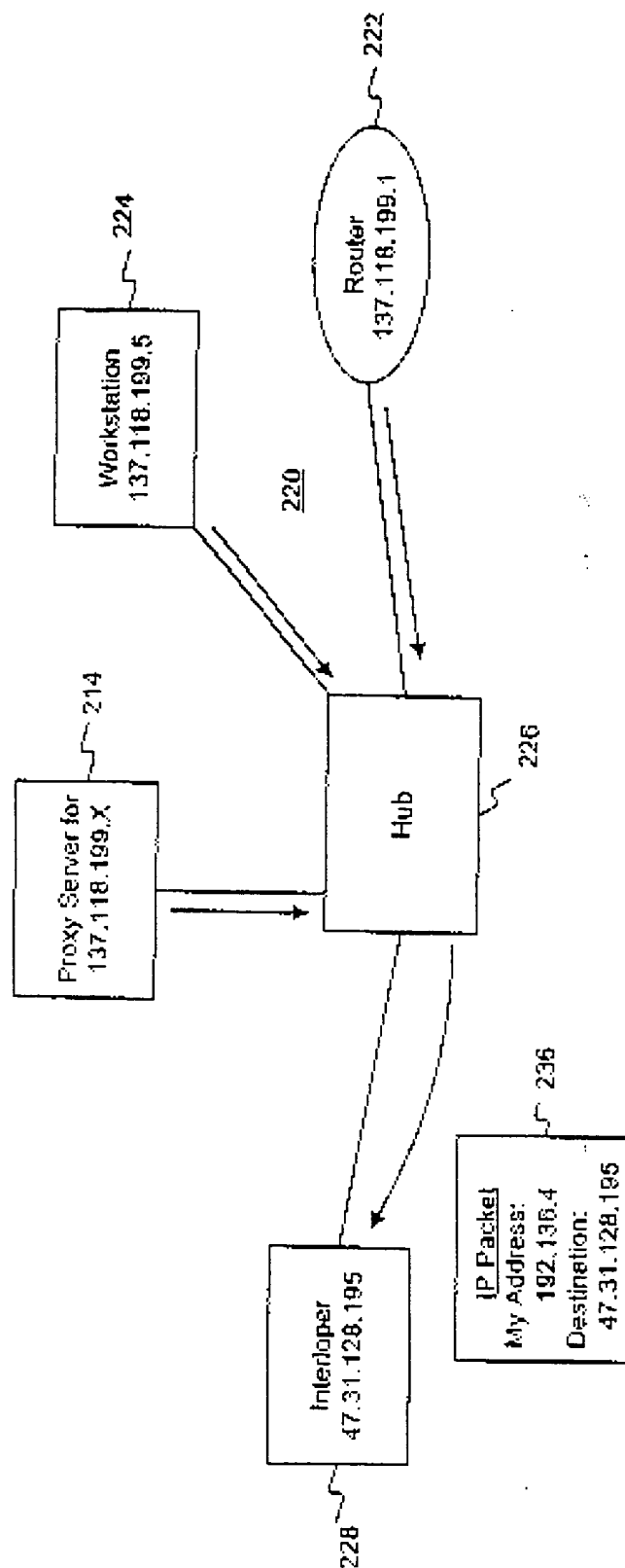


Figure 22

PROXY ARP

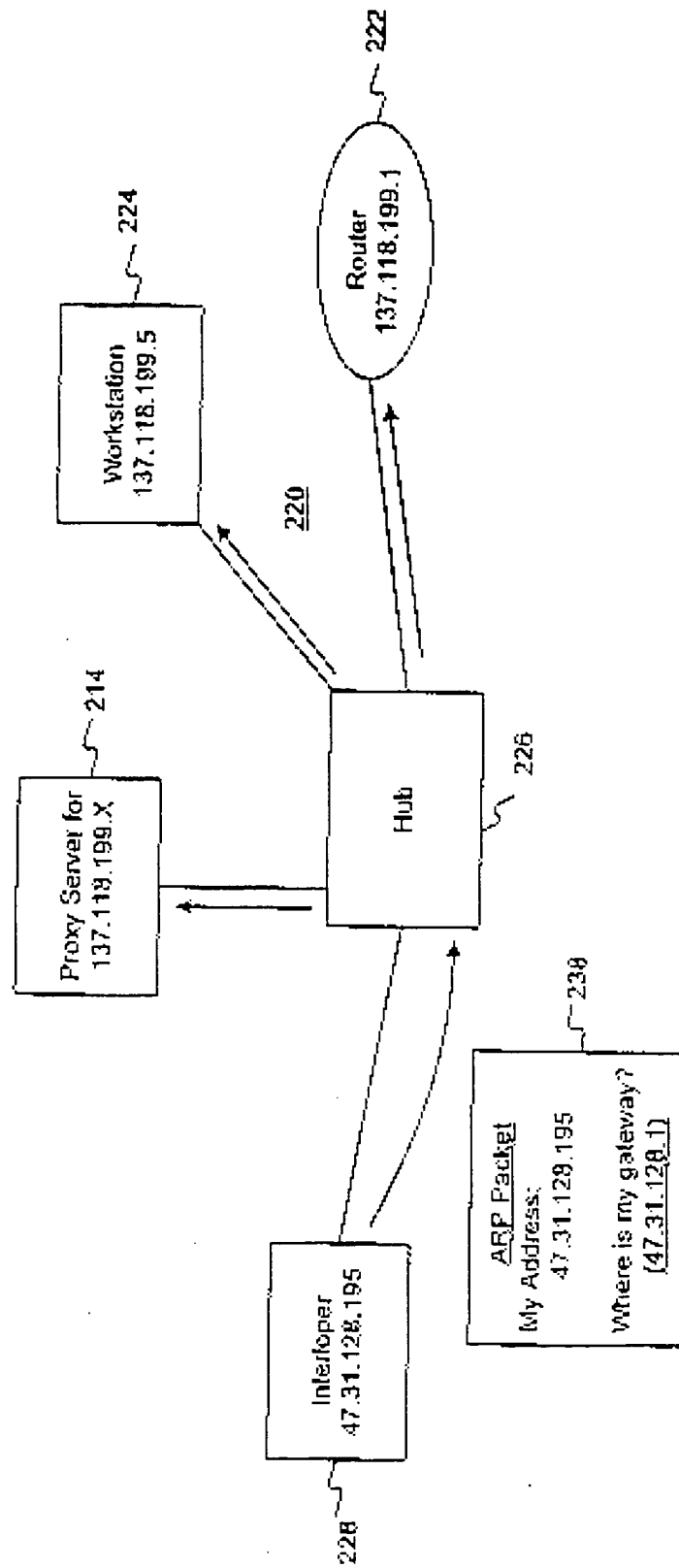


Figure 23

PROXY ARP

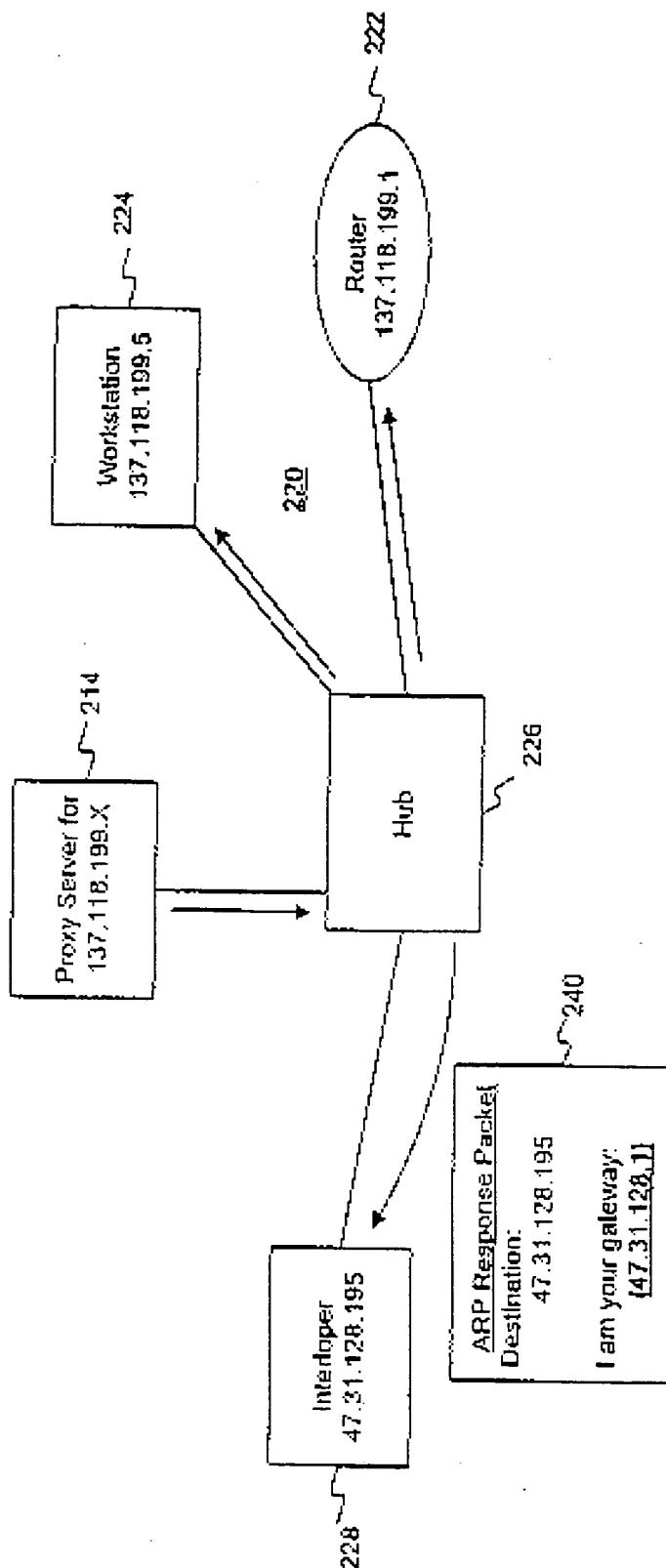


Figure 24

**This Page is Inserted by IFW Indexing and Scanning
Operations and is not part of the Official Record**

BEST AVAILABLE IMAGES

Defective images within this document are accurate representations of the original documents submitted by the applicant.

Defects in the images include but are not limited to the items checked:

- ☐ BLACK BORDERS
- ☐ IMAGE CUT OFF AT TOP, BOTTOM OR SIDES
- ☒ FADED TEXT OR DRAWING
- ☒ BLURRED OR ILLEGIBLE TEXT OR DRAWING
- ☐ SKEWED/SLANTED IMAGES
- ☐ COLOR OR BLACK AND WHITE PHOTOGRAPHS
- ☐ GRAY SCALE DOCUMENTS
- ☒ LINES OR MARKS ON ORIGINAL DOCUMENT
- ☐ REFERENCE(S) OR EXHIBIT(S) SUBMITTED ARE POOR QUALITY
- ☐ OTHER: _____

IMAGES ARE BEST AVAILABLE COPY.

As rescanning these documents will not correct the image problems checked, please do not report these problems to the IFW Image Problem Mailbox.

This Page Blank (uspto)